# Scientific and Technical Computing in C 

Stefano Tagliaventi Luca Ferraro<br>CINECA Roma - SCAI Department

## Outline

## Arithmetic

## Aggregate

 Structures Defining Types Arrays Storage \& C. More Arrays(1) Arithmetic Types and Math

Integer Types
Floating Types
Expressions
Arithmetic Conversions

## (2) Aggregate Types

## Data

- Computing $==$ manipulating data and calculating results
- Data are manipulated using internal, binary formats
- Data are kept in memory locations and CPU registers
- C is quite liberal on internal data formats
- Most CPU are similar but all have peculiarities
- C only mandates what is de facto standard
- Some details depend on the specific executing (a.k.a. target) hardware architecture and software implementation
- C Standard Library provides facilities to translate between internal formats and human readable ones
- C allows programmers to:
- think in terms of data types and named containers
- disregard details on actual memory locations and data movements


## C is a Strongly Typed Language

- Each literal constant has a type
- Dictates internal format of the data value
- Each variable has a type
- Dictates content internal format and amount of memory
- Type must be specified in a declaration before use
- Each expression has a type
- And subexpressions have too
- Depends on operators and their arguments
- Each function has a type
- That is the type of the returned value
- Specified in function declaration or definition
- If the compiler doesn't know the type, it assumes int
- Function parameters have types
- I.e. type of arguments to be passed in function calls
- Specified in function declaration or definition
- If the compiler doesn't know the types, it will accept any argument, applying some type conversion rules


## Arithmetic

Floating
Expressions Mixing Types

## Aggregate

Structures Defining Types Arrays

| Type | Sign | Conversion | Width (bits) |  | Size (bytes) |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  | Minimum | Usual | Minimum | Usual |
| signed char | +/- | \%hhd ${ }^{1}$ | 8 | 8 | 1 | 1 |
| unsigned char | + | \%hhu ${ }^{1}$ |  |  |  |  |
| short <br> short int | +/- | \%hd | 16 | 16 | 2 | 2 |
| unsigned short unsigned short int | + | \%hu |  |  |  |  |
| int | +/- | \%d | 16 | 32 | 2 | 4 |
| unsigned unsigned int | + | \% 4 |  |  |  |  |
| $\begin{aligned} & \text { long } \\ & \text { long int } \end{aligned}$ | +/- | \%ld | 32 | 32 or 64 | 4 | 4 or 8 |
| unsigned long unsigned long int | + | \%lu |  |  |  |  |
| long long ${ }^{2}$ <br> long long int ${ }^{2}$ | +/- | \%lld | 64 | 64 | 8 | 8 |
| unsigned long long ${ }^{2}$ unsigned long long int ${ }^{2}$ | + | \%11u |  |  |  |  |

Constraint: short width $\leq$ int width $\leq$ long width $\leq$ long long width

1. C99, in C89 use conversion to/from int types
2. C99

- New platform/compiler? Always check with sizeof (type)
- Values of char and short types just use less memory, they are promoted to int types in calculations

SuperComputing Applications and Innovation

## \#include <limits.h>

## Arithmetic

## Aggregate

Structures
Defining Types
Arrays

| Name | Meaning | Value |
| :---: | :---: | :---: |
| CHAR_BIT | width of any char type | $\geq 8$ |
| SCHAR_MIN | minimum value of signed char | $\leq-128$ |
| SCHAR_MAX | maximum value of signed char | $\geq 127$ |
| UCHAR_MAX | maximum value of unsigned char type | $\geq 255$ |
| SHRT_MIN | minimum value of short | $\leq-32768$ |
| SHRT_MAX | maximum value of short | $\geq 32767$ |
| USHRT_MAX | maximum value of unsigned short | $\geq 65535$ |
| INT_MIN | minimum value of int | $\leq-32768$ |
| INT_MAX | maximum value of int | $\geq 32767$ |
| UINT_MAX | maximum value of unsigned | $\geq 65535$ |
| LONG_MIN | minimum value of long | $\leq-2147483648$ |
| LONG_MAX | maximum value of long | $\geq 2147483647$ |
| ULONG_MAX | maximum value of unsigned long | $\geq 4294967295$ |
| LLONG_MIN | minimum value of long long | $\leq-9223372036854775808$ |
| LLONG_MAX | maximum value of long long | $\geq 9223372036854775807$ |
| ULLONG_MAX | maximum value of unsigned long long | $\geq 18446744073709551615$ |

- Use them to make code more portable across platforms
- New platform/compiler? Always check values


## Integer Literal Constants

## Arithmetic

- Constants have types too
- Compilers must follow precise rules to assign types to integer constants
- But they are complex
- And differ among standards
- Rule of thumb:
- write the number as is, if it is in int range
- otherwise, use suffixes U, L, UL, LL, ULL
- lowercase will do as well, but 1 is easy to misread as 1
- Remember: do not write spokes $=$ bycicles* 2 *36;
- \#define SPOKES_PER_WHEEL 36
- or declare: const int SpokesPerWheel = 36;
- and use them, code will be more readable, and you'll be ready for easy changes


## Integer Types Math

## Arithmetic

- \#include <stdlib.h> to use:

| Function | Returns |
| :--- | :--- |
| abs () | absolute value of an int |
| labs () | absolute value of a long |
| llabs () | absolute value of a long long |

- Use like: $a=a b s(b+i)+c ;$
- For values of type short or char, use abs ()


## Bitwise Arithmetic

- Integer types are encoded in binary format
- Each one is a sequence of bits, each having state 0 or 1
- Bitwise arithmetic manipulates state of each bit
- Each bit of the result of unary operator ~ is in the opposite state of the corresponding bit of the operand
- Each bit of the result of binary operators $\mid, \&$, and $\wedge$ is the OR, AND, and XOR respectively of the corresponding bits in the operands
- Precedence
- a\&b | $c^{\wedge} d \& e$ same as (a\&b) | ( $\left.c^{\wedge}(d \& e)\right)$
- ~a\&b same as (~a) \&b
- Associativity is from left to right
- a | b | csame as (a | b) | c
- As usual, precedence and associativity can be overridden using explicit ( and) , and $\mid=, \varepsilon=$, and $\wedge=$ are available

SuperComputing Applications and Innovation

## More Bitwise Arithmetic

## Arithmetic

- Left and right shifts
- a<n same as $a * 2^{n}$ modulo $2^{\text {type with in bits }}$
- a»n same as a/2n
- Precedence lower than ~ but higher than I, \& and ^
- Beware: if $n>$ type width in bits, or $n<0$, result is undefined
- Applications
- isodd = (a\&1); same as isodd = a\%2;
- b\&255 same as b\%256
- a | 15 same as $(a / 16) * 16+15$
- You have to think in base 2 to get why and if it works
- Think of the examples above ... did you get the pattern?
- 256 is $2^{8}$ and 255 is $2^{8}-1$
- 16 is $2^{4}$ and 15 is $2^{4}-1$
- a | 19 is NOT the same as $(\mathrm{a} / 20) * 20+19$

SuperComputing Applications and Innovation

## Enumerated Types

## Arithmetic

```
enum boundary {
    free_slip,
    no_slip,
    inflow,
    outflow
    };
enum boundary leftside, rightside;
enum liquid {water, mercury} fluid; //may confuse readers
leftside = free_slip;
```

- A set of integer values represented by identifiers
- Under the hood, it's an int
- free_slip is an enumeration constant with value 0
- no_slip is an enumeration constant with value 1
- inflow is an enumeration constant with value 2
- ...


## Choosing Values for Enumeration Constants

## Arithmetic

```
enum spokes {SpokesPerWheel = 36};
enum element {
    hydrogen = 1,
    helium,
    carbon = 6,
    oxygen = 8,
    fluorine
    };
```

- Enumeration constants can be given a specified value
- When the enumeration constant value is not specified:
- if it's the first in the declaration, gets the value 0
- if it's not, gets (value of the previous one+1)
- thus helium above gets 2 , and fluorine gets 9
- negative values can be used too
- A convenient way to give names to related integer constants


## Floating Types (as on Most CPUs)

Arithmetic Integers

## Aggregate

Structures Defining Types Arrays

| Type | Conversion | Width (bits) | Size (bytes) |
| :---: | :---: | :---: | :---: |
|  |  | Usual | Usual |
| float | \%f, \%E, \% $\mathrm{G}^{2}$ | 32 | 4 |
| double | \%lf, \%lE, \% $1 \mathrm{G}^{2}$ | 64 | 8 |
| long double | \%Lf, \%LE, \% $\mathrm{LG}^{2}$ | 80 or 128 | 10 or 16 |
| float _Complex ${ }^{1}$ | none | NA | 8 |
| double _Complex ${ }^{1}$ | none | NA | 16 |
| long double _Complex ${ }^{1}$ | none | NA | 20 or 32 |

Constraints:
all float values must be representable in double
all double values must be representable in long double

1. C99
2. $\% f$ forces decimal notation, $\% \mathrm{E}$ forces exponential decimal notation, \%G chooses the one most suitable to the value

- New platform/compiler? Always check with sizeof (type)
- In practice, always in IEEE Standard binary format, but not a C Standard requirement
- \#include <complex.h> and use float complex, double complex, and long double complex, if your program does not already uses the complexCINECA identifier


## Arithmetic

 Integers
## Aggregate

Structures Defining Types Arrays

| Name | Meaning | Value |
| :---: | :---: | :---: |
| FLT_EPSILON | $\min \{x \mid 1.0+x>1.0\}$ in float type | $\leq 10^{-5}$ |
| DBL_EPSILON | $\min \{x \mid 1.0+x>1.0\}$ in double type | $\leq 10^{-9}$ |
| LDBL_EPSILON | $\min \{x \mid 1.0+x>1.0\}$ in long double type | $\leq 10^{-9}$ |
| FLT_DIG | decimal digits of precision in float type | $\geq 6$ |
| DBL_DIG | decimal digits of precision in double type | $\geq 10$ |
| LDBL_DIG | decimal digits of precision in long double type | $\geq 10$ |
| FLT_MIN | minimum normalized positive number in float range | $\leq 10^{-37}$ |
| DBL_MIN | minimum normalized positive number in long range | $\leq 10^{-37}$ |
| LDBL_MIN | minimum normalized positive number in long double range | $\leq 10^{-37}$ |
| FLT_MAX | maximum finite number in float range | $\geq 10^{37}$ |
| DBL_MAX | maximum finite number in long range | $\geq 10^{37}$ |
| LDBL_MAX | maximum finite number in long double range | $\geq 10^{37}$ |
| FLT_MIN_10_EXP | minimum $x$ such that $10^{X}$ is in float range and normalized | $\leq-37$ |
| DBL_MIN_10_EXP | minimum $x$ such that $10^{X}$ is in double range and normalized | $\leq-37$ |
| LDBL_MIN_10_EXP | minimum $x$ such that $10^{x}$ is in long double range and normalized | $\leq-37$ |
| FLT_MAX_10_EXP | maximum $x$ such that $10^{x}$ is in float range and finite | $\geq 37$ |
| DBL_MAX_10_EXP | maximum $x$ such that $10^{x}$ is in double range and finite | $\geq 37$ |
| LDBL_MAX_10_EXP | maximum $x$ such that $10^{x}$ is in long double range and finite | $\geq 37$ |

- Use them to make code more portable across platforms
- New platform/compiler? Always check values
- "Normalized"? Yes, IEEE Standard allows for even smaller values, with loss of precision, and calls them "denormalized"
- "Finite"? Yes, IEEE Standard allows for infinite values


## Floating Literal Constants

- Need something to distinguish them from integers
- Decimal notation: 1.0, -17., . 125, 0.22
- Exponential decimal notation: $2 \mathrm{E} 19\left(2 \times 10^{19}\right),-123.4 \mathrm{E} 9$ $\left(-1.234 \times 10^{11}\right), .72 \mathrm{E}-6\left(7.2 \times 10^{-7}\right)$
- They have type double by default
- Use suffixes F to make them float or L to make them long double
- Lowercase will do as well, but 1 is easy to misread as 1
- Never write charge = protons*1.602176487E-19;
- \#define UNIT_CHARGE 1.602176487E-19
- or declare: const double UnitCharge $=1.602176487 \mathrm{E}-19$;
- and use them in the code to make it readable
- it will come handier when more precise measurements will be available


## double Math

## Arithmetic

## Integers

Expressions Mixing Types

## Aggregate

Structures Defining Types Arrays

| Function/Macro | Returns |
| :---: | :---: |
| HUGE_VAL ${ }^{1}$ | largest positive finite value |
| INFINITY ${ }^{1}$ | positive infinite value |
| NAN ${ }^{1}$ | IEEE quiet NaN (if supported) |
| double fabs(double x), | $\|\mathbf{x}\|$, |
| double copysign(double $x$, double $y$ ) ${ }^{1}$ | if $\mathbf{y} \neq 0$ returns $\|\mathbf{x}\| \mathbf{y} /\|\mathbf{y}\|$ else returns $\|\mathbf{x}\|$ |
| ```double floor(double x), double ceil(double x), double trunc(double x)}\mp@subsup{}{}{1}\mathrm{ , double round(double x)}\mp@subsup{}{}{1``` | $\begin{gathered} \lfloor\mathbf{x}\rfloor,\lceil\mathbf{x}\rceil \text {, } \\ \text { if } \mathbf{x}>0 \text { returns }\lfloor\mathbf{x}\rfloor \text { else returns }\lceil\mathbf{x}\rceil \text {, } \\ \text { nearest }^{2} \text { integer to } \mathbf{x} \\ \hline \end{gathered}$ |
| double fmod(double $x$, double $y$ ), | $\mathbf{x}$ mod $\mathbf{y}$ (same sign as $\mathbf{x}$ ) |
| double fdim(double $x$, double $y$ ) ${ }^{\text {1 }}$ | if $\mathbf{x}>\mathrm{y}$ returns $\mathbf{x}-\mathrm{y}$ else returns 0 |
| double nextafter (double $x$, double $y$ ) ${ }^{\text {² }}$ | next representable value after $\mathbf{x}$ toward $\mathbf{y}$ |
| $\begin{aligned} & \text { double fmin(double } x, \text { double } y)^{1} \\ & \text { double fmax(double } x, \text { double } y)^{1} \end{aligned}$ | $\begin{aligned} & \min \{\mathbf{x}, \mathbf{y}\} \\ & \max \{\mathbf{x}, \mathbf{y}\} \end{aligned}$ |
| 1. C99 <br> 2. If $\mathbf{x}$ is halfway, returns the farthest from 0 |  |

- \#include <math.h>
- Before C99, there were no fmin () or fmax ()
- Preprocessor macros have been widely used to this aim
- Use the new functions, instead
- More functions are available to manipulate values
- Mostly in the spirit of IEEE Floating Point Standard
- We encourage you to learn more about

SuperComputing Applications and Innovation

## double Higher Math

## Arithmetic

Integers
Expressions Mixing Types

## Aggregate

Structures Defining Types Arrays

| Functions | Return |
| :---: | :---: |
| ```double sqrt(double x), double cbrt(double x)}\mp@subsup{}{}{1}\mathrm{ , double pow(double x, double y), double hypot(double x, double y)}\mp@subsup{}{}{1``` | $\begin{gathered} \sqrt{\mathbf{x}} \\ \sqrt[3]{\mathbf{x}}, \\ \mathbf{x}^{\mathbf{y}} \\ \sqrt{\mathrm{x}^{2}+\mathrm{y}^{2}} \end{gathered}$ |
| double sin(double $x$ ), double cos(double $x$ ), double tan(double $x$ ), double asin(double $x$ ), double acos(double $x$ ), double atan(double $x$ ) | Trigonometric functions |
| double atan2 (double $x$, double $y$ ) | Arc tangent in ( $-\pi, \pi$ ] |
| ```double exp(double x), double log(double x), double log10(double x), double expm1(double x)}\mp@subsup{}{}{1}\mathrm{ , double log1p(double x)}\mp@subsup{}{}{1``` | $\begin{gathered} e^{\mathbf{x}} \\ \log _{e} \mathbf{x}, \log _{10} \mathbf{x} \\ e^{\mathbf{x}}-1, \log (\mathbf{x}+1) \\ \hline \end{gathered}$ |
| ```double sinh(double x), double cosh(double x), double tanh(double x), double asinh(double x)}\mp@subsup{}{}{1}\mathrm{ , double acosh(double x)}\mp@subsup{}{}{1}\mathrm{ , double atanh(double x)}\mp@subsup{}{}{1``` | Hyperbolic functions |
| double erf(double x$)^{1}$ | error function: $\frac{2}{\sqrt{\pi}} \int_{0}^{\mathbf{x}} e^{-t^{2}} d t$ |
| double erfc (double $x$ ) ${ }^{1}$ | $1-\frac{2}{\sqrt{\pi}} \int_{0}^{\mathbf{x}} e^{-t^{2}} d t$ |
| double tgamma (double x$)^{1}$, double ${ }^{\text {l gamma }}$ (double x$)^{1}$ | $\Gamma(\mathbf{x}), \log (\|\Gamma(\mathbf{x})\|)$ |
| 1. C99 |  |

- Again, \#include <math.h>


## double complex Math C99 \& C11

## Arithmetic

## Integers

Expressions Mixing Types

## Aggregate

Structures Defining Types Arrays Storage \& C. More Arrays

| Function/Macro |  |  |  |  | Returns |
| :---: | :---: | :---: | :---: | :---: | :---: |
| double | complex | CMPLX(double $x$, double $y$ ) ${ }^{\text {1 }}$ |  |  | $\mathbf{x}+i \mathbf{y}$ |
| double <br> double <br> double <br> doubl <br> doubl | complex complex complex complex complex | ```cabs(double complex z), carg(double complex z), creal (double complex z), cimag(double complex z), conj(double complex z)``` |  |  | $\|\mathbf{z}\|$ <br> Argument of $\mathbf{z}$ <br> (a.k.a. phase angle), <br> Real part of $\mathbf{z}$, <br> Imaginary part of $\mathbf{z}$, <br> Complex conjugate of $\mathbf{z}$ |
| doubl <br> doubl | complex complex | csqrt (double complex $z$ ), <br> cpow (double complex $z$, double | complex |  | $\begin{aligned} & \sqrt{\mathbf{z}}, \\ & \mathbf{z}^{\mathbf{w}} \end{aligned}$ |
| doubl <br> doubl | complex complex | cexp (double complex $z$ ), clog (double complex z) |  |  | $\begin{gathered} e^{\mathbf{z}} \\ \log _{e} \mathbf{z} \\ \hline \end{gathered}$ |
| 1. C11 |  |  |  |  |  |

- To use them, \#include <complex.h>
- You'll also get: csin(), ccos(), ctan(), casin(), cacos(), catan(), csinh(), ccosh(), ctanh(), casinh(), cacosh(), catanh()
- And I for the imaginary unit

SuperComputing Applications and Innovation

## float and long double Math

- Before C99, all functions were only for doubles
- And automatic conversion of other types was applied
- But from 1999 C is really serious about floating point math
- All functions exist also for float and long double
- Same names, suffixed by $f$ or 1
- Like acosf() for arccosine of a float
- Or cacosl() for long double complex
- Ditto for macros, like HUGE_VALF or CMPLXL ()
- If you find this annoying (it is!):
- \#include <tgmath.h>
- and use everywhere, for all real and complex types, function names for double type
- These are clever type generic processor macros, expanding to the function appropriate to the argument


## Expressions

- A fundamental concept in C
- A very rich set of operators
- Almost everything is an expression
- Even assignment to a variable
- C expressions are complicated
- Expressions can have side effects
- Not all subexpressions are necessarily computed
- Except for associativity and precedence rules, order of evaluation of subexpressions is up to the compiler
- Values of different type can be combined, and a result produced according to a rich set of rules
- Sometimes with surprising consequences
- We'll give a simplified introduction
- Subtle rules are easily forgotten
- Relying on them makes the code difficult to read
- When you'll find a puzzling piece of code, you can always look for a good manual or book


## Arithmetic Expressions

- Binary operators +, -, * (multiplication) and / have the usual meaning and behavior
- Unary operator - evaluates to the opposite of its operand
- Unary operator + evaluates to its operand
- Precedence
- $-\mathrm{a} * \mathrm{~b}+\mathrm{c} / \mathrm{d}$ same as ( (-a) *b) $+(\mathrm{c} / \mathrm{d})$
- -a + b same as (-a) + b
- Associativity of binary ones is from left to right
- $a+b+c s a m e ~ a s ~(a+b)+c$
- $\mathrm{a} * \mathrm{~b} / \mathrm{c} * \mathrm{~d}$ same as ( $(\mathrm{a} * \mathrm{~b}) / \mathrm{c})$ *d
- Explicit ( and) override precedence and associativity
- Only for integer types, \% is the modulo operator (27\%4 evaluates to 3), same precedence as /


## Hitting Limits

- All types are limited in range
- What about:
- INT_MAX + 1? (too big)
- INT_MIN*3? (too negative)
- Technically speaking, this is an arithmetic overflow
- And division by zero is a problem too
- For signed integer types, the Standard says:
- behavior and results are unpredictable
- i.e. up to the implementation
- For other types, the Standard says:
- arithmetic on unsigned integers must be exact modulo $2^{\text {type }}$ width, no overflow
- with floating types, is up to the implementation (you can get DBL_MAX, or a NaN, or an infinity)
- Best practice: NEVER rely on behaviors observed with a specific architecture and/or compiler


## Assignment Operator

- Binary operator =
- assigns the value of the right operand to the left operand
- and returns the value of the right operand
- thus $\mathrm{a}=\mathrm{b} * 2$ is an expression with value $\mathrm{b} * 2$ and the side effect of changing variable a
- $a=b * 2$; is an assignment statement
- The left operand must be something that can store a value
- In C jargon, an Ivalue
- a = 20 is OK, if $a$ is a variable
- $20=a$ is not
- Precedence is lowest (except for, operator) and associativity is from right to left
- $a=b * 2+c$ same as $a=(b * 2+c)$
- $z=a=b * 2+c$ same as $z=(a=(b * 2+c))$
- You'll read the latter form, particularly in while () statements, but avoid writing it



## More Assignment Operators

- Most binary operators offer useful shortcut forms:

| Expression | Same as |
| :--- | :--- |
| $\mathrm{a}+=\mathrm{b}$ | $\mathrm{a}=\mathrm{a}+\mathrm{b}$ |
| $\mathrm{a}-=\mathrm{b}$ | $\mathrm{a}=\mathrm{a}-\mathrm{b}$ |
| $\mathrm{a} *=\mathrm{b}$ | $\mathrm{a}=\mathrm{a} * \mathrm{~b}$ |
| $\mathrm{a} /=\mathrm{b}$ | $\mathrm{a}=\mathrm{a} / \mathrm{b}$ |
| $\mathrm{a} \%=\mathrm{b}$ | $\mathrm{a}=\mathrm{a} \% \mathrm{~b}$ |

- In heroic times, used to map some CPUs optimized instructions
- With nowadays optimizing compilers, only good to spare keystrokes
- You'll find them often, particularly in for (; ; ) statements


## More Side Effects

- Pre-increment/decrement unary operators: ++ and --
- ++i same as (i = i + 1)
- --i same as (i = i - 1)
- Post-increment/decrement unary operators: ++ and --
- i++ increments i content, but returns the original value
- i-- decrements i content, but returns the original value
- Operand must be an Ivalue
- Precedence is highest
- Quite handy in while () and for (; ; ) statements
- Easily becomes a nightmare inside expressions
- Particularly when you change the code

SuperComputing Applications and Innovation

Order of Subexpressions Evaluation

- $i$ is an int type variable whose value is 5
j = 4*i++ - 3*++i;
foo(++i, ++i);
- Which value is assigned to $j$ ?
- Could be
- Or could as well be
- Which values are passed to foo()?
- Could be foo( , )
- Or could as well be foo( , )
- Order of evaluation of subexpressions is implementation defined!
- Ditto for order of evaluation of function arguments!
- NEVER! NEVER pre/post-in/de-crement the same variable twice in a single expression, or function call!

SuperComputing Applications and Innovation

## Logical Expressions

## Arithmetic

## Integers

## Floating

- Comparison operators
- == (equal), != (not equal), >, <, >=, <=
- Compare operand values
- Return int type 0 if evaluation is false, 1 if true
- Precedence lower than arithmetic operators, higher than bitwise and logical operators
- In doubt, add parentheses, but be sober
- Logical operators
- ! is unary NOT, $\& \&$ is binary AND, I। is binary OR
- Zero operand are considered false, non zero ones true
- Return int type 0 if comparison is false, 1 if true
- Precedence of ! just lower than ++ and --
- \&\&, II: higher than = and friends
- !a\&\&b || a\&\&!b means ((!a)\&\&b) || (a\&\&(!b))
- Again: in doubt, add parentheses, but be sober


## More Logic from math . h

## Arithmetic

## Integers

Floating

- Some macros to tame floating point complexity
- isfinite()
- True if argument value is finite
- isinf()
- True if argument value is an infinity
- isnan ()
- True if argument value is a NaN
- And more, if you are really serious about floating point calculations
- Mostly in the spirit of IEEE Floating Point Standard
- Learn more about it, before using them


## Being Completely Logical

## Arithmetic

- C99 defines integer type _Bool
- Only guaranteed to store 0 or 1
- Perfect for logical (a.k.a. boolean) expressions
- Use it for "flag" variables, and to avoid surprises
- Better yet, \#include <stdbool.h>, and use type bool, and values true and false
- Watch your step!
- Simply mistype \& for \& \& or vice versa
- Simply mistype |। for I
- You'll discover, possibly after hours of debugging, that (bitwise arithmetic) != (logical arithmetic)
- C99 offers a fix to this unfortunate choice
- \#include <iso646.h>
- And use not, or, and and in place of !, |। and \&\&

SuperComputing Applications and Innovation

## Even More Side Effects

## Arithmetic

- Right operand of II and $\& \&$ is evaluated after left one
- And is not evaluated at all if:
- left one is found true for an II
- left one is found false for an $\& \&$
- Beware of "short circuit" evaluation...
- ... if the right operand is an expression with side effects!
- A life saver in preprocessor macros and a few more cases
- But makes your code less readable
- Use nested if () whenever you can
- logical-expr ? expr1 : expr2
- expr1 is only evaluated if logical-expr is true
- expr2 is only evaluated if logical-expr is false
- Again, is a life saver in preprocessor macros
- But in normal use an if () is more readable


## Mixing Types in Expressions



- A result will always be produced
- Whether this is the result you expect, it's another story
- Broadly speaking, the base concept is clear
- For each binary operator in the expression, in order of precedence and associativity:
- if both operands have the same type, fine
- otherwise, operand with narrower range is converted to type of other operand
- OK when mixing floating types
- The wider range includes the narrower one
- OK when mixing signed integer types
- The wider range includes the narrower one
- OK even when mixing unsigned integer types
- The wider range includes the narrower one


## Type Conversion Traps

- For the assignment operator:
- if both operands have the same type, fine
- otherwise, right operand is converted to left operand type
- if the value cannot be represented in the destination type, it's an overflow, and you are on your own
- We said: in order of precedence and associativity
- if a is a type long long int variable, and b is a 32 bits wide int type variable and contains value INT_MAX, in:
$\mathrm{a}=\mathrm{b} * 2$
multiplication will overflow
- and in:
$\mathrm{a}=\mathrm{b} * 2+1 \mathrm{LL}$
multiplication will overflow too
- while:
$a=b * 2 L L+1$
is OK

SuperComputing Applications and Innovation

- Think of mixing floating and integer types
- Floating types have wider range
- But not necessarily more precision
- A 32 bits float has fewer digits of precision than a 32 bits int
- And a 64 bits double has fewer digits of precision than a 64 bits int
- The result could be smaller than expected
- Think of mixing signed and unsigned integer types!
- Negative values cannot be represented in unsigned types
- Half of the values representable in an unsigned type, cannot be represented in a signed type of the same width
- So, you are in for implementation defined surprises!
- And Standard rules are quite complicated
- We spare you the gory details, simply don't do it!

SuperComputing Applications and Innovation

## Cast Your Subexpressions

- (type)
- Unsurprisingly, it's an operator
- Precedence just higher than multiplication, right-to-left associative
- Use it like (unsigned long) (sig + ned)
- Casting let you override standard conversion rules
- In previous example, you could use it like this:
$a=$ (long long int) $b * 2+1$
- Type casting is not magic
- Just instructs compiler to apply the conversion you need
- Only converts values, not type of variables you assign to
- Do not abuse it
- Makes codes unreadable
- Could be evidence of design mistakes
- Or that your C needs a refresh


# Scientific and Technical Computing in C 

Stefano Tagliaventi Luca Ferraro<br>CINECA Roma - SCAI Department

## Outline

## Arithmetic

## Integers

# (1) Arithmetic Types and Math 

(2) Aggregate Types

Structure Types
Defining New Types
Arrays
Storage Classes, Scopes, and Initializers Arrays \& Functions

## struct

## Arithmetic

Integers Floating Expressions Mixing Types

## Aggregate

Defining Types Arrays Storage \& C. More Arrays

```
struct vect3D {
        double x, y, z;
};
struct vect3D va, vb;
// REMINDER: I have to make vcross() more efficient!
struct vect3d vcross(struct vect3D u, struct vect3D v) {
        struct vect3D c;
        c.x = u.y*v.z - u.z*v.y;
        c.y = u.z*v.x - u.x*v.z;
        c.z = u.x*v.y - u.y*v.x;
    return c;
}
//...
    vc = vcross(va, vb);
```

- Aggregates a single type from named, typed components (a.k.a. members)
- The vect 3D tag must be unique among structure tags
- struct components can be independently accessed using the . binary operator

SuperComputing Applications and Innovation

## structs Are Flexible

## Arithmetic

```
struct ion {
    struct vect3D r; // position
    struct vect3D v; // velocity
    enum element an; // atomic number
    int q; // in units of elementary charges
};
struct ion a;
//...
    a.r.x += dt*a.v.x; // very low order in time...
```

- struct components can be inhomogeneous
- And they can also be structs, of course
- To access nested struct components, chain . expressions
- Best practice: order components by decreasing size
- You'll get better performances
- To know, you can use sizeof() operator on any type


## structs: a Concrete Example

## Arithmetic

## Integers

- structs are widely used in C Standard Library
- Like in struct tm, below, defined in time.h
- Used to convert from/to internal time representation time_t

```
struct tm {
    int tm_sec; // seconds after the minute [0, 60]
    int tm_min; // minutes after the hour [0, 59]
    int tm_hour; // hours since midnight [0, 23]
    int tm_mday; // day of the month [1, 31]
    int tm_mon; // months since January [0, 11]
    int tm_year; // years since 1900
    int tm_wday; // days since Sunday [0, 6]
    int tm_yday; // days since January 1 [0, 365]
    int tm_isdst; // Daylight Saving Time flag
};
```

SuperComputing Applications and Innovation

## typedef

## Arithmetic

## Integers

Floating Expressions Mixing Types

## Aggregate

Structures Defining Types

## Arrays

```
typedef struct vect3D position, velocity;
typedef enum element element; // let's spare keystrokes
typedef int charge; // I'll maybe switch to short or signed char
typedef struct ion {
    position r;
    velocity v;
    element an;
    charge q;
} ion;
ion a;
```

- typedef turns a normal declaration into a declaration of a new type (as usual, a legal identifier)
- The new type can be used as the native ones
- Great to save keystrokes
- Even better to write self-documenting code
- Shines in hiding and factoring out implementation details
- struct tags and type identifiers belong to separate sets


## typedef in C Standard Library

- typedef is widely used in C Standard Library
- Mostly to abstract details that may differ among implementations
- E.g. size_t from stddef.h
- Type of value returned by sizeof()
- Different platforms allow for different memory sizes
- size_t must be "typedefed" to an integer type able to represent the maximum possible variable size allowed by the implementation
- E.g. clock_t from time.h
- Type of value returned by clock ()
- Cast it to double, divide by CLOCK_PER_SEC, ...
- and you'll know the CPU time in seconds used by your program from its beginning

SuperComputing Applications and Innovation

## Arrays

- some_type a[n];
- declares a collection of $n$ variables of type some_type
- the variables (a.k.a. elements) are laid out contiguously in memory
- each element can be read or written using the syntax a [ integer indexing expression]
- first element is a [0], second one is a [1], last one is a [ $n-1$ ]
- You can't work on an array as a whole
- Use array elements (if allowed...) in expressions and assignments
- There is no bound checking!
- Use a negative index, or an index too big, and you are accessing something else, if any
- Compiler options to (very slowly) check every access
- A common mistake:
- to access from double a[1] to double a [ $n$ ]
- Fortran programmers beware!

SuperComputing Applications and Innovation

## Arrays of(Arrays of(Arrays of(...)))

## Arithmetic

- C has no concept of multidimensional arrays
- But array is a regular C type (you can even sizeof(double[150]))
- Thus, arrays of arrays can be declared
- A simple, practical abstraction
- Very annoying to Fortran or Matlab programmers
- int a[12][31];
- declares an array of 12 elements
- and each element is itself an array of 31 ints
- double b[130][260][260];
- declares an array of 130 elements
- and b[37] is itself an array of 260 elements
- and b[37] [201] is again an array of 260 doubles
- By the way, you can also use sizeof (b), it works


## Array Memory Layout

## Arithmetic

## Aggregate

Structures
Defining Types
int a[10];

| $a[0]$ | $a[1]$ | $a[2]$ | $a[3]$ | $a[4]$ | $a[5]$ | $a[6]$ | $a[7]$ | $a[8]$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $a[9]$ |  |  |  |  |  |  |  |  |

int b[5][2];

| $\mathrm{b}[0]$ | $\mathrm{b}[1]$ | $\mathrm{b}[2]$ | $\mathrm{b}[3]$ | $\mathrm{b}[4]$ |
| :---: | :---: | :---: | :---: | :---: |


| b [0] [0] | b [0] [1] | b [1] [0] | b [1] [1] | b [2] [0] | b [2] [1] | b [3] [0] | b [3] [1] | b [4] [0] | b [4] [1] |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |

## A Very Important Digression

- Storage duration
- To make it simple, the life time of a variable
- Also influences the part of memory where it's allocated
- Scope
- The region where a variable or function is accessible, a.k.a. "visible"
- Qualifiers
- The value in a const variable cannot be changed
- There are more, but we'll not discuss them
- Initializers
- Values assigned to a variable at declaration


SuperComputing Applications and Innovation

## Storage Duration

- A variable can be
- Automatic: it can be created when needed, and destroyed when not needed anymore
- Static: it persists for the whole duration of the program
- Variables declared outside of any functions (i.e. at file scope) are static
- By default, are automatic:
- all variables declared inside a compound statement
- function parameters
- The default can be overridden using static
- Functions are static too, because to call them you need their code to persist in memory

SuperComputing Applications and Innovation

- By default, variables declared at file scope and functions are extern
- i.e. visible to the linker, and to the whole program
- Unless you declare them to be static only
- Variables declared at file scope and functions are visible to all blocks in the same source file
- Variables declared in a block are only visible in the block and in all scopes it encloses
- Unless you declare them extern
- But in most cases that's a symptom of bad design
- A variable declared in a block hides anything declared with the same name in enclosing scopes


## Variable Initializers

- The content of an automatic variable is uninitialized until the variable is assigned a value
- Uninitialized is a polite form for "unpredictable rubbish"
- double $f=2.5$; is a practical shorthand for:

```
double f;
f = 2.5;
```

- Expressions can be used as initializers, as long as they can be computed at that point:

```
double pi = acos(-1.0);
double pihalf = pi/2.0;
```

is legal, while the following:

```
double pihalf = pi/2.0;
double pi = acos(-1.0);
```

obviously is not

## More on Variable Initializers

- structs can be initialized too, as in:

```
struct vect3D V = {0.0, 1.0, 0.0};
```

- Same for arrays, as in:

```
float rot[2][2] = {{0.0, -1.0}, {1.0, 0.0}};
```

- $\{0.0,1.0,0.0\}$ and $\{\{0.0,-1.0\},\{1.0,0.0\}\}$ are said compound literals
- By default, static variables are initialized to 0
- But they can be initialized to different values
- Expressions can also be used, with some restrictions
- For a static variable, initialization expression must be computed at compile time
- I.e. it must be a constant expression, containing only constants
- No variables, no function calls are permitted


## Arrays and Storage Classes

- Static arrays must be dimensioned with constant expressions
- Before C99, this was true for automatic arrays too
- So to use an array in a function, you had to dimension it for the largest possible amount of work
- A waste of memory and error prone
- C99 has a much better way
- Variable length arrays
- Arrays whose size is unknown until run time
- Automatic arrays can have their dimension specified by a nonconstant expression
- Every time execution enters the block, the expression is evaluated
- And the array size is determined, up to exit from the block


## Arrays as Function Arguments

- Arrays can be huge
- And usually are, in S\&T computing
- Passing them by value would be too costly
- Moreover, arrays cannot be used in assignments
- Thus a function cannot return an array
- The solution
- The address of the array is passed to a function
- And elements can be accessed by it
- (Later on, you'll understand how)
- This allows elements to be assigned to
- Thus a function has a way to "return" an array result
- A mixed blessing: allows changes to happen by mistake
- Best practice: declare an array parameter const if your only intent is reading its elements


## Averaging, the C99 Way

## Arithmetic

## Integers

## Floating

## Aggregate

- Let's write a function to average an array of doubles
- And make it generic in the array length
- Variable length array parameters come to the rescue

```
double avg(int n, const double a[n]) {
    int i;
    double sum = 0.0;
    for (i=0; i<n; ++i)
        sum += a[i];
    return sum/n;
}
```

Beware: double avg (double $a[n]$, int $n$ ) does not work!

## Averaging, the Old Way

- Before C99, there were no VLAs
- The solution was simple
- Compiler just uses type size to find the right element
- No bounds checking, no bound needed
- Many still write that way: it's equivalent, but less readable

```
double avg(int n, const double a[]) {
    int i;
    double sum = 0.0;
    for (i=0; i<n; ++i)
        sum += a[i];
    return sum/n;
}
```


## Calling avg ()

Arithmetic

- New or old style, simply pass array dimension and name
- If avg () is written using VLAs, pedantic compilers may give a warning on function call, even if it's correct: they are wrong, check with Standard document or good book

```
double mydata[N];
double mydata_avg;
// read or compute N doubles into mydata[]
mydata_avg = avg(N, mydata);
```


## Averaging Arrays of 5 Elements

Arithmetic

- Let's write a function to average arrays of 5 doubles
- And make it generic, as usual
- Again, VLA parameters come to the rescue

```
void avg5(int n, const double a[n][5], double b[5]) {
    int i, j;
    for (j=0; j<5; ++j)
        b[j] = 0;
    for (i=0; i<n; ++i)
        for (j=0; j<5; ++j)
            b[j] += a[i][j];
    for (j=0; j<5; ++j)
        b[j] /= n;
}
```

Notice: this order of loops nesting gives faster execution

## Averaging Arrays of 5 Elements, the Old Way

- Let's write a function to average arrays of 5 doubles
- And make it generic, as usual
- Again, do not specify first bound
- Again, it's equivalent

```
void avg5(int n, const double a[][5], double b[5]) {
    int i, j;
    for (j=0; j<5; ++j)
        b[j] = 0;
    for (i=0; i<n; ++i)
        for (j=0; j<5; ++j)
            b[j] += a[i][j];
    for (j=0; j<5; ++j)
        b[j] /= n;
}
```

Notice: this order of loops nesting gives faster execution

## Calling avg5 ()

Arithmetic

- New or old style, simply pass array dimension and name
- If avg5 () is written using VLAs, pedantic compilers may give a warning on function call, even if it's correct: they are wrong, check with Standard document or good book

```
double mydata[N][5];
double mydata_avg[5];
// read or compute N 5-uples of doubles into mydata[]
avg5(N, mydata, mydata_avg);
```


## Averaging Arrays of Arbitrary Length

- Let's generalize the average to set of $m$ numbers
- And make it generic, as usual
- Again, VLA parameters come to the rescue

```
    void avg(int n, int m, const double a[n][m], double b[m]) {
int i, j;
for (j=0; j<m; ++j)
        b[j] = 0;
for (i=0; i<n; ++i)
    for (j=0; j<m; ++j)
            b[j] += a[i][j];
for (j=0; j<m; ++j)
    b[j] /= n;
```

\}

Notice: this order of loops nesting gives faster execution

## Usage of type void

Arithmetic

- What type is void?
- As a return type, it tells a function returns nothing
- As a parameter, it tells no arguments are accepted
- double avg (void) \{
- Why there is no return statement in avg ()?
- It returns nothing and completes at the closing brace


## Calling Generic avg ()

## Arithmetic

Integers
Floating
Expressions Mixing Types

## Aggregate

Structures Defining Types Arrays Storage \& C.

- Again, simply pass array dimension and name
- Using casts for arrays of doubles
- If avg () is written using VLAs, pedantic compilers may give a warning on function call, even if it's correct: they are wrong, check with Standard document or good book

```
double mydata1[N][12];
double mydata1_avg[12];
double mydata2[N][7];
double mydata2_avg[7];
double mydata3[N][1];
double mydata3_avg[1];
double mydata4[N];
double mydata4_avg[1];
// read or compute N 12-uples of doubles into mydata1[]
// read or compute N 7-uples of doubles into mydata2[]
// read or compute N 1-uples of doubles into mydata3[]
// read or compute N doubles into mydata4[]
avg(N, 12, mydata1, mydatal_avg);
avg(N, 7, mydata2, mydata2_avg);
avg(N, 1, mydata3, mydata3_avg);
avg(N, 1, (double [N][1])mydata4, mydata4_avg);
```


## Matrix Algebra, the C99 Way

- Let's write a function to compute the trace of a matrix of doubles
- And make it generic in the matrix size
- Again, variable length array parameters come to the rescue
- Again, you may get warnings on calls, and they could prove wrong

```
double tr(int n, const double a[n][n]) {
    int i;
    double sum = 0.0;
    for (i=0; i<n; ++i)
        sum += a[i][i];
    return sum;
}
```

Beware: compiler will not check the array dimensions match!

## Matrix Algebra, the Old Way

## Arithmetic

- Before C99, there were no VLAs
- The solution was not that simple...
- Only the 'first dimension' of an array parameter could be left unspecified at compile time
- To understand the solution, you have to learn more


## Rights \& Credits

These slides are (C)CINECA 2014 and are released under the Attribution-NonCommercial-NoDerivs (CC BY-NC-ND) Creative Commons license, version 3.0.

Uses not allowed by the above license need explicit, written permission from the copyright owner. For more information see:
http://creativecommons.org/licenses/by-nc-nd/3.0/

Slides and examples were authored by:

- Michela Botti
- Federico Massaioli
- Luca Ferraro
- Stefano Tagliaventi

