Cineca TRAINING High Performance Computing 2017

Introduction to MPI Part II Collective Communications and communicators

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Collective communications



Collective communications is a method of communication which involves all processes in a communicator:

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- •All processes (in a communicator) call the collective function
- •Collective communications will not interfere with point-to-point
- All collective communications are blocking (in MPI 2.0)
- No tags are required
- Receive buffers must match in size (number of bytes)

It's a safe communication mode





Communications involving a group of processes. They are called by all the ranks involved in a communicator (or a group) and are of three types:

- Synchronization (e.g. Barrier)
- Data Movement (e.g. Broadcast or Gather/scatter)
- Global Computation (e.g. reductions)

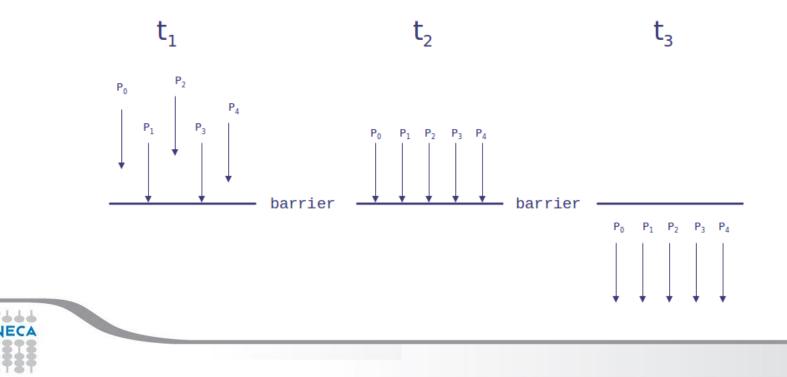




MPI Barrier

It stops all processes within a communicator until they are synchronized

int MPI_Barrier(MPI_Comm comm);

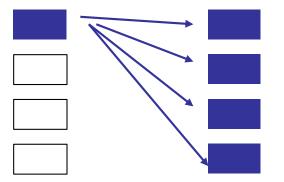


MPI Broadcast

Int MPI_Bcast (void *buf, int count, MPI_Datatype datatype, int root, MPI_Comm comm)

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Note that all processes must specify the same root and same comm.





Example

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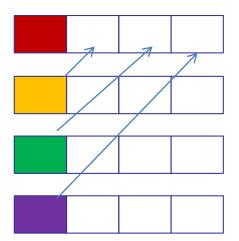
```
PROGRAM broad cast
INCLUDE 'mpif.h'
INTEGER ierr, myid, nproc, root
INTEGER status (MPI STATUS SIZE)
REAL A(2)
CALL MPI INIT(ierr)
CALL MPI COMM SIZE (MPI COMM WORLD, nproc, ierr)
CALL MPI COMM RANK (MPI COMM WORLD, myid, ierr)
root = 0
IF ( myid .EQ. 0 ) THEN
      a(1) = 2.0
      a(2) = 4.0
END IF
CALL MPI BCAST(a, 2, MPI REAL, 0, MPI COMM WORLD, ierr)
WRITE (6, \overline{*}) myid, ': a(1) = ', a(1), 'a(\overline{2}) = ', \overline{a(2)}
CALL MPI FINALIZE (ierr)
END PROGRAM broad cast
```

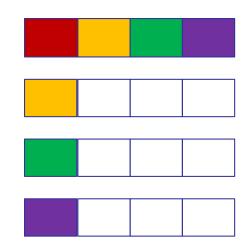
MPI Gather



Each process, root included, sends the content of its send buffer to the root process. The root process receives the messages and stores them in the rank order.

int MPI_Gather(void *sendbuf, int sendcnt, MPI_Datatype sendtype, void *recvbuf, int recvcnt, MPI_Datatype recvtype, int root, MPI_Comm comm)





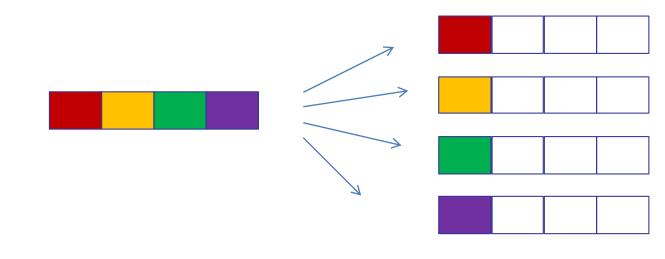


MPI Scatter



The root sends a message. The message is split into n equal segments, the i-th segment is sent to the i-th process in the group and each process receives this message.

int MPI_Scatter(void *sendbuf, int sendcnt, MPI_Datatype sendtype, void *recvbuf, int recvcnt, MPI_Datatype recvtype, int root, MPI_Comm comm)



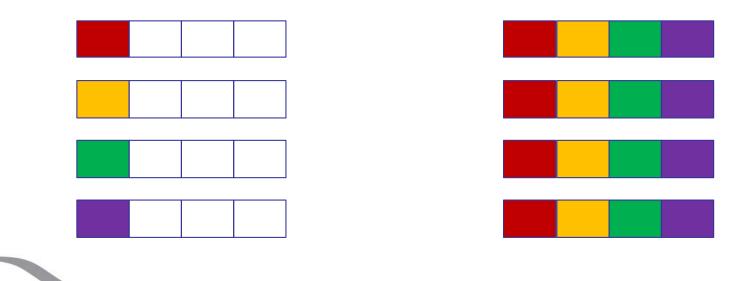




There are possible combinations of collective functions. For example, **MPI Allgather** is a combination of a gather + a broadcast *int MPI_Allgather(void *sendbuf, int sendcount, MPI_Datatype sendtype,*

void *recvbuf, int recvcount, MPI_Datatype recvtype,

MPI_Comm comm)





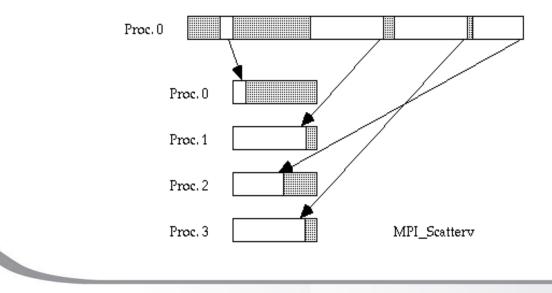


For many collective functions there are extended functionalities.

For example it's possible to define the length of arrays to be scattered or gathered with

MPI_Scatterv

MPI_Gatherv

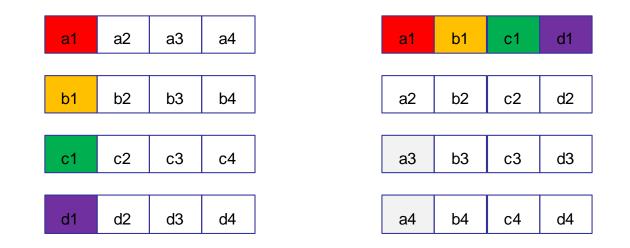


MPI All to all

This function makes a redistribution of the content of each process in a way that each process know the buffer of all others. It is a way to implement the matrix data transposition.

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int MPI_Alltoall(void *sendbuf, int sendcount, MPI_Datatype sendtype, void *recvbuf, int recvcount, MPI_Datatype recvtype, MPI_Comm comm)





Reduction

Reduction operations permits us to

- Collect data from each process
- Reduce the data to a single value
- Store the result on the root process (MPI_Reduce) or

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• Store the result on all processes (MPI_Allreduce)



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Predefined reduction operations

МРІ ор	Function
MPI_MAX	Maximum
MPI_MIN	Minimum
MPI_SUM	Sum
MPI_PROD	Product
MPI_LAND	Logical AND
MPI_BAND	Bitwise AND
MPI_LOR	Logical OR
MPI_BOR	Bitwise OR
MPI_LXOR	Logical exclusive OR
MPI_BXOR	Bitwise exclusive OR
MPI_MAXLOC	Maximum and location
MPI_MINLOC	Minimum and location



```
PROGRAM reduce
INCLUDE 'mpif.h'
INTEGER ierr, myid, nproc, root
REAL A(2), res(2)
CALL MPI INIT(ierr)
CALL MPI COMM SIZE (MPI COMM WORLD, nproc, ierr)
CALL MPI COMM RANK (MPI COMM WORLD, myid, ierr)
root = 0
a(1) = 2.0
a(2) = 4.0
CALL MPI REDUCE (a, res, 2, MPI REAL, MPI SUM, root,
MPI COMM WORLD, ierr)
IF ( myid .EQ. 0 ) THEN
WRITE (6, *) myid, ': res(1) = ', res(1), 'res(2) = ', res(2)
END IF
CALL MPI FINALIZE (ierr)
END
```

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Performance issues

Much hidden communication takes place with collective communication.

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- Hardware vendors work hard to provide optimized collective calls but performances will vary according to implementation.
- Because of forced synchronization, collective communications may not always be the best solution.

Some studies show that around 80% transfer time is in collectives.



MPI communicators and groups



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Many users are familiar with the mostly used communicator:

MPI_COMM_WORLD

A communicator can be thought as a handle to a group.

- a group is a ordered set of processes
 - each process is associated with a rank
 - ranks are contiguous and start from zero

Groups allow collective operations to be operated on a subset of processes

The group routines are primarily used to specify which processes should be used to construct a communicator.



Intracommunicators

are used for communications within a single group

Intercommunicators

are used for communications between two disjoint groups



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Group management:

- -All group operations are local
- -Groups are not initially associated with communicators
- -Groups can only be used for message passing within a communicator
- -We can access groups, construct groups, destroy groups, i.e. groups/communicators are dynamic they can be created and destroyed during program execution.



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Typical usage:

- 1. Extract handle of global group from MPI_COMM_WORLD using MPI_Comm_group
- 2. Form new group as a subset of global group using MPI_Group_incl
- 3. Create new communicator for new group using MPI_Comm_create
- 4. Determine new rank in new communicator using MPI_Comm_rank
- 5. Conduct communications using any MPI message passing routine
- 6. When finished, free up new communicator and group (optional) using MPI_Comm_free and MPI_Group_free



Group constructors

Group constructors are used to create new groups from existing ones (initially from the group associated with MPI_COMM_WORLD; you can use mpi_comm_group to get this).

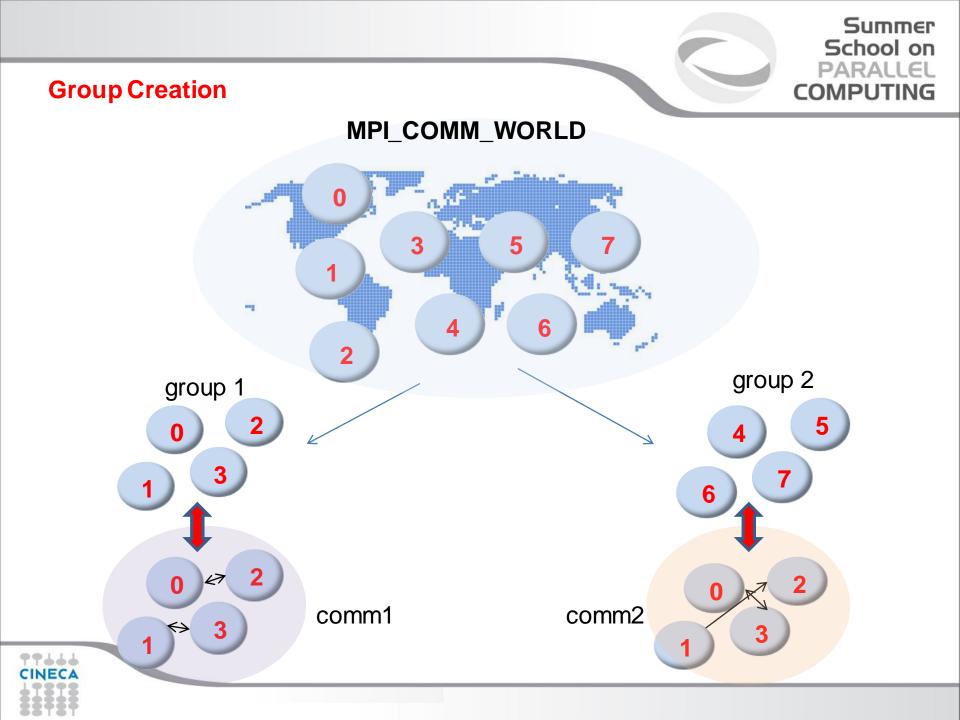
Group creation is a local operation: no communication is needed

After the creation of a group, no communicator has been associated to this group, and hence no communication is possible within the new group



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Group accessors:

-MPI_GROUP_SIZE

This routine returns the number of processes in the group

-MPI_GROUP_RANK

This routine returns the rank of the calling process inside a given group



-MPI_COMM_GROUP(comm,group,ierr)



This routine returns the group associated with the communicator comm

-MPI_GROUP_UNION(group_a, group_b, newgroup, ierr)

This returns the ensemble union of group_a and group_b

-MPI_GROUP_INTERSECTION(group_a, group_b, newgroup, ierr)

This returns the ensemble intersection of group_a and group_b

-MPI_GROUP_DIFFERENCE(group_a, group_b, newgroup, ierr)

This returns in newgroup all processes in group_a that rare not in group_b, ordered as in group_a





-MPI_GROUP_INCL(group, n, ranks, newgroup, ierr)

This routine creates a new group that consists of all the n processes with ranks ranks[0]... ranks[n-1]

Example:

group = $\{a,b,c,d,e,f,g,h,i,j\}$ n = 5 ranks = $\{0,3,8,6,2\}$

newgroup = $\{a,d,i,g,c\}$





-MPI_GROUP_EXCL(group,n,ranks,newgroup,ierr)

This routine returns a newgroup that consists of all the processes in the group after removing processes with ranks: ranks[0]..ranks[n-1]

Example:

group = $\{a,b,c,d,e,f,g,h,i,j\}$ n = 5 ranks = $\{0,3,8,6,2\}$

newgroup = $\{b,e,f,h,j\}$



Communicator management

Communicator access operations are local, not requiring interprocess communication

Communicator constructors are collective and may require interprocess communications

We will cover in depth only intracommunicators, giving only some notions about intercommunicators.



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Communicator accessors

-MPI_COMM_SIZE(comm,size,ierr)

Returns the number of processes in the group associated with the comm

-MPI_COMM_RANK(comm,rank,ierr)

Returns the rank of the calling process within the group associated with the comm

-MPI_COMM_COMPARE(comm1,comm2,result,ierr)

Returns:

- MPI_IDENT if comm1 and comm2 are the same handle
- MPI_CONGRUENT if comm1 and comm2 have the same group attribute

- MPI_SIMILAR if the groups associated with comm1 and comm2 have the same members but in different rank order

- MPI_UNEQUAL otherwise



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Communicator constructors

-**MPI_COMM_DUP**(comm, newcomm, ierr) This returns a communicator newcomm identical to the communicator comm

-MPI_COMM_CREATE(comm, group, newcomm,ierr) This collective routine must be called by all the process involved in the group associated with comm. It returns a new communicator that is associated with the group. MPI_COMM_NULL is returned to processes not in the group. Note that group must be a subset of the group associated with comm!



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A practical example:

CALL MPI_COMM_RANK (...) CALL MPI_COMM_SIZE (...)

CALL MPI_COMM_GROUP (MPI_COMM_WORLD, wgroup, ierr)

define something..

CALL MPI_COMM_GROUP_EXCL(wgroup...., newgroup...) CALL MPI_COMM_CREATE(MPI_COMM_WORLD,newgroup,newcomm,ierr)





-MPI_COMM_SPLIT(comm, color, key, newcomm, ierr)

This routine creates as many new groups and communicators as there are distinct values of color.

(processes in the same color are in the same communicator).

The **rank**ings in the new groups are determined by the value of the key.

MPI_UNDEFINED is used as the color for processes to not be included in any of the new groups



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int MPI_Comm_split(MPI_Comm old_comm, int color, int key, MPI_Comm *new_comm)

For a 2D logical grid, create subgrids of rows and columns

c**logical 2D topology with nrow rows and mcol columns irow = Iam/mcol !! logical row number jcol = mod(Iam, mcol) !! logical column number comm2D = MPI_COMM_WORLD call MPI_Comm_split(comm2D, irow, jcol, row_comm, ierr) call MPI_Comm_split(comm2D, jcol, irow, col_comm, ierr)

Figure a. 2D logical Grid		Figure b. 3 Row Subgrids		Figure c. 2 Column Subgrids				
	(0)	(1)	(0) (0)	(1) (1)		(0) (0)	(1) (0)	
	(2)	(3)	(2) (0)	(3) (1)		(2) (1)	(3) (1)	
	(4)	(5)	(4) (0)	(5) (1)		(4) (2)	(5) (2)	

Iam	0	1	2	3	4	5
irow	0	0	1	1	2	2
jcol	0	1	0	1	0	1



http://scv.bu.edu/~kadin/alliance/communicators/MPI_Comm_split.html



Destructors

The communicators and groups from a process' viewpoint are just handles. Like all handles, there is a limited number available: you could (in principle) run out!

-MPI_GROUP_FREE(group, ierr) -MPI_COMM_FREE(comm,ierr)



Intercommunicators

Intercommunicators are associated with 2 groups of disjoint processes.

Intercommunicators are associated with a remote group and a local group

The target process (destination for send, source for receive) is its rank in the remote group.

A communicator is either intra or inter, never both



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Now we can seriously start to work...

