



Programmazione Avanzata

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Bugs and Prevention

Testing

Static analysis

Run-time analysis







As soon as we started programming, we found to our surprise that it wasn't as easy to get programs right as we had thought. Debugging had to be discovered. I can remember the exact instant when I realized that a large part of my life from then on was going to be spent in finding mistakes in my own programs.!

Maurice Wilkes discovers debugging, 1949.





Testing-Debugging



- ► TESTING: finds errors.
- DEBUGGING: localizes and repairs them.

TESTING DEBUGGING CYCLE: we test, then debug, then repeat.





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TESTING DEBUGGING CYCLE: we test, then debug, then repeat.

Program testing can be used to show the presence of bugs, but never to show their absence! Edsger Dijkstra







Defect: An incorrect program code







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Infection chain





The programmer creates a defect in the program code (also known as bug or fault).





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 $Defect \Longrightarrow Infection$

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- ► The defect causes an in infection in the program state.





Infection chain



 $\mathsf{Defect} \Longrightarrow \mathsf{Infection} \Longrightarrow \mathsf{Failure}$

- The programmer creates a defect in the program code (also known as bug or fault).
- ► The defect causes an in infection in the program state.
- ► The infection creates a failure an externally observable error.





Tracking down defect



Failure

► A Failure is visible to the end user of a program. For example, the program prints an incorrect output.





Tracking down defect



Infection \leftarrow Failure

- ► A Failure is visible to the end user of a program. For example, the program prints an incorrect output.
- Infection is the underlying state of the program at runtime that leads to a Failure. For example, the program might display the incorrect output because the wrong value is stored in avariable.





Tracking down defect



Defect \leftarrow Infection \leftarrow Failure

- A Failure is visible to the end user of a program. For example, the program prints an incorrect output.
- Infection is the underlying state of the program at runtime that leads to a Failure. For example, the program might display the incorrect output because the wrong value is stored in avariable.
- A Defect is the actual incorrect fragment of code that the programmer wrote; this is what must be changed to fix the problem.





First of all..



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- Preventing bugs.
- Detecting/locating bugs.





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Ca. 80 percent of software development costs spent on identifying and correcting defects.

It is much more expensive (in terms of time and effort) to detect/locate exisisting bugs, than prevent them in the first place.





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- We have an agreed common goal: to raise the standard of software (and software development) around the world by promoting Code Inspection as a prerequisite to Software Testing in the production and delivery cycle.
- Our strong hope is that eventually the idea of inspecting code before testing will be as self-evident as testing software before using it.





The Fundamental question



How can I prevent Bugs?





The Fundamental question



How can I prevent Bugs?

- Design.
- Good writing.
- Self-checking code.
- Test scaffolding.







Programming is a design activity. It's a creative act, not mechanical code generation.



Preventing bugs via design



Programming is a design activity. It's a creative act, not mechanical code generation.

- Good modularization.
- Strong encapsulation/information hiding.
- Clear,simple pre- and post-processing.
- Requirements of operations should testable.







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- The smaller the component, the easier it is to understand.
- Every component has interfaces with other components, and all interfaces are sources of confusion.







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- Large components reduce external interfaces but have complicated internal logic that may be difficult or impossible to understand.
- The art of software engineering is setting component size and boundaries at points that balance internal complexity against interface complexity to achieve an overall complexity minimization. A study showed that when routines averaged 100 to 150 lines each, code was more stable and required less changes.





Good modularization...



- reduces complexity
- avoids duplicate code
- facilitates reusable code
- limits effects of the changes
- facilitates test
- results in easier implementation







- The principle of information hiding suggests that modules be characterized by design decisions that each hides from all others.
- The information contained within a module is inaccesible to other modules that have no need for such information.
- Hiding implies that effective modularity can be achieved by defining a set of independent modules that communicate with one another only that information necessary to achieve software function.
- Because most data and procedure are hidden from other parts of software, inadvertent errors introduced during modification are less likely to propagate to other locations within the software.





Cohesion & Coupling



- Cohesion: a measure of how closely the instructions in a module are related to each other.
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- Avoid global variables.
- As a general rule, you should always aim to create modules that have strong cohesion and weak coupling.
- The routines with the highest coupling-to-cohesion ratios had 7 times more errors than those with the lowest ratios.





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Spaghetti code is characterized by very strong coupling.





Preventing bugs via good writing



- Clarity is more important than efficiency: clarity of writing and style.
 - Use simple expressions, not complex
 - Use meaningful names
 - Clarity of purpose for:
 - Functions.
 - Loops.
 - Nested constructs.






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- Comments,comments,comments.
- Small size of functions, routines.







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► Comments,comments,comments.

 Small size of functions, routines.
 A study at IBM found that the most error-prone routines were those larger than 500 lines of code.







- Use the smallest acceptable scope for:
 - Variable names.
 - Static/local functions.
- Use named intermediate values.
- Avoid "Magic numbers".
- Generalize your code.
- Document your program.
- Write standard language.





Clarity of writing



```
do j=1,n
    do i=1,n
        v(i,j)=(i/j)*(j/i)
        end do
end do
.....
```





Clarity of writing



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do j=1,n
    do i=1,n
        v(i,j)=(i/j)*(j/i)
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end do
.....
```

```
.....
do j=1, n
    do i=1, n
    v(i,j)=0.0
    v(i,i)=1.0
    end do
end do
.....
```







The name of a variable, function, etc. should answer all the big questions:







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- what it does







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- why it exists
- what it does
- how it is used.







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If a name requires a comment, then the name does not reveal its intent.







Purpose of variables	Good names	Bad names
	Good descriptors	Bad descriptors
Running total of	runningTotal,checkTotal	written,ct,checks,
checks written to date	nChecks	CHKTTL,x,x1,x2
Velocity of a bullet train	velocity,trainVelocity velocityInMph	velt,v,tv, x,x1,x2,train
Current date	currentDate,todaysDate	cd,current,c,x,x1,x2,dat
Lines per page	LinesPerPage	lpp,lines,l,x,x1,x2





Optimum name length



- Too long: numberOfPeopleOnTheUsOlympicteam numberOfSeatsInTheStadium maximumNumberOfPointsModernOlympics
- Too short: n,np,ntn n,ns,nsisd m,mp,max,points
- Just right: numberTeammembers,teamMemberCount numSeatInStadium,seatCount teamPointsMax,pointsRecord





Optimum name length



- Too long: numberOfPeopleOnTheUsOlympicteam numberOfSeatsInTheStadium maximumNumberOfPointsModernOlympics
- Too short: n,np,ntn n,ns,nsisd m,mp,max,points
- Just right: numberTeammembers,teamMemberCount numSeatInStadium,seatCount teamPointsMax,pointsRecord The effort required to debug a program is minimized when variables had names averaging 10 to 16 characters long.







Make sure comments and code are agree.







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- Don't duplicate code in a comment.
- Format a program to help the reader understand it.







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- Format a program to help the reader understand it.
- Indent to show the logical structure of a program.





Conversion of mass mixing ratio in units of g/kg to volume mixing ratio in units of ppmv

REAL(8) :: ppmv				
REAL(8) :: Mixing_Ratio				
REAL(8) :: Molecular_Weight				
· · ·				
<pre>ppmv = 1.0e+03 * Mixing_Ratio * 28.9648 / Molecular_Weight</pre>				





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ppmv = 1.0e+03 * Mixing_Ratio * 28.9648 / Molecular_Weight
```

```
REAL(8), PARAMETER :: G_TO_KG = 1.0d-03
REAL(8), PARAMETER :: PPV_TO_PPMV = 1.0d+06
REAL(8), PARAMETER :: SCALE_FACTOR = G_TO_KG * PPV_TO_PPMV
REAL(8), PARAMETER :: MW_DRYAIR = 28.9648
...
REAL(8) :: ppmv
REAL(8) :: Mixing_Ratio
REAL(8) :: Molecular_Weight
...
ppmv = SCALE_FACTOR * Mixing_Ratio * MW_DRYAIR / Molecular_Weight
```





Avoid magic number



- Changes can be made more reliably.
- Changes can be made more easily.
- ► Your code is more readable.



Preventing bugs via Self-checking code



- Checking assumptions.
- "assert" macro.

nputing Applications and Innovation

- Custom "assert" macros.
- Assertions about intermediate values.
- Preconditions, postcondition.



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Defensive programming.





Checking assumptions



- That an input parameter's value falls within its expected range (or an output parameters' value does)
- That the value of an input-only variable is not changed by a routine.
- That a pointer is non-NULL.
- That an array or other container passed into a routine can contain at least X data number of data elements.
- That a table has been initialized to contain real values.
- ► Etc.





Assertions



- An assertion is a code (usually routine or macro) that allows a program to check itself as runs.
- When an assertion is true, that means everything is operating as expected
- When it's false, that means it has detected an unexpected error in the code.
- ► Use error handling code for conditions you expected to occur.





Error handling technique



- Return a neutral value.
- Substitute the closest legal value.
- Log a warning message to a file.
- Return an error code.
 - Set a value of a status variable.
 - Return status as the function's return value.
 - Throw an exception using the language's build-in exception mechanism.
- Display an error message wherever the error is encountered.
- Handle the error in whatever way works best locally.
- Shutdown.





From "The Elements of Programming Style" Kernighan and Plauger

- Write clearly don't be too clever.
- Say what you mean, simply and directly.
- Use library functions whenever feasible.
- Avoid too many temporary variables.
- Write clearly don't sacrifice clarity for "efficiency".
- Let the machine do the dirty work.
- Replace repetitive expressions by calls to common functions.
- Parenthesize to avoid ambiguity.
- Choose variables names that won't be confused.
- Avoid unnecessary branches.





- ► If a logical expression is hard to understand, try transforming it.
- Choose a data representation that makes the program simple.
- Write first in easy-to-understand pseudo language; then translate into whatever language you have to use.
- Modularize. Use procedures and functions.
- Avoid gotos completely if you keep the program readable.
- Don't patch bad code rewrite it.
- Write and test a big program in small pieces.
- Use recursive procedures for recursively-defined data structures.
- Test input for plausibility and validity.
- Make sure input doesn't violate the limits of the program.





- Terminate input by end-of-file marker, not by count.
- Identify bad input; recover if possible.
- Make input easy to prepare and output self-explanatory.
- Use uniform input formats.
- Make input easy to proofread.
- Use self-identifying input. Allow defaults. Echo both on output.
- Make sure all variables are initialized before use.
- Don't stop at one bug.
- Use debugging compilers.
- Watch out for off-by-one errors.





- Take care to branch the right way on equality.
- Avoid multiple exits from the loops.
- Make sure your code does "nothing" gracefully.
- Test programs at their boundary values.
- Check some answers by hand.
- 10.0 times 0.1 is hardly ever 1.0.
- 7/8 is zero while 7.0/8.0 is not zero.
- Don't compare floating point numbers solely equality.
- Make it right before you make it faster.
- Make it fail-safe before you make faster.





- Make it clear before you make it faster.
- Don't sacrifice clarity for small gains of "efficiency".
- Let your compiler do the simple optimizations.
- Don't strain to re-use code; reorganize instead.
- Make sure special cases are truly special.
- Keep it simple to make it faster.
- Don't diddle code to make it faster find a better algorithm.
- Instrument yout programs. Measure before making "efficiency" changes.
- Make sure comments and code agree.
- Don't just echo the code with comments make every comment count.





- Don't comment bad code rewrite it.
- Use variable names taht mean something.
- Use statement labels that mean something.
- Format a program to help the reader understand it.
- Document your data layouts.
- Don't over-comment.








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Run-time analysis







Input

Read three integer values from the command line. The three values represent the lengths of the sides of a triangle.







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Output

Tell whether the triangle is: Scalene Isosceles Equilater.







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Tell whether the triangle is: Scalene Isosceles Equilater.

Create a set of test cases for this program. ("The art of sotware testing" G.J. Myers)







Q:1 Do you have a test case with three integers greater than zero such that the sum of two of the numbers is less than the third ?







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- (4,1,2) is an invalide triangle.
- (a,b,c) with a > b+c



Define valide triangles a < b + c







Q:2 Do you have a test case with some permutations of previous test?







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(1,4,2) (4,1,2)

Fulfill above definition, but are still invalid.

Patch definition of valid triangles: $a < b + c \ b < a + c \ and \ c < a + b$







Q:3 Do you have a test case with three integers greater than zero such that the sum of two numbers is equal to the third?







Q:3 Do you have a test case with three integers greater than zero such that the sum of two numbers is equal to the third? (4,2,2) is invalid triangle with equal sum.



Fulfill above definition, but is invalid: a < b+c and b < a+c and c < a+b







Do you have a test case:

- 4. with some permutations of previous test? (2,4,2) (2,2,4)
- 5. that represents a valid scalene triangle? (3,4,5)
- 6. that represents a valid equilateral triangle? (3,3,3)
- 7. that represents a valid isosceles triangle? (4,3,3)
- 8. with some permutations of previous test? (3,4,3) (3,3,4)
- 9. in which one side has a zero value? (0,4,3)
- 10. in which one side has a negative value? (-1,4,3)
- 11. in which all sides are zero? (0,0,0)
- 12. specifying at least one noninteger value? (2,2.5,4)
- 13. specifying the wrong number of values? (2,3) or (2,3,5,4)

Q:14 For each test case did you specify the expected output from the program in addition to the input values?





About the example



- A set of test case that satisfies these conditions does not guarantee that all possibile errors would be found.
- An adeguate test of this program should expose at least these errors.
- Higly qualified professional programmers score, on the average, 7.8 out of a possibile 14.







Syntax

- Definition: errors in grammar (violations of the "rules" for forming legal language statements).
- Examples: undeclared identifiers, mismatched parentheses in an expression, an opening brace without a matching closing brace, etc.
- Occur: an error that is caught in the process of compiling the program.







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- Occur: an error that is caught in the process of compiling the program.

The easiest errors to spot by a compiler or some aid in checking the syntax.

As you get more practice using a language, you naturally make fewer errors, and will be able to quickly correct those that do occur.





Runtime

- Definition: "Asking the computer to do the impossible!"
- Examples: division by zero, taking the square root of a negative number, referring to the 101th on a list of only 100 items, deferencing a null pointer, etc.
- Occur: an error that is not detected until the program is executed, and then causes a processing error to occur.







Runtime

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- Examples: division by zero, taking the square root of a negative number, referring to the 101th on a list of only 100 items, deferencing a null pointer, etc.
- Occur: an error that is not detected until the program is executed, and then causes a processing error to occur. They are not easy to spot because they are not syntax or grammar errors, they are subtle errors and develop in the course of program's execution.
 - Avoiding exceptions and correcting program behaviour is also largely a matter of experience.
 - To prevent use defensive programming.







Logic (semantic, meaning)

- Definition: the program compiles (no syntax errors) and runs to a normal completion (no runtime erros), but the output is wrong. A statement may be syntactically correct, but mean something other than what we intended. Therefore it has a different effect, causing the program output to be wrong.
- Examples: improper initialization of variables, performing arithmetic operations in the wrong order, using an incorrect formula to compute a value, etc.
- Occur: an error that affect how the code works.







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- Examples: improper initialization of variables, performing arithmetic operations in the wrong order, using an incorrect formula to compute a value, etc.
- Occur: an error that affect how the code works. It is a type of error that only the programmer can recognize.Finding and correcting logic errors in a program is known as debugging.









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Definitions



Static analysis refers to a method of examinig software that allows developers to discover dangerous programming pratices or potential errors in source code, without actually run the code.





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- Using compiler options.
- Using static analyzer.





Compiler and errors









Compiler and errors





Compiler checks:

- Syntax.
- Semantic.





Lexical (Syntactic) errors



- Characters in the source that aren't in the alphabet of the language.
- Words in the source that aren't in the vocabulary of the language.





Syntactic errors



- Comment delimiters that have been put in wrong place or omitted.
- Literal delimiters that have been put in wrong place or omitted.
- Keywords that have been misspelled.
- Required punctuation that is missing.
- Construct delimiters such as parentheses or braces that have been missplaced.
- Blank or tab characters that are missing.
- Blank or tab characters that shouldn't occur where they're found.





Semantic errors



- Names that aren't declared.
- Operands of the wrong type for the operator they're used with.
- Values that have the wrong type for the name to which they're assigned.
- Procedures that are invoked with the wrong number of arguments.
- Procedures that are invoked with the wrong type of arguments.
- Function return values that are the wrong type for the context in which they're used.



Semantic errors



- Code blocks that are unreachable.
- Code blocks that have no effect.
- Local variables that are used before being initialized or assigned.
- Local variables that are initialized or assigned but not used.
- Procedures that are never invoked.
- Procedures that have no effect.
- Global variables that are used before being initialized or assigned.
- ► Global variables that are initialized or assigned, but not used.





Compilers



- Not all compilers find the same defects.
- The more information a compilers has, the more defects it can find.
- Some compilers operate in "forgiving" mode but have "strict" or "pedantic" mode, if you request it.









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Static analysis:example 1 part one



```
1 #include <stdio.h>
2 int main (void)
3 {
4     printf ("Two plus two is %f\n", 4);
5     return 0;
6 }
```





Static analysis:example 1 part one



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```

<ruggiero@shiva ~/CODICI>gcc bad.c

<ruggiero@shiva ~/CODICI>./a.out

Two plus two is 0.000000





Static analysis:example 1 part two



<ruggiero@shiva ~/CODICI>gcc -Wall bad.c

```
bad.c: In function main:
bad.c:4: warning: format '%f' expects type 'double',
but argument 2 has type 'int'
```





Static analysis:example 1 part two



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```





gcc options



-Wall

turns on the most commonly-used compiler warnings option

-Waddress -Warray-bounds (only with -O2) -Wc++0x-compat -Wchar-subscripts -Wimplicit-int -Wimplicit-function-declaration -Wcomment -Wformat -Wmain (only for C/ObjC and unless -ffreestanding) -Wmissing-braces -Wnonnull -Wparentheses -Wpointer-sign -Wreorder -Wreturn-type -Wsequence-point -Wsign-compare (only in C++) -Wstrict-aliasing -Wstrict-overflow=1 -Wswitch -Wtrigraphs -Wuninitialized -Wunknown-pragmas -Wunused-function -Wunused-label -Wunused-value -Wunused-variable -Wvolatile-register-var




Oher compilers



[ruggiero@matrix1 ~]\$ pgcc bad.c





Oher compilers



[ruggiero@matrix1 ~]\$ pgcc bad.c

[ruggiero@matrix1 ~]\$ icc bad.c

bad.c(4): warning #181: argument is incompatible
with corresponding format string conversion
 printf ("Two plus two is %f\n", 4);
 ^





check.c:source



```
#include <stdio.h>
1
2
    int main (void)
3
4
               return 0;
5
6
    ł
    void checkVal(unsigned int n) {
7
8
               if (n < 0) {
                             /* Do something... */
9
10
                 else if (n \ge 0) {
11
                               /* Do something else... */
12
13
                             }
14
    }
```





check.c:compilation



<ruggiero@shiva:~> gcc -Wall check.c





check.c:compilation



<ruggiero@shiva:~> gcc -Wall check.c

<ruggiero@shiva:~> gcc -Wall -Wextra check.c

check.c: In function ?checkVal?: check.c:8: warning: comparison of unsigned expression < 0 is always false check.c:11: warning: comparison of unsigned expression >= 0 is always true





gcc options



-Wextra (-W) reports the most common programming errors and less-serious but potential problem

```
-Wclobbered -Wempty-body
-Wignored-qualifiers -Wmissing-field-initializers
-Wmissing-parameter-type (C only) -Wold-style-declaration (C only)
-Woverride-init -Wsign-compare
-Wtype-limits -Wuninitialized (only with -O1 and above)
-Wunused-parameter (only with -Wunused or -Wall)
```





check1.c:source

y;



1	<pre>#include <stdio.h></stdio.h></pre>
2	int main (void)
3	{
4	double $x = 10.0;$
5	double $y = 11.0;$
6	double $z = 0.0;$
7	if (x == y) {
8	z = x *
9	}
10	return 0;
11	}





check1.c:source



1	<pre>#include <stdio.h></stdio.h></pre>
2	int main (void)
3	{
4	double $x = 10.0;$
5	double $y = 11.0;$
6	double $z = 0.0;$
7	if (x == y) {
8	$\mathbf{z} = \mathbf{x} \star \mathbf{y};$
9	}
10	return 0;
11	}

<ruggiero@shiva:~> gcc -Wall -Wextra check1.c

<ruggiero@shiva:~> gcc -Wall -Wextra -Wfloat-equal check1.c

```
checkl.c: In function main:
checkl.c:8: warning: comparing floating point
with == or != is unsafe
```



gcc: other compiler options



- -Wno-div-by-zero -Wsystem-headers -Wfloat-equal -Wtraditional (C only)
- -Wdeclaration-after-statement (C only) -Wundef -Wno-endif-labels -Wshadow
- -Wlarger-than-len -Wpointer-arith -Wbad-function-cast (C only) -Wcast-align
- -Wwrite-strings -Wconversion -Wsign-compare -Waggregate-return -Wstrict-prototypes
- (C only) -Wold-style-definition (C only) -Wmissing-prototypes (C only)
- -Wmissing-declarations (C only) -Wcast-qual -Wmissing-field-initializers
- -Wmissing-noreturn -Wmissing-format-attribute -Wno-multichar
- -Wno-deprecated-declarations -Wpacked -Wpadded -Wredundant-decls
- -Wnested-externs (C only) -Wvariadic-macros -Wunreachable-code -Winline
- -Wno-invalid-offsetof (C++ only) -Winvalid-pch -Wlong-long -Wdisabled-optimization
- -Wno-pointer-sign

	SCAI	testinit.c:source	SAPIENZA UNIVERSITÀ DI ROMA
1	<pre>int main() {</pre>		
2	int v[16];		
3	int i,j,k;		
4	j=i;		
5	v[i]= 42;		
6	return 0 ;		
7	}		



	SCAI	testinit.c:source	SAPIENZA UNIVERSITÀ DI ROMA
1	<pre>int main() {</pre>		
2	int v[16];		
3	int i, j, k;		
4	j=i;		
5	v[i]= 42;		
6	return 0 ;		
7	}		

ruggiero@shiva:~> gcc -Wall -Wextra testinit.c



CINECA	SCAI	testinit.c:source	SAPIENZA UNIVERSITÀ DI ROMA
Supercompt	int main() {		
	int main() (
2	int v[16];		
3	int i,j,k;		
4	j=i;		
5	v[i]= 42;		
6	return 0 ;		
7	}		

ruggiero@shiva:~> gcc -Wall -Wextra testinit.c

ruggiero@shiva:~> gcc -01 -Wall -Wextra -Wuninitialized testinit.c

testinit.c: In function main testinit.c:4: warning: unused variable k testinit.c:5: warning: i is used uninitialized in this function



	SCAI	testinit.c:source	SAPIENZA UNIVERSITÀ DI ROMA
SuperCompu	int main () (
	fint main() {		
2	int v[16];		
3	int i,j,k;		
4	j=i;		
5	v[i]= 42;		
6	return 0 ;		
7	}		

ruggiero@shiva:~> gcc -Wall -Wextra testinit.c

ruggiero@shiva:~> gcc -01 -Wall -Wextra -Wuninitialized testinit.c

testinit.c: In function main testinit.c:4: warning: unused variable k testinit.c:5: warning: i is used uninitialized in this function

[ruggiero@matrix1 ~]\$ icc testinit.c

testinit.c(5): warning #592: variable "i" is used before its value is set j=i;







program par implicit none
integer, parameter :: hacca=10
call sub(hacca)
end
subroutine sub(hacca)
implicit none
integer hacca
hacca=hacca+1
write(*,*) hacca
return
end







> xlf par.f -o par.x

> ./par.x

11









> ./par.x

11

<ruggiero@ife2 ~>ifort par.f -o par.x



SCA Compilation: default options



> xlf par.f -o par.x

> ./par.x

SuperComputing Applications and Innovation

11

<ruggiero@ife2 ~>ifort par.f -o par.x

<ruggiero@ife2 ~> ./par.x

forrtl: se	evere (174): SIGSEGV,	segmentation fault o	ccurred	
Image	PC	Routine	Line Sou	irce
par.x	0000000004)2688 Unknown	Unknown Unk	nown
par.x	0000000004	02670 Unknown	Unknown Unk	nown
par.x	0000000004)262A Unknown	Unknown Unk	nown
libc.so.6	0000036FF8	LC40B Unknown	Unknown Unk	nown
par.x	0000000004	0256A Unknown	Unknown Unk	nown







<ruggiero@ife2 ~> man ifort

? -assume noprotect_constants

Tells the compiler to pass a copy of a constant actual argument. This copy can be modified by the called routine, even though the Fortran standard prohibits such modification. The calling routine does not see any modification to the constant. The default is -assume protect_constants, which passes the constant actual argument. Any attempt to modify it results in an error.







<ruggiero@ife2 ~> man ifort

? -assume noprotect_constants

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<ruggiero@ife2 ~>ifort par.f -assume noprotect_constants -o par.x





<ruggiero@ife2 ~> man ifort

? -assume noprotect_constants

Tells the compiler to pass a copy of a constant actual argument. This copy can be modified by the called routine, even though the Fortran standard prohibits such modification. The calling routine does not see any modification to the constant. The default is -assume protect_constants, which passes the constant actual argument. Any attempt to modify it results in an error.

<ruggiero@ife2 ~>ifort par.f -assume noprotect_constants -o par.x

<ruggiero@ife2 ~>./par.x

11









Bugs and Prevention

Testing

Static analysis Compiler options Static analyzer

Run-time analysis









- Open Source Static Analysis Tool developed at University of Virginia by Professor Dave Evans
- Based on Lint
- www.splint.org
- ▶ splint [-option -option ...] filename [filename ...]



splint can detect with just source code...



Unused declarations.

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- Type inconsistencies.
- Variables used before being assigned.
- Function return values that are ignored.
- Execution paths with no return.
- Switch cases that fall through.
- Apparent infinite loops.



SAPI

- Dereferencing pointers with possible null values.
- Using storage that is undefined or partly undefined.
- Returning storage that is undefined or partly defined.
- Type mismatches.
- Using deallocated storage.

CAsplint can detect with annotation information



Memory leaks.

outing Applications and Innovation

- Inconsistent modification of caller visible states.
- Violations of information hiding.
- Undefined program behaviour due to evaluation order, incomplete logic, infinite loops, statements with no effect, and so on.
- Problematic uses of macros.





Static analysis:example 1 part one



```
1 #include <stdio.h>
2 int main (void)
3 {
4     printf ("Two plus two is %f\n", 4);
5     return 0;
6 }
```





check.c:source



```
#include <stdio.h>
1
2
    int main (void)
3
4
               return 0;
5
6
    ł
    void checkVal(unsigned int n) {
7
8
               if (n < 0) {
                             /* Do something... */
9
10
                 else if (n \ge 0) {
11
                               /* Do something else... */
12
13
                             }
14
    }
```





splint:examples



<ruggiero@shiva:~> splint-3.1.2/bin/splint bad.c

Finished checking --- no warnings





splint:examples



<ruggiero@shiva:~> splint-3.1.2/bin/splint bad.c

Finished checking --- no warnings

<ruggiero@shiva:~> splint-3.1.2/bin/splint check.c

```
splint 3.1.2 --- 28 Mar 2008
check.c: (in function checkVal)
check.c:8:15: Comparison of unsigned value involving zero: n < 0
An unsigned value is used in a comparison with zero in a way that is
a bug or confusing. (Use -unsignedcompare to inhibit warning)
check.c:11:22: Comparison of unsigned value involving zero: n >= 0
Finished checking --- 2 code warnings
```





check1.c:source



1	<pre>#include <stdio.h></stdio.h></pre>
2	int main (void)
3	{
4	double $x = 10.0;$
5	double $y = 11.0;$
6	double $z = 0.0;$
7	if (x == y) {
8	$\mathbf{z} = \mathbf{x} \star \mathbf{y};$
9	}
10	return 0;
11	}

<ruggiero@shiva:~> gcc -Wall -Wextra check1.c

<ruggiero@shiva:~> gcc -Wall -Wextra -Wfloat-equal check1.c

```
checkl.c: In function main:
checkl.c:8: warning: comparing floating point
with == or != is unsafe
```





splint:examples



<ruggiero@shiva:~> splint-3.1.2/bin/splint check1.c

```
Splint 3.1.2 --- 28 Mar 2008
```

check1.c: (in function main) check1.c:8:14: Dangerous equality comparison involving double types: x == y Two real (float, double, or long double) values are compared directly using == or != primitive. This may produce unexpected results since floating point representations are inexact. Instead, compare the difference to FLT_EPSILON or DBL_EPSILON. (Use -realcompare to inhibit warning)

Finished checking --- 1 code warning





check1.c: rewritten



1	int main (void)
2	{
3	double $x = 10.0;$
4	double y = 11.0;
5	double $z = 0.0;$
6	double norma=0.1e-16;
7	double epsilon;
8	epsilon=x-y;
9	if (epsilon < norma) {
10	$z = x \star y;$
11	}
12	
13	return 0;
14	}





average.c: source



```
#include <stdio.h>
1
    int main (void) {
2
      int size = 5;
3
4
      int a;
      float total;
5
6
      float art[size];
7
8
     for(a = 0 ; a < size; ++ a) {</pre>
     art[a] = (float) a;
9
      total += art[a]; }
10
      printf(" %f\n" , total / (float) size);
11
      return 0;
12
13
    }
```







<ruggiero@shiva:~> gcc -Wall -Wextra average.c

<ruggiero@shiva:~> ./a.out

1.999994







<ruggiero@shiva:~> gcc -Wall -Wextra average.c

<ruggiero@shiva:~> ./a.out

1.999994

<ruggiero@shiva:~> gcc -Wall -Wextra -O2 average.c

average.c: In function main: average.c:5: warning: total may be used uninitialized in this function

<ruggiero@shiva:~> ./a.out

CINECA

nan



average.c: splint



<ruggiero@shiva:~> splint-3.1.2/bin/splint average.c

Splint 3.1.2 --- 28 Mar 2008

average.c: (in function main) average.c:10:4 Variable total used before definition An rvalue is used that may not be initialized to a value on some execution path. (Use -usedef to inhibit warning) average.c:11:20: Variable total used before definition

Finished checking --- 2 code warnings




average.c: new version



```
#include <stdio.h>
1
    int main (void) {
2
      int size = 5;
3
4
      int a;
      float total;
5
6
      float art[size];
7
      total=0;
8
     for(a = 0 ; a < size; ++ a) {</pre>
       art[a] = (float) a;
9
      total += art[a]; }
10
     printf(" %f\n" , total / (float) size);
11
      return 0;
12
13
    }
```





assign.c: source and compilation



```
#include <stdio.h>
1
2
   main()
3
   ł
          int a=0;
4
          while (a=1)
5
               printf("hello\n");
6
7
          return 0;
8
   }
```





assign.c: source and compilation



```
#include <stdio.h>
1
2
   main()
3
   ł
4
          int a=0:
          while (a=1)
5
               printf("hello\n");
6
          return 0;
7
8
   }
```

<ruggiero@shiva:~> gcc -Wall -Wextra assign.c

```
assigna.c:3: warning: return type defaults to int
assign.c: In function main:
assign.c:5: warning: suggest parentheses around assignment used as
truth value
```





assign.c: splint



<ruggiero@shiva:~>splint-3.1.2/bin/splint assign.c

Splint 3.1.2 --- 28 Mar 2008
assign.c: (in function main)
assign.c:(in function main)
The condition test is an assignment expression. Probably, you mean
to use == instead of =. If an assignment is intended, add an extra
parentheses nesting (e.g., if ((a = b)) ...) to suppress this message.
(Use -predassign to inhibit warning)
assign.c:5:14: Test expression for while not boolean, type int: a=1
Test expression type is not boolean or int. (Use -predboolint
to inhibit warning)
Finished checking --- 2 code warnings





memory.c: source and compilation



```
1
    #include <stdlib.h>
2
    int main()
3
4
    ł
             int *p = malloc(5*sizeof(int));
5
6
7
8
9
             *p = 1;
             free(p);
10
             return 0;
11
12
    }
```





memory.c: source and compilation



1	
2	<pre>#include <stdlib.h></stdlib.h></pre>
3	int main()
4	-{
5	<pre>int *p = malloc(5*sizeof(int));</pre>
6	
7	
8	
9	*p = 1;
0	free (p) ;
11	return 0;
12	}

<ruggiero@shiva:~> gcc -Wall -Wextra memory.c





memory.c: splint



<ruggiero@shiva:~> splint-3.1.2/bin/splint memory.c

Splint 3.1.2 --- 28 Mar 2008

memory.c: (in function main)
memory.c:9:10: Dereference of possibly null pointer p: *p
A possibly null pointer is dereferenced. Value is either
the result of a function which may return null (in which
case, code should check it is not null), or a global,
parameter or structure field declared with the null
qualifier. (Use -nullderef to inhibit warning)
memory.c:5:18: Storage p may become null

Finished checking --- 1 code warning





memory.c: new version



```
#include <stdlib.h>
1
   #include <stdio.h>
2
   int main()
3
    ł
4
            int *p = malloc(5*sizeof(int));
5
            if (p == NULL) {
6
                 fprintf(stderr, "error in malloc");
7
                 exit(EXIT_FAILURE);
8
             } else *p = 1;
9
            free(p);
10
            return 0;
11
12
    }
```





out_b.c: source



1	<pre>#include <stdio.h></stdio.h></pre>
2	#define N 5
3	int main (void)
4	{
5	<pre>int t[N];</pre>
6	int i;
7	
8	i=6;
9	t[i] = i+1;
10	
11	return 0;
12	}





out_b.c: splint



<ruggiero@shiva:~> splint-3.1.2/bin/splint out_b.c





out_b.c: splint



<ruggiero@shiva:~> splint-3.1.2/bin/splint out_b.c

<ruggiero@shiva:~> splint-3.1.2/bin/splint +bounds out_b.c

```
Splint 3.1.2 --- 28 Mar 2008
out_b.c: (in function main)
out_b.c:9:9: Likely out-of-bounds store: t[i]
    Unable to resolve constraint:
    requires 4 >= 6
    needed to satisfy precondition:
    requires maxSet(t @ out_b.c:9:9) >= i @ out_b.c:9:11
    A memory write may write to an address beyond the allocated buffer. (Use
    -likelyboundswrite to inhibit warning)
```

Finished checking --- 1 code warning





error.f: source



1		REAL FUNCTION COMPAV (SCORE, COUNT)
2		INTEGER SUM, COUNT, J, SCORE (5)
3		DO 30 I = 1, COUNT
4		SUM = SUM + SCORE(I)
5	30	CONTINUE
6		COMPAV = SUM/COUNT
7		write(*,*) compav
8		END
9		PROGRAM AVENUM
10		PARAMETER (MAXNOS=10)
11		INTEGER I, COUNT
12		REAL NUMS (MAXNOS), AVG
13		COUNT = 0
14		DO 80 I = 1, MAXNOS
15		READ (5, *, END =100) NUMS(I)
16		COUNT = COUNT + 1
17	80	CONTINUE
18	100	AVG = COMPAV (NUMS, COUNT)
19		END







<ruggiero@ife2 /ftnchek>gfortran -Wall -Wextra -O2 error.f

In file error.f:2

```
INTEGER SUM,COUNT,J,SCORE(5)
1
CWarning: Unused variable j declared at (1)
error.f: In function compav:
error.f:2: warning: sum may be used uninitialized in this function
```







<ruggiero@ife2/ftnchek>pgf90 -Minform=inform error.f

PGF90-I-0035-Predefined intrinsic sum loses intrinsic property (error.f: 5) PGF90-I-0035-Predefined intrinsic count loses intrinsic property (error.f: 16)







<ruggiero@ife2 /ftnchek>ifort -warn all error.f

fortcom:	Warning: error.f, line 4: This name has not been given an explicit type. DO 30 I = 1,COUNT	[I]
fortcom:	<pre>Info: error.f, line 2: This variable has not been used. [J] INTEGER SUM,COUNT,J,SCORE(5)</pre>	
fortcom:	Warning: error.f, line 13: This name has not been given an explicit type. PARAMETER(MAXNOS=10)	[MAXNOS]
fortcom: 100	Warning: error.f, line 21: This name has not been given an explicit type. AVG = COMPAV(NUMS, COUNT)	[COMPAV]







<ruggiero@ife2 /ftnchek>g95 -Wall -Wextra error.f

```
In file error.f:6
 30
             CONTINUE
Warning (142): Nonblock DO statement at (1) is obsolescent
In file error.f:1
      REAL FUNCTION COMPAV (SCORE, COUNT)
Warning (163): Actual argument 'score' at (1) does not have an INTENT
In file error.f:1
      REAL FUNCTION COMPAV (SCORE, COUNT)
Warning (163): Actual argument 'count' at (1) does not have an INTENT
In file error.f:2
             INTEGER SUM, COUNT, J, SCORE (5)
Warning (137): Variable 'j' at (1) is never used and never set
```







```
In file error.f:20
80
            CONTINUE
Warning (142): Nonblock DO statement at (1) is obsolescent
In file error.f:21
            AVG = COMPAV (NUMS, COUNT)
Warning (165): Implicit interface 'compav' called at (1)
In file error.f:15
           REAL NUMS (MAXNOS), AVG
Warning (112): Variable 'avg' at (1) is set but never used
In file error.f:21
100
            AVG = COMPAV (NUMS, COUNT)
In file error.f:1
      REAL FUNCTION COMPAV (SCORE, COUNT)
Warning (155): Inconsistent types (REAL(4)/INTEGER(4))
in actual argument lists at (1) and (2)
```





SAPIENZA UNIVERSITÀ DI ROMA

<ruggiero@ife2/ftnchek>./bin/ftnchek error.f

```
FINCHEK Version 3.3 November 2004
File error f.
                     COMPAV = SUM/COUNT
Warning near line 7 col 21 file error.f:
 integer quotient expr SUM/COUNT converted to real
Warning in module COMPAV in file error.f:
Variables declared but never referenced:
    J declared at line 2 file error.f
Warning in module COMPAV in file error.f:
 Variables may be used before set:
    SUM used at line 5 file error.f
    SUM set at line 5 file error.f
Warning in module AVENUM in file error.f:
Variables set but never used:
    AVG set at line 21 file error.f
```



SCAI www.dsm.fordham.edu/ftnchek

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```
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```

```
0 syntax errors detected in file error.f
4 warnings issued in file error.f
Warning: Subprogram COMPAV argument data type mismatch at position 1:
    Dummy arg SCORE in module COMPAV line 1 file
error.f is type intg
Actual arg NUMS in module AVENUM line 21 file
error.f is type real
Warning: Subprogram COMPAV argument arrayness mismatch at position 1:
    Dummy arg SCORE in module COMPAV line 1 file
error.f has 1 dim size 5
    Actual arg NUMS in module AVENUM line 21 file
error.f has 1 dim size 10
```





. . .

Static analyzer for Fortran



- Forcheck can handle up Fortran 95 and some Fortran 2003. item Cleanscape FortranLint can handle up to Fortran 95.
- plusFORT is a multi-purpose suite of tools for analyzing and improving Fortran programs.







- ► 40% false positive reports of correct code.
- ► 40% multiple occurence of same problem.
- ► 10% minor or cosmetic problems.
- ► 10% serious bugs, very hard to find by other methods.

From **"The Developer's Guide to Debugging"** T. Grotker, U. Holtmann, H. Keding, M. Wloka





Static analyzer: Lessons learned



- Do not ignore compiler warnings, even if they appear to be harmless.
- Use multiple compilers to check the code.
- ► Familiarize yourself with a static checker.
- ► Reduce static checker errors to (almost) zero.
- Rerun all test cases after a code cleanup.
- Doing regular sweeps of the source code will pay off in long term.

From **"The Developer's Guide to Debugging"** T. Grotker, U. Holtmann, H. Keding, M. Wloka









Bugs and Prevention

Testing

Static analysis

Run-time analysis Memory checker





Runtime signals



- When a job terminates abnormally, it usually tries to send a signal (exit code) indicating what went wrong.
- ► The exit code from a job is a standard OS termination status.
- ► Typically, exit code 0 (zero) means successful completion.
- Your job itself calling exit() with a non-zero value to terminate itself and indicate an error.
- The specific meaning of the signal numbers is platform-dependent.



Runtime signals



You can find out why the job was killed using:

[ruggiero@matrix1 ~]\$ kill -1

```
1)
   SIGHUP
            2) SIGINT
                       3) SIGQUIT
                                    4)
                                       SIGILL
 5)
    SIGTRAP
             6)
                SIGABRT
                          7)
                             SIGBUS
                                     8)
                                        SIGFPE
    SIGKILL 10) SIGUSR1 11) SIGSEGV 12) SIGUSR2
 9)
13)
    SIGPIPE 14) SIGALRM 15) SIGTERM 16) SIGSTKFLT
    SIGCHLD 18) SIGCONT 19) SIGSTOP 20) SIGTSTP
17)
21)
    SIGTTIN 22) SIGTTOU 23) SIGURG 24) SIGXCPU
25)
    SIGXFSZ 26) SIGVTALRM 27) SIGPROF 28) SIGWINCH
29) SIGIO 30) SIGPWR 31) SIGSYS 34) SIGRTMIN
. . . .
```





Runtime signals



To find out what all the "kill -l" words mean:

[ruggiero@matrix1 ~]\$ man 7 signal

Signal	Value	Action	Comment
SIGHUP	1	Term	Hangup detected on controlling terminal or death of controlling process
SIGINT	2	Term	Interrupt from keyboard
SIGQUIT	3	Core	Quit from keyboard
SIGILL	4	Core	Illegal Instruction
SIGABRT	6	Core	Abort signal from abort(3)





Action description



- Term Default action is to terminate the process.
- Ign Default action is to ignore the signal.
- Core Default action is to terminate the process and dump the core.
- Stop Default action is to stop the process.
- Cont Default action is to continue the process if is currently stopped.





Common runtime signals



Signal name	OS signal name	Description
Floating point exception	SIGFPE	The program attempted
		an arithmetic operation
		with values that
		do not make sense
Segmentation fault	SIGSEGV	The program accessed
		memory incorrectly
Aborted	SIGABRT	Generated by the runtime
		library of the program
		or a library it uses,
		after having detecting
		a failure condition.





FPE example









FPE example





[ruggiero@matrix1 ~]\$ gcc fpe_example.c

[ruggiero@matrix1 ~]\$./a.out

Floating exception





SEGV example









SEGV example



```
1 main()
2 {
3 int array[5]; int i;
4 for(i = 0; i < 255; i++) {
5 array[i] = 10;}
6 return 0;
7 }</pre>
```

[ruggiero@matrix1 ~]\$ gcc segv_example.c

[ruggiero@matrix1 ~]\$./a.out

Segmentation fault





ABORT example









ABORT example





[ruggiero@matrix1 ~]\$ gcc abort_example.c

[ruggiero@matrix1 ~]\$./a.out

```
a.out: abort_example.c:5: main: Assertion `i!=0' failed.
Abort
```





Common runtime errors



- Allocation Deallocation errors (AD).
- Array conformance errors (AC).
- Array Index out of Bound (AIOB).
- Language specific errors (LS).
- Floating Point errors (FP).
- Input Output errors (IO).
- Memory leaks (ML).
- Pointer errors (PE).
- String errors (SE).
- Subprogram call errors (SCE).
- Uninitialized Variables (UV).





Useful link



- Iowa State University's High Performance Computing Group
- ▶ Run Time Error Detection Test Suites for Fortran, C, and C++
- http://rted.public.iastate.edu


SCA Grading Methodology: score



- ▶ 0.0: is given when the error was not detected.
- ► 1.0: is given for error messages with the correct error name.
- 2.0: is given for error messages with the correct error name and line number where the error occurred but not the file name where the error occurred.
- 3.0: is given for error messages with the correct error name, line number and the name of the file where the error occurred.
- 4.0: s given for error messages which contain the information for a score of 3.0 but less information than needed for a score of 5.0.
- 5.0: is given in all cases when the error message contains all the information needed for the quick fixing of the error.





Grading Methodology : an example



!**** ! coj ! Jai	**************************************	**************************************
! and ! Li	d Ziyu Zhang, All r censed under the Edu	ights reserved. ucational Community License version 1.0.
! Se	e the full agreement	t at http://rted.public.iastate.edu/ .
! * * * * ! * * * *	***************************************	***************************************
1	Name of the test	F H 1 1 b f90
1	Name of the cest.	
1	Summary:	allocation/deallocation error
! !	Test description:	deallocation twice for allocatable array in a subroutine
1		contains in a main program
1	Support files:	Not needed
!	Env. requirements:	Not needed
! ! !	Keywords:	deallocation error subroutine contains in a main program





Grading Methodology: an example



```
1
     Last modified:
                       1/17/2005
     Programmer:
                      Ying Xu, Iowa State Univ.
-----
              ****************
                                             *****
program tests
  implicit none
 integer :: n=10, m=20
 double precision :: var
 call sub(n,m,var)
 print *, var
 contains
   subroutine sub(n.m.var)
     integer, intent(in) :: n,m
     double precision, intent(inout) :: var
     double precision, allocatable :: arr(:,:)
                                               ! DECLARE
```





Grading Methodology: an example







Grading Methodology: an example



Real message (grade 1.0)

Fortran runtime error: Internal: Attempt to DEALLOCATE unallocated memory.

Ideal message (grade 5.0)

ERROR: unallocated array At line 52 column 17 of subprogram 'sub' in file 'F_H_1_1_b.f90', the argument 'arr' in the DEALLOCATE statement is an unallocated array. The variable is declared in line 41 in subprogram 'sub' in file 'F_H_1_1_b.f90'.





Fortran Results



Compiler	AC	A D	AIOB	LS	FP	10
gcc-4.3.2	1	0.981481	3.40025	2.88235	0	2.33333
gcc-4.3.2	1	1.38889	3.40025	2.88235	0	2.33333
gcc-4.4.3	1	1.38889	3.40025	2.88235	0	2.33333
gcc-4.6.3	1	1.27778	0.969504	0.94117	0	2.33333
gcc-4.7.0	1	1.38889	0.969504	0.94117	0.714286	2.33333
g95.4.0.3	0.421053	1.22222	3.60864	3.82353	0.571428	2.66667
intel-10.1.021	0.421053	1.42593	3.45362	2.82353	0.571428	2.11111
intel-11.0.074	0.421053	1.68519	3.446	2.82353	0.571428	2.11111
intel-11.1	1	1.90741	3.47649	2.88235	1.42857	2.33333
intel-12.0.2	0.421053	1.62963	3.44727	2.82353	0.571428	2.11111
intel-14.0.1	0.421053	1.62963	3.44854	2.82353	0.571428	2.11111
open64.4.2.3	3	0.888889	2.63405	3	0	1
pgi-7.2-5	0.421053	0.388889	3.8526	3.82353	0	2.44444
pgi-8.0-1	0.421053	0.388889	3.8526	3.82353	0	2.44444
pgi-10.3	0.421053	0.388889	3.8526	3.82353	0	2.44444
pgi-11.8	0.421053	0.388889	3.8526	3.82353	0	2.44444
pgi-12.8	1	1	3.8831	3.82353	1	2.61111
pgi-13.10	1	1	3.8831	3.82353	1	2.72222
pgi-14.1.0	0.421053	0.5	3.8526	3.82353	0	2.61111
sun.12.1	3	2.77778	3.00381	3	2	2.16667
sun.12.1+bcheck	3	2.77778	3.03431	3	0.285714	2.16667





Fortran Results



Compiler	ML	PE	SE	SCE	UV
gcc-4.3.2	0	3.49609	3.25	0	0.0159236
gcc-4.3.2	0	3.49609	3.25	0	0.0286624
gcc-4.4.3	0	3.49609	3.25	0	0.0286624
gcc-4.6.3	0	1	3.25	0.166667	0.22293
gcc-4.7.0	0	1	3.25	0.166667	0.130573
g95.4.0.3	1	3	3.43333	0	0.0159236
intel-10.1.021	0	3.5625	0	0.166667	0.286624
intel-11.0.074	0	3.55469	0	0.166667	0.299363
intel-11.1	1	3.55469	1	1	1.07643
intel-12.0.2	0	3.55469	0	0.166667	0.292994
intel-14.0.1	0	3.55469	0	0.166667	0.292994
open64.4.2.3	0	3.3625	0	0.0833333	0.286624
pgi-7.2-5	0	4	0	0	0.022293
pgi-8.0-1	0	4	0	0	0.022293
pgi-10.3	0	4	0	0	0.0254777
pgi-11.8	0	4	0	0	0.0127389
pgi-12.8	1	4	1	1	1
pgi-13.10	1	4	1	1	1
pgi-14.1.0	0	4	0	0	0.143312
sun.12.1	0	3.03125	3	0	0.022293
sun.12.1+bcheck	1.25	3.03125	3	1	0.640127









Compiler	AD	AloB	LS	FP	10
gcc-4.3.2	0.44	0.00925926	0.0416667	0.2	0.0666667
gcc-4.4.3	0.44	0.00925926	0.0416667	0.2	0.0666667
gcc-4.6.3	0.44	0.00925926	0.0416667	0.2	0.2
gcc-4.7.0	0.44	0.00925926	0.0416667	0.2	0.2
intel-10.1.021	0.52	0.00925926	0.0416667	0	0.2
intel-11.0.074	0.52	0.00925926	0.0416667	0	0.2
intel-11.1	0.68	1	1	1	1
intel-12.0.2	0.44	0.00925926	0.0416667	0	0.2
intel-14.0.1	0.4	0.00925926	0.0416667	0	0.2
open64-4.2.3	0.52	0.00925926	0.0416667	0	0.2
pgi-7.2-5	0.44	0.00925926	0.0416667	0	0.2
pgi-8.0-1	0.44	0.00925926	0.0416667	0	0.2
pgi-10.3	0.44	0.00925926	0.0416667	0	0.2
pgi-11.8	0.44	0.00925926	0.0416667	0	0.2
pgi-12.8	0.68	2.40741	2.16667	0	1
pgi-13.10	0.68	2.40741	2.16667	0	1
pgi-14.1-0	0.48	1.93519	1.625	0	0.2
sun-12.1	0.44	0.00925926	0	0	0.2
sun-12.1+bcheck	0.16	0	0	0	0

CINECA



C Results



Compiler	ML	PE	SE	SCE	UV
gcc-4.3.2	0.0166667	0.0166667	0.05	0	0
gcc-4.4.3	0.0166667	0.0166667	0.05	0	0
gcc-4.6.3	0.0166667	0.0166667	0.05	0	0
gcc-4.7.0	0.0166667	0.0166667	0.05	0	0
intel-10.1.021	0.0666667	0.0166667	0.05	0	0
intel-11.0.074	0.0666667	0.0166667	0.05	0	0
intel-11.1	1	1	1	1	1
intel-12.0.2	0.0666667	0.0166667	0.05	0	0
intel-14.0.1	0.133333	0.0166667	0.05	0	0
open64-4.2.3	0.0666667	0.0166667	0.05	0	0
pgi-7.2-5	0.0666667	0.0166667	0.05	0	0
pgi-8.0-1	0.0666667	0.0166667	0.05	0	0
pgi-10.3	0.0666667	0.0166667	0.05	0	0
pgi-11.8	0.0666667	0.0166667	0.05	0	0
pgi-12.8	1	1	1	1	1
pgi-13.10	1	1	1	1	1
pgi-14.1-0	0.133333	0.0166667	0.05	0	0
sun-12.1	0.0666667	0.0166667	0.05	0	0
sun-12.1+bcheck	1.13333	0.025	0	0	0





C++ Results



Compiler	AD	AloB	FP	IO	ML	PE	SE	UV
gcc-4.3.2	0.44	0.00925926	0	0	0.0666667	0.0166667	0.05	0
gcc-4.4.3	0.44	0.00925926	0	0	0.0666667	0.0166667	0.05	0
gcc-4.6.3	0.44	0.00925926	0	0.2	0.0666667	0.0166667	0.05	0
gcc-4.7.0	0.44	0.00925926	0	0.2	0.0666667	0.0166667	0.05	0
intel-10.1.021	0.330275	0.00903614	0	0.0714286	0.047619	0.0254777	0.05	0
intel-11.0.074	0.330275	0.00903614	0	0.0714286	0.047619	0.0254777	0.05	0
intel-11.1	0.926606	1	1	1	1	1	1	1
intel-12.0.2	0.311927	0.00903614	0	0.0714286	0.047619	0.0254777	0.05	0
intel-14.1.0	0.4	0.00925926	0	0.2	0.133333	0.0166667	0.05	0
open64-4.2.3	0.330275	0.00903614	0	0	0.047619	0.0254777	0.05	0
pgi-7.2.5	0.311927	0.00903614	0	0.0714286	0.047619	0.0254777	0.05	0
pgi-8.0.1	0.311927	0.00903614	0	0.0714286	0.047619	0.0254777	0.05	0
pgi-10.3	0.311927	0.00903614	0	0.0714286	0.047619	0.0254777	0.05	0
pgi-11.8	0.311927	0.00903614	0	0.0714286	0.047619	0.0254777	0.05	0
pgi-12.8	0.926606	2.23795	1	1	1	1	1	1
pgi-13.1	0.926606	2.23494	1	1	1	1	1	1
pgi-14.1-0	0.321101	1.69277	0	0.0714286	0.0714285	0.0254777	0.05	0
sun-12.1	0.311927	0.00903614	0	0	0.047619	0.0254777	0.05	0
sun-12.1+bcheck	0.247706	0.0150602	0	0	1.16667	0.0191083	0	0



SuperComputing Applications and Innovation



Fortran

gcc	-frange-check -O0 -fbounds-check -g -ffpe-trap=invalid,zero,overflow -fdiagnostics-show-location=every-line
g95	-O0 -fbounds-check -g -ftrace=full
intel	-O0 -C -g -traceback -ftrapuv -check
open64	-C -g -O0
pgi	-C -g -Mchkptr -O0
sun	-g -C -O0 -xs -ftrap=%all -fnonstd -xcheck=%all

С

-	
gcc	-O0 -g -fbounds-check -ftrapv
intel	-O0 -C -g -traceback
open64	-g -C -O0
pgi	-g -C -Mchkptr -O0
sun	-g -C -O0 -xs -ftrap=%all -fnonstd -xcheck=%all

C++

gcc	-O0 -g -fbounds-check -ftrapv
intel	-O0 -C -g -traceback
open64	-g -C -O0
pgi	-g -C -Mchkptr -O0
sun	-g -C -O0 -xs -ftrap=%all -fnonstd -xcheck=%all









Bugs and Prevention

Testing

Static analysis

Run-time analysis Memory checker





Fixing memory problems



- Memory leaks are data structures that are allocated at runtime, but not deallocated once they are no longer needed in the program.
- Incorrect use of the memory management is associated with incorrect calls to the memory management: freeing a block of memory more than once, accessing memory after freeing...
- Buffer overruns are bugs where memory outside of the allocated boundaries is overwritten, or corrupted.
- Uninitialized memory bugs: reading uninitialized memory.





Valgrind



- Open Source Software, available on Linux for x86 and PowerPc processors.
- Interprets the object code, not needed to modify object files or executable, non require special compiler flags, recompiling, or relinking the program.
- Command is simply added at the shell command line.
- ► No program source is required (black-box analysis).

www.valgrind.org





Valgrind:tools



- Memcheck: a memory checker.
- ► Callgrind: a runtime profiler.
- ► Cachegrind: a cache profiler.
- Helgrind: find race conditions.
- Massif: a memory profiler.





Why should I use Valgrind?



- Valgrind will tell you about tough to find bugs.
- Valgrind is very through.
- You may be tempted to think that Valgrind is too picky, since your program may seem to work even when valgrind complains. It is users' experience that fixing ALL Valgrind complaints will save you time in the long run.





Why should I use Valgrind?



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But...

Valgrind is kind-of like a virtual x86 interpeter. So your program will run 10 to 30 times slower than normal. Valgrind won't check static arrays.



Use of uninitialized memory:test1.c



- Local Variables that have not been initialized.
- ► The contents of malloc's blocks, before writing there.





Use of uninitialized memory:test1.c



- Local Variables that have not been initialized.
- ► The contents of malloc's blocks, before writing there.

```
#include <stdlib.h>
1
   int main()
2
3
         int p,t,b[10];
4
         if (p==5)
5
              t=p+1;
6
7
              b[p]=100;
         return 0;
8
9
   }
```





Use of uninitialized memory:test1.c



- Local Variables that have not been initialized.
- ► The contents of malloc's blocks, before writing there.

1	<pre>#include <stdlib.h></stdlib.h></pre>
2	<pre>int main()</pre>
3	{
4	<pre>int p,t,b[10];</pre>
5	if (p==5) ERROR
6	t=p+1;
7	b[p]=100; ERROR
8	return 0;
9	}





Valgrind output







Valgrind output



```
==7879== Memcheck, a memory error detector.
==7879== Conditional jump or move depends on uninitialised value(s)
==7879==
            at 0x8048399: main (test1.c:5)
==7879==
==7879== Use of uninitialised value of size 4
==7879==
            at 0x80483A7: main (test1.c:7)
==7879==
==7879== Invalid write of size 4
==7879== at 0x80483A7: main (test1.c:7)
==7879== Address 0xCEF8FE44 is not stack'd, malloc'd or (recently) free'd
==7879==
==7879== Process terminating with default action of signal 11 (SIGSEGV)
==7879== Access not within mapped region at address 0xCEF8FE44
==7879==
           at 0x80483A7: main (test1.c:7)
==7879==
==7879== ERROR SUMMARY: 3 errors from 3 contexts (suppressed: 3 from 1)
==7879== malloc/free: in use at exit: 0 bytes in 0 blocks.
==7879== malloc/free: 0 allocs, 0 frees, 0 bytes allocated.
==7879== For counts of detected errors, rerun with: -v
==7879== All heap blocks were freed -- no leaks are possible.
Segmentation fault
```





Illegal read/write test2.c



```
#include <stdlib.h>
1
2
    int main()
3
    ł
         int *p,i,a;
4
         p=malloc(10*sizeof(int));
5
              p[11]=1;
6
              a=p[11];
7
8
              free(p);
         return 0;
9
10
    }
```





Illegal read/write test2.c



```
#include <stdlib.h>
1
2
    int main()
    ł
3
         int *p,i,a;
4
         p=malloc(10*sizeof(int));
5
             p[11]=1; ERROR
6
              a=p[11]; ERROR
7
              free(p);
8
         return 0;
9
10
    }
```





Illegal read/write: output



```
==8081== Invalid write of size 4
==8081==
            at 0x804840A: main (test2.c:6)
==8081== Address 0x417B054 is 4 bytes after a block of size 40 alloc'd
==8081==
            at 0x40235B5: malloc (in /usr/lib/valgrind/x86-linux/vgpreload_memcheck.so)
==8081==
            by 0x8048400: main (test2.c:5)
==8081==
==8081== Invalid read of size 4
==8081==
            at 0x8048416: main (test2.c:7)
==8081== Address 0x417B054 is 4 bytes after a block of size 40 alloc'd
==8081==
            at 0x40235B5: malloc (in /usr/lib/valgrind/x86-linux/vgpreload memcheck.so)
==8081==
            by 0x8048400: main (test2.c:5)
==8081==
==8081== ERROR SUMMARY: 2 errors from 2 contexts (suppressed: 3 from 1)
==8081== malloc/free: in use at exit: 0 bytes in 0 blocks.
==8081== malloc/free: 1 allocs, 1 frees, 40 bytes allocated.
==8081== For counts of detected errors, rerun with: -v
==8081== All heap blocks were freed -- no leaks are possible.
```





Invalid free:test3.c



1	<pre>#include <stdlib.h></stdlib.h></pre>
2	int main()
3	{
4	int *p,i;
5	<pre>p=malloc(10*sizeof(int));</pre>
6	for (i=0;i<10;i++)
7	p[i]=i;
8	free(p);
9	free(p);
10	return 0;
11	}





Invalid free:test3.c



1	<pre>#include <stdlib.h></stdlib.h></pre>
2	int main()
3	{
4	int *p,i;
5	<pre>p=malloc(10*sizeof(int));</pre>
6	for(i=0;i<10;i++)
7	<pre>p[i]=i;</pre>
8	free(p);
9	free(p); ERROR
10	return 0;
11	}





Invalid free: output



```
==8208== Invalid free() / delete / delete[]
==8208==
           at 0x40231CF: free (in /usr/lib/valgrind/x86-linux/vgpreload_memcheck.so)
==8208==
           by 0x804843C: main (test3.c:9)
==8208== Address 0x417B028 is 0 bytes inside a block of size 40 free'd
==8208==
           at 0x40231CF: free (in /usr/lib/valgrind/x86-linux/vgpreload_memcheck.so)
==8208==
           by 0x8048431: main (test3.c:8)
==8208==
==8208== ERROR SUMMARY: 1 errors from 1 contexts (suppressed: 3 from 1)
==8208== malloc/free: in use at exit: 0 bytes in 0 blocks.
==8208== malloc/free: 1 allocs, 2 frees, 40 bytes allocated.
==8208== For counts of detected errors, rerun with: -v
==8208== All heap blocks were freed -- no leaks are possible.
```







- If allocated with malloc,calloc,realloc,valloc or memalign, you must deallocate with free.
- If allocated with new[], you must dealloacate with delete[].
- ► If allocated with **new**, you must deallocate with **delete**.







- If allocated with malloc,calloc,realloc,valloc or memalign, you must deallocate with free.
- If allocated with new[], you must dealloacate with delete[].
- ► If allocated with **new**, you must deallocate with **delete**.

```
#include <stdlib.h>
1
    int main()
2
3
          int *p,i;
4
          p=(int*)malloc(10*sizeof(int));
5
               for (i=0; i<10; i++)</pre>
6
                  p[i]=i;
7
               delete(p);
8
          return 0;
9
10
```







- If allocated with malloc,calloc,realloc,valloc or memalign, you must deallocate with free.
- If allocated with new[], you must dealloacate with delete[].
- ► If allocated with **new**, you must deallocate with **delete**.

```
#include <stdlib.h>
1
    int main()
2
3
          int *p,i;
4
          p=(int*)malloc(10*sizeof(int));
5
               for (i=0; i<10; i++)</pre>
6
7
                  p[i]=i;
              delete(p);
                                ERROR
8
          return 0;
9
10
```







```
....
=8330== Mismatched free() / delete / delete []
==8330== at 0x4022EE6: operator delete(void*) (in /usr/lib/valgrind/x86-linux/vgpreload_memch
=8330==address 0x4292028 is 0 bytes inside a block of size 40 alloc'd
=8330==at 0x4025B5: malloc (in /usr/lib/valgrind/x86-linux/vgpreload_memcheck.so)
=8330==by 0x80484C0: main (test4.c:5)
=8330==by 0x80484C0: main (test4.c:5)
=8330==malloc/free: in use at exit: 0 bytes in 0 blocks.
=8330==malloc/free: 1 allocs, 1 frees, 40 bytes allocated.
=8330==For counts of detected errors, rerun with: -v
=8330==All heap blocks were freed -- no leaks are possible.
```







```
#include <stdlib.h>
1
   #include <unistd.h>
2
    int main()
3
4
         int *p;
5
         p=malloc(10);
6
          read(0,p,100);
7
              free(p);
8
         return 0;
9
10
    }
```







```
#include <stdlib.h>
1
   #include <unistd.h>
2
    int main()
3
4
         int *p;
5
         p=malloc(10);
6
          read(0,p,100); ERROR
7
              free(p);
8
         return 0;
9
10
    }
```







```
...
=-18007== Syscall param read(buf) points to unaddressable byte(s)
=-18007== at 0x4EEC240: __read_nocancel (in /lib64/libc-2.5.so)
=-18007== by 0x40056F: main (test5.c:7)
=-18007== Address 0x517d04a is 0 bytes after a block of size 10 alloc'd
=-18007== at 0x4C21166: malloc (vg_replace_malloc.c:236)
=-18007== by 0x400555: main (test5.c:6)
```














	and the second
1	#Include <stdlib.h></stdlib.h>
2	int main()
3	{
4	int *p,i;
5	<pre>p=malloc(5*sizeof(int));</pre>
6	for(i=0; i<5;i++)
7	p[i]=i;
8	free (p) ;
9	return 0;
10	}





Memory leak detection: output



ruggiero@shiva:> valgrind --tool=memcheck --leak-check=full ./t6

```
==8237== ERROR SUMMARY: 0 errors from 0 contexts (suppressed: 3 from 1)
   ==8237== malloc/free: in use at exit: 20 bytes in 1 blocks.
   ==8237== malloc/free: 1 allocs, 0 frees, 20 bytes allocated.
   ==8237== For counts of detected errors, rerun with: -v
   ==8237== searching for pointers to 1 not-freed blocks.
   ==8237== checked 65,900 bytes.
    ==8237==
   ==8237== 20 bytes in 1 blocks are definitely lost in loss record 1 of 1
   ==8237== at 0x40235B5: malloc (in /usr/lib/valgrind/x86-linux/vgpreload memcheck.so)
           by 0x80483D0: main (test6.c:5)
==8237==
   ==8237==
   ==8237== LEAK SUMMARY:
   ==8237== definitely lost: 20 bytes in 1 blocks.
   ==8237==
                 possibly lost: 0 bytes in 0 blocks.
   ==8237== still reachable: 0 bytes in 0 blocks.
   ==8237==
                     suppressed: 0 bytes in 0 blocks.
```





What won't Valgrind find?



1	<pre>int main()</pre>
2	{
3	char x[10];
4	x[11]='a';
5	}





What won't Valgrind find?



1	int main()
2	{
3	char x[10];
4	x[11]='a';
5	}

- Valgrind doesn't perform bound checking on static arrays (allocated on stack).
- Solution for testing purposes is simply to change static arrays into dinamically allocated memory taken from the heap, where you will get bounds-checking, though this could be a message of unfreed memory.





sum.c: source



```
#include <stdio.h>
1
    #include <stdlib.h>
2
    int main (int argc, char* argv[]) {
3
4
        const int size=10;
        int n, sum=0;
5
6
        int* A = (int*)malloc( sizeof(int)*size);
7
8
        for (n=size; n>0; n--)
9
            A[n] = n;
        for(n=0; n<size; n++)</pre>
10
            sum+=A[n];
11
        printf("sum=%d\n", sum);
12
13
        return 0;
14
    1
```







ruggiero@shiva:~> gcc -00 -g -fbounds-check -ftrapv sum.c

ruggiero@shiva:~> ./a.out

sum=45





Valgrind:example



ruggiero@shiva:~> valgrind --leak-check=full --tool=memcheck ./a.out

```
==21579== Memcheck, a memory error detector.
==21791==Invalid write of size 4
==21791==at 0x804842A: main (sum.c:9)
==21791==Address 0x417B050 is 0 bytes after a block of size 40 alloc'd
==21791==at 0x40235B5: malloc (in /usr/lib/valgrind/x86-linux/vgpreload_memcheck.so)
==21791==by 0x8048410: main (sum.c:6)
==21791==Use of uninitialised value of size 4
==21791== at 0x408685B; itoa word (in /lib/libc-2.5.so)
==21791==bv 0x408A581: vfprintf (in /lib/libc-2.5.so)
==21791==by 0x4090572: printf (in /lib/libc-2.5.so)
==21791==bv 0x804846B; main (sum.c:12)
==21791==
==21791==Conditional jump or move depends on uninitialised value(s)
==21791==at 0x4086863: itoa word (in /lib/libc-2.5.so)
==21791==by 0x408A581: vfprintf (in /lib/libc-2.5.so)
==21791==by 0x4090572: printf (in /lib/libc-2.5.so)
==21791==by 0x804846B: main (sum.c:12)
==21791==40 bytes in 1 blocks are definitely lost in loss record 1 of 1
==21791==at 0x40235B5: malloc (in /usr/lib/valgrind/x86-linux/vgpreload memcheck.so)
==21791==by 0x8048410: main (sum.c:6)
==21791==
```









1	<pre>#include <stdio.h></stdio.h></pre>
2	<pre>#include <stdlib.h></stdlib.h></pre>
3	int main (void)
4	{
5	int i;
6	<pre>int *a = (int*) malloc(9*sizeof(int));</pre>
7	
8	for (i=0; i<=9; ++i){
9	a[i] = i;
10	<pre>printf ("%d\n ", a[i]);</pre>
11	}
12	
13	free (a) ;
14	return 0;
15	}







ruggiero@shiva:~> icc -C -g outbc.c

ruggiero@shiva:~> ./a.out









ruggiero@shiva:~> pgcc -C -g outbc.c

ruggiero@shiva:~> ./a.out







Electric Fence



- Electric Fence (efence) stops your program on the exact instruction that overruns (or underruns) a malloc() memory buffer.
- GDB will then display the source-code line that causes the bug.
- It works by using the virtual-memory hardware to create a red-zone at the border of each buffer - touch that, and your program stops.
- Catch all of those formerly impossible-to-catch overrun bugs that have been bothering you for years.





Electric Fence



ruggiero@shiva:~> icc -g outbc.c libefence.a -o outbc -lpthread

ruggiero@shiva:~> ./outbc



SCA Esercitazione:Run time errors



- http://rted.public.iastate.edu/Serial/tests.html
- Nella directory RTED leggere il file Readme
- Modificare il compilatore e le relative opzioni per cercare di ottenere il valore piu' alto possibile
- Posso migliorare i risultati usando valgrind?

