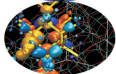
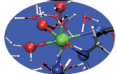
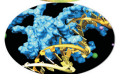
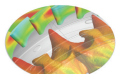


# Introduction to Fortran 90

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CINECA Bologna - SCAI Department

Bologna, 21-24 October 2013





## Part I

# A Fortran Survey 1

Program main unit, source formats, comments, declarations and instructions. Fundamental operators, expressions, conditional constructs, loops, functions: arguments passing, intent, interface, intrinsic and external functions. Modules: contains and use. Intrinsic types: integer, real, complex, logical, and parameter. I/O base.



Introduction

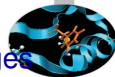
Fortran Basics

More Fortran Basics

Integer Types and Iterating

More on Compiling and Linking

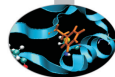
Homeworks



- ▶ Developed in the 50s among the earliest high level languages (HLL)
- ▶ Widely and rapidly adopted in the area of numerical, scientific, engineering and technical applications
- ▶ First standard in 1966: Fortran 66
  - ▶ The first of all programming language standards
- ▶ Second standard in 1978: Fortran 77
- ▶ Third standard in 1991: Fortran 90
  - ▶ Adds new, modern features such as structured constructs, array syntax and ADT
  - ▶ Extended and revised in 1997: Fortran 95
  - ▶ Further extended with published Technical Reports
- ▶ Fourth standard in 2004: Fortran 2003
  - ▶ Major revision, incorporates TRs, adds many new features (OO!), still not fully supported
- ▶ Fifth standard in 2010: Fortran 2008

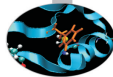


- ▶ Strongly oriented to *number crunching*
- ▶ Efficient language, highly optimized code
  - ▶ Basic data types and operators mapping "naturally" to CPUs
  - ▶ Translated by a compiler to machine language
  - ▶ Language rules allow for aggressive, automatic optimization
  - ▶ Facilities to build new data types from the basic ones
  - ▶ Flexible flow control structures mapping the most common numerical computing use cases
- ▶ Scientific computing specialized syntax
  - ▶ A wealth of math data types and functions available as intrinsics of the language
  - ▶ Compact, readable array syntax to operate on many values as a whole



► Why Fortran is bad

- Current standard embodies four different language versions,...
- ... all of them still alive in legacy codes
- Non-numeric computing in Fortran is a real pain
- There are more C than Fortran programmers
- GUI and DB accesses are best programmed in C
- C99 partly addressed numerical computing needs



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## ▶ Why Fortran is good

- ▶ Fortran is highly tuned for numerical computation
- ▶ Fortran is older and more “rigid” than C, compilers optimize better
- ▶ Much better than C at managing user defined data types
- ▶ Object-oriented features are now part of the language
- ▶ Provides facilities for interoperability with C and other languages



- ▶ Teach you the fundamentals of modern Fortran
- ▶ For both reading (old and new) and writing (new) programs
- ▶ Showing common idioms
- ▶ Illustrating and demonstrating many of the extensions introduced in the more recent standards
- ▶ Illustrating best practices
- ▶ Blaming bad ones
- ▶ Making you aware of the typical traps





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- ▶ **Neither a substitute for personal practice**



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Compiling and Linking Your First Program

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More Types and Choices

Wrapping it Up 1

More Fortran Basics

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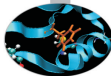
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! roots of a 2nd degree equation with real coefficients

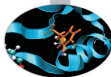
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  delta = sqrt(b**2 - 4.0*a*c) ! square root of discriminant
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**! roots of a 2nd degree equation with real coefficients**

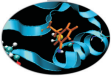
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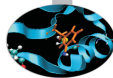
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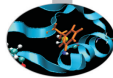


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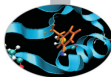


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  - ▶ Variable contents
  - ▶ Algorithms
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  - ▶ Tricks





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- ▶ Best practice: do comment your code!
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  - ▶ Tricks
- ▶ Best practice: do not over-comment your code!
  - ▶ Obvious comments obfuscate code and annoy readers
  - ▶ ! **square root of discriminant** is a bad example



! roots of a 2nd degree equation with real coefficients

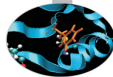
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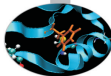
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  - ▶ Procedures (subroutines and functions)
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- ▶ Best practice: always mark unit beginning and ending with its type and name
  - ▶ Makes your readers (including you) happier



! roots of a 2nd degree equation with real coefficients

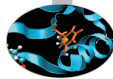
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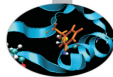


- ▶ **real :: x1, x2** declares two variables
  - ▶ Named memory locations where values can be stored
  - ▶ Declared by specifying a data type, an optional attribute list, and a comma-separated list of names
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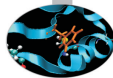


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- ▶ A legal *name* must be used for a variable:
  - ▶ Permitted characters: **a-z, A-Z, 0-9, \_**
  - ▶ The first one cannot be a digit (e.g. **x1** is a valid name, **1x** is not)
  - ▶ At most 31 characters are permitted (63 in Fortran 2003)
  - ▶ A good advice: do not exceed 31 characters in a name





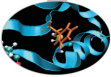
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- ▶ Beware: Fortran is CaSe insenSITIVE!



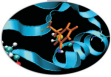
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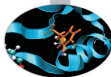
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- ▶ Best practice: it is strongly recommended to turn off implicit declarations with **`implicit none`**, at the beginning of each program unit
  - ▶ Improves readability and clarity: each variable has its type declared
  - ▶ Mistyped names can be caught by the compiler as undeclared variables



! roots of a 2nd degree equation with real coefficients

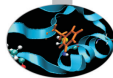
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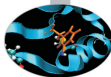
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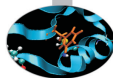
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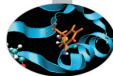
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- ▶ `x1 = x1 + delta` is a statement assigning the value of expression `x1 + delta` to variable `x1`
- ▶ By the way, expressions can be passed as argument to functions, as to `sqrt ()`: their value will be computed and passed to the function



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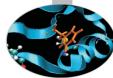


- ▶ Fortran lets you write programs in a high-level, human-readable language
- ▶ Computer CPUs do not directly understand this language
- ▶ You need to translate your code into machine-level instructions for your CPU architecture
- ▶ **Compilers take care of that translation and generate machine code that can be actually executed by a CPU**



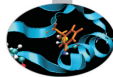
- ▶ Compilers are sophisticated tools, made up of many components
- ▶ When compiler is invoked to generate executable code, three main steps are performed:
  1. parsing of source files, various kinds of analysis and transformations, optimization and *assembly* files creation
  2. machine-code generation and object file creation
    - ▶ an object file is an organized collection of all symbols (variables, functions...) used or referenced in the code
  3. linking and executable creation
- ▶ Options are provided to execute each step separately, take a look at the manual of your favourite compiler, there's a lot to learn!

# Compile your first Fortran program !



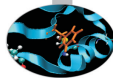
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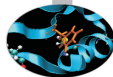
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- ▶ Many more available on the market (Intel, PGI, Pathscale, IBM XL Fortran, Sun Studio Fortran, g95, Lahey, NAG, etc)
- ▶ Let's use **gfortran** to compile and run our examples and exercises

# Compile your first Fortran program !

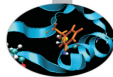


- ▶ GNU compiler collection includes **gfortran** compiler, supporting Fortran 95 and some features of the 2003 standard
- ▶ Many more available on the market (Intel, PGI, Pathscale, IBM XL Fortran, Sun Studio Fortran, g95, Lahey, NAG, etc)
- ▶ Let's use **gfortran** to compile and run our examples and exercises
  - ▶ Compile with:

```
user@cinca$> gfortran second_degree_eq.f90
```

An executable file named **a.out** (**a.exe** under Windows) will be generated





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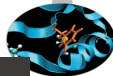
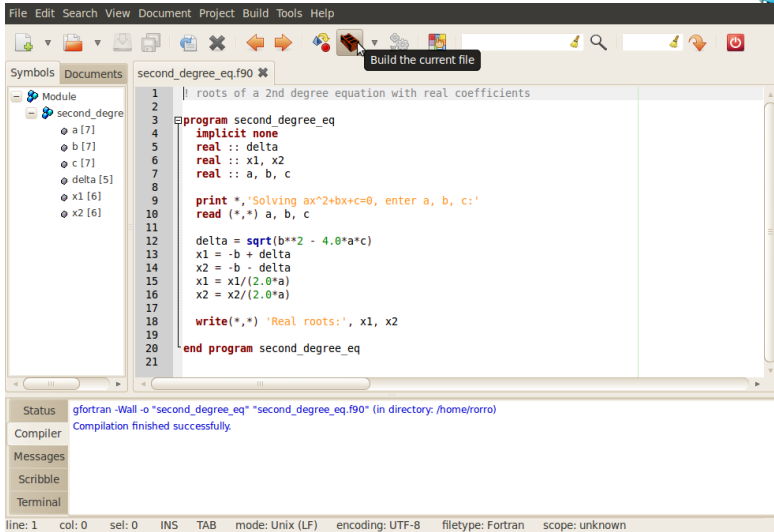
An executable file named **a.out** (**a.exe** under Windows) will be generated

- ▶ Run the program under GNU/Linux with:

```
user@cineca$> ./a.out
```

or under Windows with:

```
C:\Documents and Settings\user> a.exe
```

File Edit Search View Document Project Build Tools Help

Build the current file

second\_degree\_eq.f90

```

1  | roots of a 2nd degree equation with real coefficients
2
3  program second_degree_eq
4  implicit none
5  real :: delta
6  real :: x1, x2
7  real :: a, b, c
8
9  print *, 'Solving ax^2+bx+c=0, enter a, b, c:'
10 read (*,*) a, b, c
11
12 delta = sqrt(b**2 - 4.0*a*c)
13 x1 = -b + delta
14 x2 = -b - delta
15 x1 = x1/(2.0*a)
16 x2 = x2/(2.0*a)
17
18 write(*,*) 'Real roots:', x1, x2
19
20 end program second_degree_eq
21

```

Status: gfortran -Wall -o "second\_degree\_eq" "second\_degree\_eq.f90" (in directory: /home/orroro)  
 Compiler: Compilation finished successfully.  
 Messages:  
 Scribble  
 Terminal

line: 1 col: 0 sel: 0 INS TAB mode: Unix (LF) encoding: UTF-8 filetype: Fortran scope: unknown



! roots of a 2nd degree equation with real coefficients

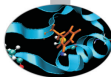
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Introduction

**Fortran Basics**

My First Fortran Program

Compiling and Linking Your First Program

**Making Choices**

More Types and Choices

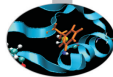
Wrapping it Up 1

More Fortran Basics

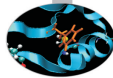
Integer Types and Iterating

More on Compiling and Linking

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if (delta < 0.0) then  
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end if  
delta = sqrt(delta)
```

- ▶ Try it now!



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- ▶ Try it now!
- ▶ Did you check that normal cases still work? Good.



▶ **if** (*logical-condition*) **then**  
*block of statements*  
**end if**

- ▶ Executes *block of statements* only if *logical-condition* is true
- ▶ Comparison operators: == (equal), /= (not equal), >, <, >=, <=
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if (delta < 0.0) then  
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to:

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if (delta < 0.0) stop 'No real roots!'
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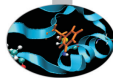
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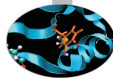
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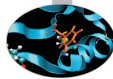
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but this is not that good...

- ▶ In general, Fortran disregards white space, but proper indentation visualizes program control flow



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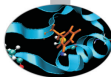
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! roots of a 2nd degree equation with real coefficients

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program second_degree_eq
  implicit none
  real :: delta
  real :: rp
  real :: a, b, c

  print *, 'Solving ax^2+bx+c=0, enter a, b, c: '
  read(*,*) a, b, c

  delta = b*b - 4.0*a*c
  if (delta < 0.0) stop 'No real roots!'
  delta = sqrt(delta)/(2.0*a)

  rp = -b/(2.0*a)

  print *, 'Real roots: ', rp+delta, rp-delta

end program second_degree_eq
  
```

# And Now Make it More Complex!



! roots of a 2nd degree equation with real coefficients

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program second_degree_eq
  implicit none
  real :: delta, rp, a, b, c
  logical :: rroots

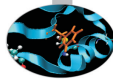
  print *, 'Solving ax^2+bx+c=0, enter a, b, c: '
  read(*,*) a, b, c
  delta = b*b - 4.0*a*c
  rroots = .true.
  if (delta < 0.0) then
    delta = -delta
    rroots = .false.
  end if
  delta = sqrt(delta)/(2.0*a)
  rp = -b/(2.0*a)
  if (rroots) then
    print *, 'Real roots: ', rp+delta, rp-delta
  else
    print *, 'Complex roots: ', rp, '+', delta, 'i ', &
      rp, '-', delta, 'i'
  end if
end program second_degree_eq
  
```



- ▶ **logical** type represents logical values
  - ▶ Can be `.true.` or `.false.`

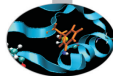


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! roots of a 2nd degree equation with real coefficients

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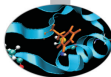
program second_degree_eq
  implicit none
  real :: delta, rp, a, b, c
  logical :: roots

  print *, 'Solving ax^2+bx+c=0, enter a, b, c: '
  read(*,*) a, b, c
  delta = b*b - 4.0*a*c
  roots = .true.
  if (delta < 0.0) then
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  end if
  delta = sqrt(delta)/(2.0*a)
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```





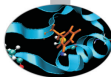
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- ▶ Again, use proper indentation
- ▶ And Fortran statements cannot exceed one line, unless it ends with an `&`



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  implicit none
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  end if
end program second_degree_eq
```



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program second_degree_eq
  implicit none
  complex :: delta
  complex :: z1, z2
  real    :: a, b, c

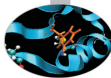
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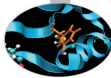
  z1 = (-b+delta)/(2.0*a)
  z2 = (-b-delta)/(2.0*a)

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end program second_degree_eq
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- ▶ Most math functions like **sqrt ()** work for complex type too!
  - ▶ Returning correct results, instead of NaNs
- ▶ And so do **read**, **write**, and **print**
- ▶ **(1.5, 2.3)** is *Fortranese* for  $1.5 + 2.3i$



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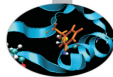
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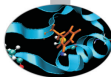
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    if (c == 0.0) then
      write(0,*) 'A trivial identity!'
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    end if
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stop
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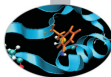
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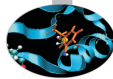
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  - ▶ And **stop error-message** is equivalent to: **write(0,\*) error-message**  
**stop**
- ▶ Best practice: if your program has to fail, always have it fail in a controlled way





- ▶ Let's give names to `if` constructs:

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no2nd: if (a == 0.0) then
  nolst: if (b == 0.0) then
    no0th: if (c == 0.0) then
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      write(0,*) 'Plainly absurd!'
    end if no0th
  else nolst
    write(0,*) 'Too simple problem!'
  end if nolst
  stop
end if no2nd
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  stop
end if no2nd
```

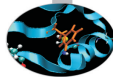
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    write(0,*) 'Too simple problem!'
  end if nolst
  stop
end if no2nd
```

- ▶ Giving names to constructs makes program logic more explicit
- ▶ Names are for readability purposes only, do not enforce pairing rules



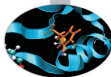
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  nolst: if (b == 0.0) then
    no0th: if (c == 0.0) then
      write(0,*) 'A trivial identity!'
    else no0th
      write(0,*) 'Plainly absurd!'
    end if no0th
  else nolst
    write(0,*) 'Too simple problem!'
  end if nolst
  stop
end if no2nd
```

- ▶ Giving names to constructs makes program logic more explicit
- ▶ Names are for readability purposes only, do not enforce pairing rules
- ▶ Best practice: always give names to constructs which span many lines of code or are deeply nested



- ▶ The one on the right, is the statement  $Z(I) = Y + W(I)$
- ▶ The one in the middle, is an IBM punch card reader
- ▶ The one on the left, is a complete Fortran source program
- ▶ But you'll only encounter these in museums, nowadays



## C ROOTS OF A 2ND DEGREE EQUATION WITH REAL COEFFICIENTS

```
PROGRAM EQ2DEG
```

```
IMPLICIT NONE
```

```
REAL DELTA
```

```
REAL RP
```

```
REAL A, B, C
```

```
PRINT *, 'SOLVING AX^2+BX+C=0, ENTER A, B, C: '
```

```
READ(*,*) A, B, C
```

```
DELTA = B*B - 4.0*A*C
```

```
IF (DELTA.LT.0.0) STOP 'NO REAL ROOTS!'
```

```
DELTA = SQRT(DELTA) / (2.0*A)
```

```
RP = -B / (2.0*A)
```

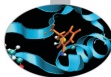
```
PRINT *, 'REAL ROOTS: ', RP+DELTA, RP-DELTA
```

```
END
```



- ▶ Code is all capitals
  - ▶ First computers had only uppercase letters
- ▶ Fixed source form
  - ▶ The legacy of punch cards
  - ▶ Comment lines must be marked with a C or \* in first column
  - ▶ First six columns on each line are reserved for labels and to mark continuation lines
  - ▶ Columns after the 72nd are ignored (cause of really nasty bugs!)
- ▶ No double colon on variable declarations
  - ▶ And no way to initialize a variable at declaration, for that matter
  - ▶ More on this later
- ▶ And this example is not that different...

# A Bottle of Fortran, Vintage Year 1963



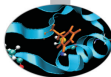
```

C      SOLUTION OF QUADRATIC EQUATION
C      (P. 122 OF A FORTRAN PRIMER BY E. ORGANICK)
1 READ INPUT TAPE 5, 51, ANAME, N
51 FORMAT(A6,I2)
   WRITE OUTPUT TAPE 6,52, ANAME
52 FORMAT(1H1,33HROOTS OF QUADRATIC EQUATIONS FROM A6)
   DO 21 I = 1, N
   READ INPUT TAPE 5, 53, A, B, C
53 FORMAT(3F10.2)
   WRITE OUTPUT TAPE 6,54, I, A, B, C
54 FORMAT(1H0,8HSET NO. I2/5H A = F8.2,12X,4HB = F8.2,12X,4HC = F8.2)
   IF(A) 10, 7, 10
   7 RLIN = -C/B
   WRITE OUTPUT TAPE 6, 55, RLIN
55 FORMAT(7H LINEAR,25X,4HX = F10.3)
   GO TO 21
10 D = B**2 - 4.*A*C
   IF(D) 12, 17, 17
12 COMPR = -B/(2.*A)
   COMP1 = SQRTF(-D)/(2.*A)
   COMP2= -COMP1
   WRITE OUTPUT TAPE 6, 56, COMPR, COMP1, COMP2
56 FORMAT(8H COMPLEX,21X,7HR(X1)= F10.3,11X,7HI(X1)= F10.3,/1H ,28X,
17HR(X2)= F10.3,11X,7HI(X2)= F10.3)
16 GO TO 21
17 REAL1 = (-B + SQRTF(D))/(2.*A)
   REAL2 = (-B - SQRTF(D))/(2.*A)
20 WRITE OUTPUT TAPE 6, 57, REAL1, REAL2
57 FORMAT(6H REAL 25X,5HX1 = F10.3,13X,5HX2 = F10.3)
21 CONTINUE
   WRITE OUTPUT TAPE 6, 58, ANAME
58 FORMAT(8HOEND OF A6)
   GO TO 1
   END
    
```





- ▶ Write new code in free source form
  - ▶ No limits on beginning of program statements
  - ▶ Each line may contain up to 132 default characters
  - ▶ Comments can be added at end of line
  - ▶ And it comes for free: just give your source file name an `.f90` extension
- ▶ Use new language features
  - ▶ Like new styles for declarations
  - ▶ Or naming of constructs
  - ▶ They are more powerful and readable
- ▶ We'll focus on modern Fortran programming style
  - ▶ Making you aware of differences you are most likely to encounter
  - ▶ Look at compiler manuals or reference books to tame very old codes



Introduction

**Fortran Basics**

My First Fortran Program

Compiling and Linking Your First Program

Making Choices

More Types and Choices

**Wrapping it Up 1**

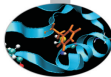
More Fortran Basics

Integer Types and Iterating

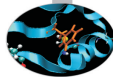
More on Compiling and Linking

Homeworks

# A Fortran Program is Made of: I



- ▶ **Comments**
  - ▶ Compiler disregards them, but humans do not
  - ▶ Please, use them
  - ▶ Do not abuse them, please
- ▶ **Program units**
  - ▶ One, at least: **program**
  - ▶ Some of them (functions) are intrinsic to the language
- ▶ **Variables**
  - ▶ Named memory location you can store values into
  - ▶ Must be declared
- ▶ **Variables declarations**
  - ▶ Give name to memory location you can store values into
  - ▶ An initial value can be specified



- ▶ Expressions
  - ▶ Compute values to store in variables
  - ▶ Compute values to pass to functions and statements
- ▶ Statements
  - ▶ Units of executable work
  - ▶ Whose execution can be controlled by other constructs
- ▶ **if** statements and constructs
  - ▶ Allow for conditional and alternative execution
  - ▶ For both single statements and blocks of



- ▶ Use free source form
- ▶ **implicit none** statement
  - ▶ Turn off implicit declarations
- ▶ Use proper indentation
  - ▶ Compilers don't care about
  - ▶ Readers visualize flow control
- ▶ Give names to complex control structures, readers will appreciate
- ▶ Do non-regression testing
  - ▶ Whenever functionalities are added
  - ▶ Whenever you rewrite a code in a different way
- ▶ Fail in a controlled way
  - ▶ Giving feedback to humans



Introduction

Fortran Basics

More Fortran Basics

My First Fortran Functions

Making it Correct

Making it Robust

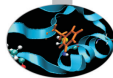
Copying with Legacy

Wrapping it Up 2

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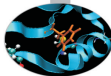
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Homeworks



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function theta(x) !Heaviside function, useful in DSP
  implicit none
  real :: theta
  real :: x

  theta = 1.0
  if (x < 0.0 ) theta = 0.0
end function theta
```

```
function sinc(x) !sinc function as used in DSP
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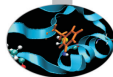
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# My First Fortran Functions



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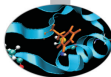
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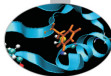
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  - ▶ Function name must be a legal Fortran name
  - ▶ Functions specialty is performing computations and returning a value



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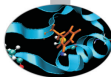
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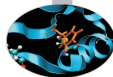


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- ▶ How to return a value
  - ▶ Just assign it to the function name, as if it were a variable
  - ▶ But this doesn't force function termination
  - ▶ Multiple assignments can be done
  - ▶ The last assigned value before function execution is complete will be returned

# My First Fortran Functions



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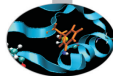


- ▶ Functions have arguments
  - ▶ Declared like variables inside the function
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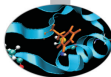


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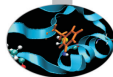
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- ▶ What if a dummy argument has the same name of a variable elsewhere in the program?
  - ▶ No conflicts of sort, they are completely independent
- ▶ Variables can be defined inside functions
  - ▶ Again, they are local, thus completely independent from the rest of the program

# My First Fortran Functions



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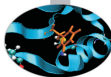
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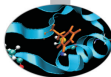
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- ▶ Fortran sports a wealth (over a hundred!) of predefined functions and procedures
- ▶ These are termed *intrinsic*
  - ▶ **acos** (**x**) returns the arc cosine of  $x$  such that  $|x| \leq 1$  in the range  $0 \leq \arccos(x) \leq \pi$
  - ▶ **sin** (**x**) returns the sine function value of  $x$  in radians
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  - ▶ **abs** (**x**) returns the absolute value of  $x$
- ▶ What's this **external** keyword?
- ▶ It's one of the many attributes you can give to something you define
  - ▶ **external** tells the compiler **theta** is an external (i.e. non intrinsic) function
  - ▶ So the compiler is not forced to guess what it is from its use
  - ▶ And that way, masters can override intrinsic functions



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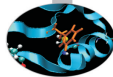
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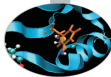
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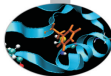


- ▶ The `parameter` attribute is used to declare named constants
  - ▶ I.e. variables that cannot be modified after initialization (compiler will bark if you try)

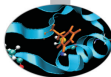




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- ▶ In initialization expressions:
  - ▶ only constants (possibly other `parameters`) can be used
  - ▶ only intrinsic operators or functions are allowed
- ▶ Best practice: always give name to constants
  - ▶ Particularly if unobvious, like `1.0/137.0`
  - ▶ It also helps to centralize updates (well, not for  $\pi$ )



Introduction

Fortran Basics

**More Fortran Basics**

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**Making it Correct**

Making it Robust

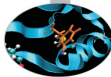
Copying with Legacy

Wrapping it Up 2

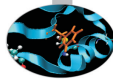
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More on Compiling and Linking

Homeworks



- ▶ Let's put the code in a file named `dsp.f90`
- ▶ Best practice: always put different groups of related functions in different files
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  - ▶ You can always pass all source files to the compiler
  - ▶ And you'll learn to do better ...



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  - ▶ You can always pass all source files to the compiler
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- ▶ And let's write a program to test all functions
  - ▶ And be distrustful, check again actual arguments after all function calls
- ▶ Best practice: always write a special purpose program to test each subset of functions
  - ▶ Best to include in the program automated testing of all relevant cases
  - ▶ Let's do by hand with I/O for now, to make it short



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function theta(x) !Heaviside function, useful in DSP
  implicit none
  real :: theta
  real :: x

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```

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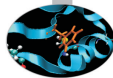
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  real :: abs_t, half_tau
  real, external :: theta

  abs_t = abs(t)
  half_tau = 0.5*tau
  rect = 0.5
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end function rect
```





- ▶ We have collected DSP functions in `dsp.f90` source file

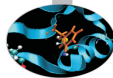


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# DSP test program



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```
program dsp_test
```

```
    real :: i,j,k
```

```
end program dsp_test
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- ▶ Let's write a `dsp_test.f90` program:

```
program dsp_test  
  
  real :: i,j,k  
  real :: rtheta, rsinc, rrect  
  real, external :: theta, sinc, rect
```

```
end program dsp_test
```



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program dsp_test

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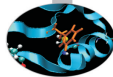
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program dsp_test

  real :: i,j,k
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  print *, 'Enter i, j, k:'
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  rtheta = theta(i)
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end program dsp_test
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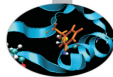
  print *, 'Enter i, j, k:'
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  rtheta = theta(i)
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  write(*,*) 'theta(', i, ')= ', rtheta
  write(*,*) 'sinc(', i, ')= ', rsinc
  write(*,*) 'rect(', j, ', ', k, ')= ', rrect

end program dsp_test
```





- ▶ Let's build our test program putting all together:

```
user@cineca$> gfortran dsp.f90 dsp_test.f90 -o dsp_test
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- ▶ `-o` option specifies the name `dsp_test` for the executable



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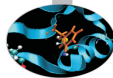
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user@cineca$> ./dsp_test
Enter i, j, k:
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theta( -3.1415927 ) = 0.0000000
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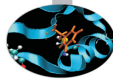
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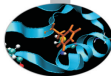
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- ▶ Something is going wrong, isn't it?
  - ▶ Seems like one function changed its actual argument!



- ▶ Let's put the code in a file named `dsp.f90`
- ▶ Best practice: always put different groups of related functions in different files
  - ▶ Helps to tame complexity
  - ▶ You can always pass all source files to the compiler
  - ▶ And you'll learn to do better ...
- ▶ And let's write a program to test all functions
  - ▶ And be distrustful, check again actual arguments after all function calls
- ▶ Best practice: always write a special purpose program to test each subset of functions
  - ▶ Best to include in the program automated testing of all relevant cases
  - ▶ Let's do by hand with I/O for now, to make it short



```
function theta(x) !Heaviside function, useful in DSP
  implicit none
  real :: theta
  real, intent(in) :: x

  theta = 1.0
  if (x < 0.0 ) theta = 0.0
end function theta
```

```
function sinc(x) !sinc function as used in DSP
  implicit none
  real :: sinc
  real, intent(in) :: x
  real, parameter :: pi = acos(-1.0)

  x = x*pi
  sinc = 1.0
  if (x /= 0.0) sinc = sin(x)/x
end function sinc
```

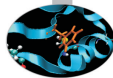
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function rect(t, tau) !generalized rectangular function, useful in DSP
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  real :: rect
  real, intent(in) :: t, tau
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  abs_t = abs(t)
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- ▶ Try to recompile `dsp`. £90...

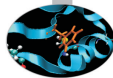




- ▶ Try to recompile `dsp.f90`...
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user@cinca$> gfortran -o dsp_test dsp_test.f90 dsp.f90
dsp.f90:16.2:

    x = x*pi
    1
Error: Cannot assign to INTENT(IN) variable 'x' at (1)
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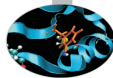
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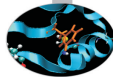
- ▶ Got a compiler error message? Good!

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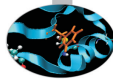
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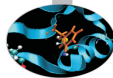
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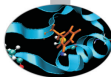


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  - ▶ **intent (inout)** (default) for those you plan to do both



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function theta(x) !Heaviside function, useful in DSP
  implicit none
  real :: theta
  real, intent(in) :: x

  theta = 1.0
  if (x < 0.0 ) theta = 0.0
end function theta
```

```
function sinc(x) !sinc function as used in DSP
  implicit none
  real :: sinc, xpi
  real, intent(in) :: x
  real, parameter :: pi = acos(-1.0)

  xpi = x*pi
  sinc = 1.0
  if (xpi /= 0.0) sinc = sin(xpi)/xpi
end function sinc
```

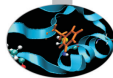
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function rect(t, tau) !generalized rectangular function, useful in DSP
  implicit none
  real :: rect
  real, intent(in) :: t, tau
  real :: abs_t, half_tau
  real, external :: theta

  abs_t = abs(t)
  half_tau = 0.5*tau
  rect = 0.5
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end function rect
```



- ▶ Way much better!

```
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Enter i, j, k:
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- ▶ Now comment out `real :: i, j, k` in `dsp_test.f90`, recompile and rerun

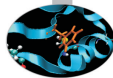


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- ▶ Now comment out `real :: i, j, k` in `dsp_test.f90`, recompile and rerun
- ▶ Now add `implicit none` to `dsp_test.f90` and do it again





Introduction

Fortran Basics

**More Fortran Basics**

My First Fortran Functions

Making it Correct

**Making it Robust**

Copying with Legacy

Wrapping it Up 2

Integer Types and Iterating

More on Compiling and Linking

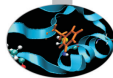
Homeworks



- ▶ Try to pass **integer** variables as actual arguments to **theta()**, **sinc()**, and **rect()**



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- ▶ We can make it aware using **interface** blocks



```
program dsp

  implicit none

  real :: i,j,k

  real, external :: theta, sinc, rect

  print *, 'Enter i, j, k:'
  read(*,*) i, j, k

  write(*,*) 'theta(', i, ')= ', theta(i)
  write(*,*) 'sinc(', i, ')= ', sinc(i)
  write(*,*) 'rect(', j, ', ', k, ')= ', rect(j,k)

end program dsp
```



```
program dsp

  implicit none

  real :: i,j,k

  interface
    function theta(x)
      real :: theta, x
    end function theta
  end interface

  interface
    function sinc(x)
      real :: sinc, x
    end function sinc
  end interface

  interface
    function rect(t, tau)
      real :: rect, t, tau
    end function rect
  end interface

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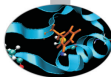


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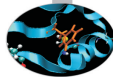


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  - ▶ But life is too short to modify interfaces spread around 56 program units
  - ▶ Good, but still error prone, no better way?

## use Modules, Instead!



- ▶ Modules are the Fortran way to complete and robust management of sets of related routines and more



```
module dsp
  implicit none
contains
  function theta(x) !Heaviside function, useful in DSP
    real :: theta
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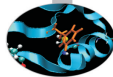
    theta = 1.0
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  end function theta

  function sinc(x) !sinc function as used in DSP
    real :: sinc, xpi
    real, intent(in) :: x
    real, parameter :: pi = acos(-1.0)

    xpi = x*pi
    sinc = 1.0
    if (xpi /= 0.0) sinc = sin(xpi)/xpi
  end function sinc

  function rect(t, tau) !generalized rectangular function, useful in DSP
    real :: rect
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    abs_t = abs(t)
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  end function rect
end module dsp
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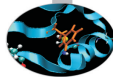
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- ▶ Try it now!





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    real, intent(in) :: x

    theta = 1.0
    if (x < 0.0 ) theta = 0.0
  end function theta

  function sinc(x) !sinc function as used in DSP
    real :: sinc, xpi
    real, intent(in) :: x
    real, parameter :: pi = acos(-1.0)

    xpi = x*pi
    sinc = 1.0
    if (xpi /= 0.0) sinc = sin(xpi)/xpi
  end function sinc

  function rect(t, tau) !generalized rectangular function, useful in DSP
    real :: rect
    real, intent(in) :: t, tau
    real :: abs_t, half_tau

    abs_t = abs(t)
    half_tau = 0.5*tau
    rect = 0.5
    if (abs_t /= half_tau) rect = theta(half_tau-abs_t)
  end function rect
end module dsp
```

# use Modules, Instead!



- ▶ Modules are the Fortran way to complete and robust management of sets of related routines and more
- ▶ Interfaces are automatically defined for each procedure a module contains
- ▶ To use **theta ()**, **sinc ()**, and **rect ()** in a program unit:
  - ▶ just add a **use dsp** statement
  - ▶ before you declare anything else in the unit
- ▶ Try it now!
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- ▶ But there is a lot more to say about modules

# Modules Give You Fine Control



- ▶ A nice colleague handed you the `dsp` module...
- ▶ but you prefer your own version of `rect ()` , which returns 1 on borders:
  - ▶ don't change the module source



- ▶ A nice colleague handed you the `dsp` module...
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  - ▶ **use `dsp, only : theta, sinc`** and keep using your own `rect ()`



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and keep using your own **rect ()**
- ▶ or you already have a function called **theta ()** , called all over your code, and don't want to change it:





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  - ▶ rename the `theta ()` function in `dsp` like this:  
`use dsp, heaviside=>theta`



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- ▶ or maybe both:
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function rect(t, tau)
  implicit none
  real :: rect
  real, intent(in) :: t, tau
  real :: abs_t, half_tau

  abs_t = abs(t)
  half_tau = 0.5*tau
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end function rect
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- ▶ What if `rect ()` is passed a negative argument for `tau`?



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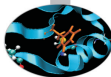
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end function rect
```

- ▶ What if `rect ()` is passed a negative argument for `tau`?
  - ▶ Wrong results
- ▶ Taking the absolute value of `tau` it's a possibility
- ▶ But not a good one, because:
  - ▶ a negative rectangle width is nonsensical
  - ▶ probably flags a mistake in the calling code
  - ▶ and a zero rectangle width is also a problem



```
function rect(t, tau)
  implicit none
  real :: rect
  real, intent(in) :: t, tau
  real :: abs_t, half_tau

  if (tau <= 0.0) stop 'rect() non positive second argument'
  abs_t = abs(t)
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end function rect
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- ▶ A known approach...

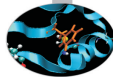




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end function rect
```

- ▶ A known approach...
- ▶ but too rude!
  - ▶ No clue at the argument value
  - ▶ No clue at which call to `rect()` was wrong
  - ▶ And stopping a program in a procedure, called by another procedure, called by another procedure, ..., is widely reputed bad programming practice



```
module dsp
  implicit none
  integer :: dsp_info
  integer, parameter :: DSPERR_DOMAIN = 1
contains
  function theta(x) !Heaviside function, useful in DSP
! code as in previous examples...
  end function theta

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  function rect(t, tau) !generalized rectangular function, useful in DSP
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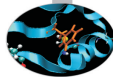
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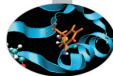
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- ▶ Yes, a module can define variables, too
- ▶ And they will be accessible to all program units using it
- ▶ And yes, **integer** it's another Fortran type
  - ▶ For variables hosting integer numerical values
  - ▶ More on this later...



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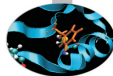
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- ▶ And yes, **return** forces function execution to terminate and return to calling unit



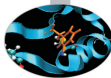
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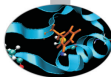
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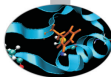


- ▶ Set a module variable to a constant corresponding to the error class
- ▶ And return a sensible result



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- ▶ Then a wise user would do something like this:

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dsp_info = 0
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- ▶ Note: even if Fortran ignores case, constants are often highlighted using all capitals



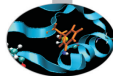
```
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contains

! ...

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  real :: rect
  real, intent(in) :: t, tau
  integer, intent(out) :: info
  real :: abs_t, half_tau

  info = 0
  if (tau <= 0.0) then
    info = DSPERR_DOMAIN
    rect = 0.0
    return
  end if

  abs_t = abs(t)
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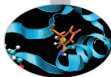
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## Using Arguments to Return Error Codes



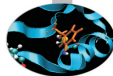
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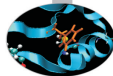
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r = rect(x, width, rect_info)
if (rect_info == DSPERR_DOMAIN) then
    ! take corrective action or fail gracefully
end if
```

- ▶ But this is annoying when the arguments are guaranteed to be correct
  - ▶ **info** can be given the **optional** attribute
  - ▶ and omitted when you feel it's safe: **rect(x, 5.0)**



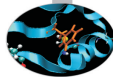
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! ...

function rect(t, tau, info) !generalized rectangular function, useful in DSP
  real :: rect
  real, intent(in) :: t, tau
  integer, intent(out), optional :: info
  real :: abs_t, half_tau

  if (present(info)) info = 0
  if (tau <= 0.0) then
    if (present(info)) info = DSPERR_DOMAIN
    rect = 0.0
    return
  end if

  abs_t = abs(t)
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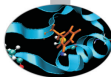
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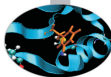
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## Total Robustness



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- ▶ Best strategy: return a NaN and set `dsp_info` in these bad cases
  - ▶ And do it also for non positive values of `tau`
  - ▶ But then the floating point environment configuration should be checked, proper floating point exceptions set...



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  - ▶ And do it also for non positive values of `tau`
  - ▶ But then the floating point environment configuration should be checked, proper floating point exceptions set...
- ▶ Being absolutely robust is difficult
  - ▶ Too advanced stuff to cover in this course
  - ▶ But not an excuse, some robustness is better than none
  - ▶ It's a process to do in steps
  - ▶ Always comment in your code bad cases you don't cover yet!





Introduction

Fortran Basics

**More Fortran Basics**

My First Fortran Functions

Making it Correct

Making it Robust

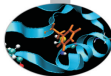
**Copying with Legacy**

Wrapping it Up 2

Integer Types and Iterating

More on Compiling and Linking

Homeworks



```
FUNCTION SINC(X)
  IMPLICIT NONE
  REAL SINC, X, XPI
  REAL PI
  PARAMETER (PI = 3.1415926)

  XPI = X*PI
  SINC = 1.0
  IF (XPI .NE. 0.0) SINC = SIN(XPI)/XPI
END

FUNCTION RECT(T, TAU)
  IMPLICIT NONE
  REAL RECT, T, TAU
  REAL ABS_T, HALF_TAU
  REAL THETA
  EXTERNAL THETA
  INTEGER DSPINFO
  COMMON /DSP/ DSPINFO

  IF (TAU .LE. 0.0) THEN
    DSPINFO = 1
    RECT = 0.0;
    RETURN
  END IF

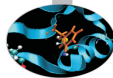
  ABS_T = ABS(T)
  HALF_TAU = 0.5*TAU
  RECT = 0.5
  IF (ABS_T .NE. HALF_TAU) RECT = THETA(HALF_TAU-ABS_T)
END
```

# Many Things are Missing



- ▶ Strange looking relational operators
- ▶ No attributes
  - ▶ Declarations spread over many lines, error prone
- ▶ No initialization expressions
  - ▶ You had to type in the actual number
- ▶ No **intent** i.e. no defense from subtle bugs
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  - ▶ You had to type in the actual number
- ▶ No **intent** i.e. no defense from subtle bugs
- ▶ No **interface**
- ▶ No easy way to share variables among program units
  - ▶ To share you had to use **common** statements
  - ▶ And type in variable types and **common** statements in each unit
  - ▶ And the smallest mistake can turn into a nightmare

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  - ▶ Declarations spread over many lines, error prone
- ▶ No initialization expressions
  - ▶ You had to type in the actual number
- ▶ No **intent** i.e. no defense from subtle bugs
- ▶ No **interface**
- ▶ No easy way to share variables among program units
  - ▶ To share you had to use **common** statements
  - ▶ And type in variable types and **common** statements in each unit
  - ▶ And the smallest mistake can turn into a nightmare
- ▶ Bottom line:
  - ▶ Is **common** good or bad? The jury is still out
  - ▶ We'll not cover them, but you'll encounter them
  - ▶ Read the fine print, or better switch to modules, they are **way much better**



- ▶ You are lucky, and inherit a 4000 lines of code library, coming from the dark ages
  - ▶ Tested and tried



- ▶ You are lucky, and inherit a 4000 lines of code library, coming from the dark ages
  - ▶ Tested and tried
- ▶ But no interface
  - ▶ Thus no compiler checks when you call it
  - ▶ And rewriting a working code in modern language is soooo dangerous...



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  - ▶ And rewriting a working code in modern language is soooo dangerous...
- ▶ Modules come to rescue
  - ▶ They don't need to include the actual code
  - ▶ But they can publish an interface for code which is elsewhere
  - ▶ And then you can use the module in calling program units





```
module dspmod

  implicit none

  interface
    function theta(x)
      real :: theta
      real, intent(in) :: x
    end function theta
  end interface

  interface
    function sinc(x)
      real :: sinc
      real, intent(in) :: x
    end function sinc
  end interface

  interface
    function rect(t, tau)
      real :: rect
      real, intent(in) :: t, tau
    end function rect
  end interface

end module dspmod
```



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My First Fortran Functions

Making it Correct

Making it Robust

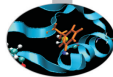
Copying with Legacy

**Wrapping it Up 2**

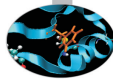
Integer Types and Iterating

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- ▶ A program can be subdivided in more source files
- ▶ Functions and their arguments
- ▶ Arguments are passed to functions by reference
- ▶ **intent** attribute is precious to prevent subtle bugs
- ▶ Intrinsic and external procedures are two different things
- ▶ **parameter** variables
- ▶ Explicit interfaces
- ▶ Modules allow complete management of procedures
- ▶ Modules allow access to variables from many program units
- ▶ Modules can be used to make proper use of legacy, reliable codes



- ▶ Always name constants
- ▶ Test every function you write
  - ▶ Writing specialized programs to do it
- ▶ Use language support and compiler to catch mistakes
- ▶ Use explicit interfaces
- ▶ Use modules
- ▶ Describe all attributes of a variable at declaration
- ▶ Anticipate causes of problems
  - ▶ Find a rational way to react
  - ▶ Fail predictably and in a user friendly way
  - ▶ Robustness it's a long way to do in steps
  - ▶ Comment in your code issues still to address



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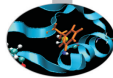
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Homeworks



## ► Euclid's Algorithm

1. Take two integers  $a$  and  $b$
  2. Let  $r \leftarrow a \bmod b$
  3. Let  $a \leftarrow b$
  4. Let  $b \leftarrow r$
  5. If  $b$  is not zero, go back to step 2
  6.  $a$  is the GCD
- 
- Let's implement it and learn some more Fortran



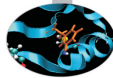
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module number_theory
  implicit none
contains
  function gcd(a, b) ! Greatest Common Divisor
    integer :: gcd
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    gcd = a
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# The Integer Type



- ▶ As we said, **integer** means that a value is an integer
  - ▶ Only integer values, positive, negative or zero
  - ▶ On most platforms, **integer** means a 32 bits value, ranging from  $-2^{31}$  to  $2^{31} - 1$

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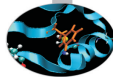
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  - ▶ Try with **kind(0)**, to know the size of a normal **integer**



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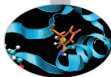


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  - ▶ And works for real values too, or values of any type, for that matter
  - ▶ More on this later
- ▶ Want to know more?
  - ▶ Intrinsic function **huge(0)** returns the greatest positive value an **integer** can assume
  - ▶ Again, we'll be back at this





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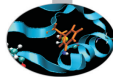
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  implicit none
contains
  function gcd(a, b) ! Greatest Common Divisor
    integer :: gcd
    integer, intent(in) :: a, b
    integer :: gb, t

    gcd = a
    gb = b

    do
      t = mod(gcd,gb)
      gcd = gb
      if (t == 0) exit
      gb = t
    end do
  end function gcd

  function lcm(a, b) ! Least Common Multiple
    integer :: lcm
    integer, intent(in) :: a, b

    lcm = a*b/gcd(a,b)
  end function lcm
end module number_theory
```

# Iterating with `do . . . end do`

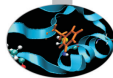


► `do`

*block of statements*

`end do`

1. Executes again and again the *block of statements*
2. And does this forever...
3. ... unless `exit` is executed, forcing execution to proceed at code following `end do`



► `do`

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► In this specific example:



▶ `do`

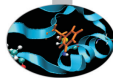
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- ▶ the code following `end do` is the end of the function
- ▶ thus, we could use `return` instead of `exit`, which is legal,
- ▶ but generally regarded bad practice



▶ **do**

*block of statements*

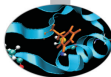
**end do**

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▶ Best practice: do not bail out of a function from inside a loop, particularly a long one



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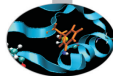
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Homeworks



- ▶ Put the code in file `numbertheory.f90`
- ▶ Write a program to test both `gcd()` and `lcm()` on a pair of integer numbers
- ▶ Test it:
  - ▶ with pairs of small positive integers
  - ▶ with the following pairs: 15, 18; -15, 18; 15, -18; -15, -18; 0, 15; 15, 0; 0, 0





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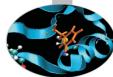
    lcm = a*b/gcd(a,b)
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```

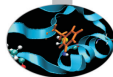


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  - ▶ with pairs of small positive integers
  - ▶ with the following pairs: 15, 18; -15, 18; 15, -18; -15, -18; 0, 15; 15, 0; 0, 0
- ▶ In some cases, we get wrong results or runtime errors
  - ▶ Euclid's algorithm is only defined for positive integers

## Let's Generalize to the Whole Integer Set

- ▶  $\text{gcd}(a, b)$  is non negative, even if  $a$  or  $b$  is less than zero
  - ▶ Taking the absolute value of  $a$  and  $b$  using `abs ()` will do





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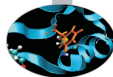
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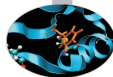
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  - ▶ Conditional statements will do





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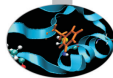
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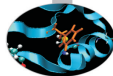
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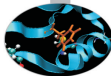


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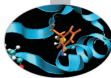
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  - ▶ and with the pair: 1000000, 1000000



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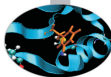
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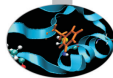
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# Beware of Type Ranges



- ▶  $a*b/\text{gcd}(a,b)$  same as  $(a*b) / \text{gcd}(a,b)$

# Beware of Type Ranges



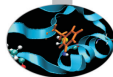
- ▶  $a*b/\text{gcd}(a,b)$  same as  $(a*b) / \text{gcd}(a,b)$
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- ▶ Easy fix:  $\text{gcd}(a,b)$  is an exact divisor of  $b$



## GCD &amp; LCM: Preventing Overflow



```
module number_theory
  implicit none
contains
  function gcd(a, b) ! Greatest Common Divisor
    integer :: gcd
    integer, intent(in) :: a, b
    integer :: gb, t

    gcd = abs(a)
    gb = abs(b)

    if (a == 0) gcd = gb
    if (a == 0 .or. b == 0) return

    do
      t = mod(gcd,gb)
      gcd = gb
      if (t == 0) exit
      gb = t
    end do
  end function gcd

  function lcm(a, b) ! Least Common Multiple
    integer :: lcm
    integer, intent(in) :: a, b

    if (a == 0 .and. b == 0) then
      lcm = 0 ; return
    end if

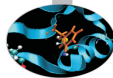
    lcm = a*(b/gcd(a,b))
  end function lcm
end module number_theory
```



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- ▶ Easy fix:  $\text{gcd}(a, b)$  is an exact divisor of  $b$
- ▶ Try and test it:
  - ▶ with pairs of small positive integers
  - ▶ on the following pairs: 15, 18; -15, 18; 15, -18; -15, -18; 0, 15; 15, 0; 0, 0
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  - ▶ with the pair: 1000000, 1000000
  - ▶ and let's test also with: 1000000, 1000001



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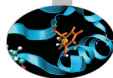
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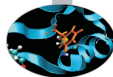


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  - ▶ or `-1` if none is wide enough
- ▶ `integer` accepts an optional *kind type parameter*
  - ▶ `integer(kind=selected_int_kind(9)) :: di` usually makes `di` a 32 bits wide variable
  - ▶ `integer(kind=selected_int_kind(18)) :: wi` makes `wi` a 64 bits wide variable
  - ▶ `integer(selected_int_kind(18)) :: wi` will also do

## GCD &amp; LCM: Let's Make Kind Explicit



```
module number_theory
  implicit none
  contains
    function gcd9(a, b) ! Greatest Common Divisor
      integer(selected_int_kind(9)) :: gcd9
      integer(selected_int_kind(9)), intent(in) :: a, b
      integer(selected_int_kind(9)) :: gb, t

      gcd9 = abs(a)
      gb = abs(b)

      if (a == 0) gcd9 = gb
      if (a == 0 .or. b == 0) return

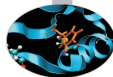
      do
        t = mod(gcd9, gb)
        gcd9 = gb
        if (t == 0) exit
        gb = t
      end do
    end function gcd9

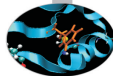
    function lcm9(a, b) ! Least Common Multiple
      integer(selected_int_kind(9)) :: lcm9
      integer(selected_int_kind(9)), intent(in) :: a, b

      if (a == 0 .and. b == 0) then
        lcm9 = 0 ; return
      end if

      lcm9 = a*(b/gcd9(a,b))
    end function lcm9
  end module number_theory
```

- ▶ And let's add support for a wider integer range





```
! add right after: end function lcm9

function gcd18(a, b) ! Greatest Common Divisor
  integer(selected_int_kind(18)) :: gcd18
  integer(selected_int_kind(18)), intent(in) :: a, b
  integer(selected_int_kind(18)) :: gb, t

  gcd18 = abs(a)
  gb = abs(b)

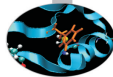
  if (a == 0) gcd18 = gb
  if (a == 0 .or. b == 0) return

  do
    t = mod(gcd18,gb)
    gcd18 = gb
    if (t == 0) exit
    gb = t
  end do
end function gcd18

function lcm18(a, b) ! Least Common Multiple
  integer(selected_int_kind(18)) :: lcm18
  integer(selected_int_kind(18)), intent(in) :: a, b

  if (a == 0 .and. b == 0) then
    lcm18 = 0 ; return
  end if

  lcm18 = a*(b/gcd18(a,b))
end function lcm18
```

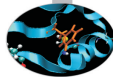


- ▶ And let's add support for a wider integer range
- ▶ Wait!
  - ▶ Now we have to remember to call the right function, depending on the integer kind
  - ▶ But this is not Fortran style: we didn't have to change the call to intrinsic **abs ()**, it's name is generic
  - ▶ Can we do better?



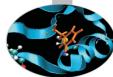


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  - ▶ `interface` blocks come to rescue

# GCD & LCM: Making it Generic



```
module number_theory
  implicit none

  private gcd9, lcm9, gcd18, lcm18

  interface gcd
    module procedure gcd9, gcd18
  end interface

  interface lcm
    module procedure lcm9, lcm18
  end interface

contains

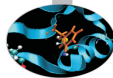
  function gcd9(a, b) ! Greatest Common Divisor
  ! code as before
  end function gcd9

  function lcm9(a,b) ! Least Common Multiple
  ! code as before
  end function lcm9

  function gcd18(a, b) ! Greatest Common Divisor
  ! code as before
  end function gcd18

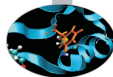
  function lcm18(a,b) ! Least Common Multiple
  ! code as before
  end function lcm18

end module number_theory
```



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  - ▶ Beware: specific functions under a same generic interface must differ in type of at least one argument

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module number_theory
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  end interface

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    module procedure lcm9, lcm18
  end interface

contains

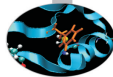
  function gcd9(a, b) ! Greatest Common Divisor
  ! code as before
  end function gcd9

  function lcm9(a,b) ! Least Common Multiple
  ! code as before
  end function lcm9

  function gcd18(a, b) ! Greatest Common Divisor
  ! code as before
  end function gcd18

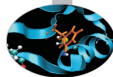
  function lcm18(a,b) ! Least Common Multiple
  ! code as before
  end function lcm18

end module number_theory
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    module procedure gcd9, gcd18
  end interface

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    module procedure lcm9, lcm18
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contains

  function gcd9(a, b) ! Greatest Common Divisor
    ! code as before
  end function gcd9

  function lcm9(a,b) ! Least Common Multiple
    ! code as before
  end function lcm9

  function gcd18(a, b) ! Greatest Common Divisor
    ! code as before
  end function gcd18

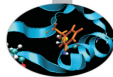
  function lcm18(a,b) ! Least Common Multiple
    ! code as before
  end function lcm18

end module number_theory
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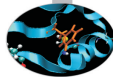


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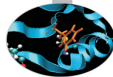




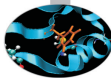
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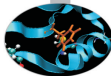
Wider Integer Types

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```
FUNCTION GCD18(A, B)
  INTEGER*8 GCD18, A, B
  INTEGER*8 GB, T

  GCD18 = A
  GB = B

1   T = MOD(GCD18,GB)
   GCD18 = GB
   IF (T .EQ. 0) GO TO 2
   GB = T
   GO TO 1

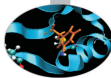
2   CONTINUE
END

FUNCTION LCM18(A, B)
  INTEGER*8 LCM18, A, B
  INTEGER*8 GCD18
  EXTERNAL GCD18

  LCM18 = A*B/GCD18(A,B)
END
```



- ▶ No structured endless loops
  - ▶ Labels and `GO TO`s were used instead
- ▶ **CONTINUE** was a no-op
  - ▶ Used to mark destination of jumps
  - ▶ No comment
- ▶ **INTEGER\*8** was used to declare an 8 bytes integer variable
  - ▶ Absolutely non standard
  - ▶ As are **INTEGER\*1**, **INTEGER\*2**, **INTEGER\*4**, **REAL\*4**, **REAL\*8**, **COMPLEX\*8**, **COMPLEX\*16**
- ▶ Many dialects
  - ▶ Many proprietary extensions used to be developed
  - ▶ And then copied among vendors for compatibility reasons
  - ▶ Many extensions were eventually standardized
  - ▶ But not all of them!
  - ▶ They still lurk around, and can be tempting: resist!



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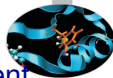
More on Compiling and Linking

Homeworks



- ▶ There are many integer types
  - ▶ With implementation dependent ranges
  - ▶ Selectable by kind type parameters
  - ▶ Whose limits can be devised using **huge ()** or **range ()**
- ▶ Library functions have generic names, good for most types
- ▶ And you can write your own generic interfaces
- ▶ Behavior on integer overflow is implementation defined
  - ▶ Some control is possible using parentheses
- ▶ Blocks of statements can be iterated forever...
  - ▶ ... and **exit** gets off the roundabout
- ▶ Logical conditions can be combined using **.or.** and **.and.** operators





- ▶ Do not rely on type sizes, they are implementation dependent
- ▶ Do not leave a function from inside a loop
- ▶ Think of intermediate results in expressions: they can overflow or underflow
- ▶ Be consistent with Fortran approach
  - ▶ E.g. writing generic interfaces
  - ▶ Even if it costs more work
  - ▶ Even if it costs learning more Fortran
  - ▶ Once again, you can do it in steps
  - ▶ You'll appreciate it in the future
- ▶ Hide implementation details as much as possible
  - ▶ You'll never regret
- ▶ Resist the temptation of old Fortran or non standard extensions
  - ▶ Will pay back in the future



Introduction

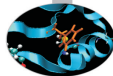
Fortran Basics

More Fortran Basics

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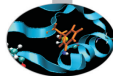
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Homeworks



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# Compiler Errors and Warnings



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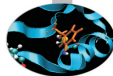


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- ▶ If given earlier ...
- ▶ Something is an error if not in Fortran 95 standard
  - ▶ Use **-std=f95** to force reference standard



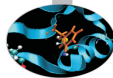
Creating an executable from source files is in general a three phase process:

- ▶ pre-processing:
  - ▶ each source file is read by the pre-processor
    - ▶ substitute (**#define**) MACROS
    - ▶ insert code by **#include** statements
    - ▶ insert or delete code evaluating **#ifdef**, **#if** ...
- ▶ compiling:
  - ▶ each source file is translated into an object code file
    - ▶ an object code file is an organised collection of symbols, referring to variables and functions defined or used in the source file
- ▶ linking:
  - ▶ object files should be combined together to build a single executable program
  - ▶ every symbol should be resolved
    - ▶ symbols can be defined in your object files
    - ▶ or available in other object code (external libraries)



- ▶ When you give the command:

```
user@cineca$> gfortran dsp.f90 dsp_test.f90
```



- ▶ When you give the command:

```
user@cineca$> gfortran dsp.f90 dsp_test.f90
```

- ▶ It's like going through three steps



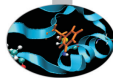
- ▶ When you give the command:

```
user@cineca$> gfortran dsp.f90 dsp_test.f90
```

- ▶ It's like going through three steps
- ▶ Pre-processing

```
user@cineca$> gfortran -E -cpp dsp.f90  
user@cineca$> gfortran -E -cpp dsp_test.f90
```

- ▶ **-E -cpp** option, tells **gfortran** to stop after pre-process
- ▶ Simply calls **cpp** (automatically invoked if the file extension is **F90**)
- ▶ Output sent to standard output



- ▶ When you give the command:

```
user@cineca$> gfortran dsp.f90 dsp_test.f90
```

- ▶ It's like going through three steps
- ▶ Pre-processing

```
user@cineca$> gfortran -E -cpp dsp.f90  
user@cineca$> gfortran -E -cpp dsp_test.f90
```

- ▶ **-E -cpp** option, tells **gfortran** to stop after pre-process
- ▶ Simply calls **cpp** (automatically invoked if the file extension is **F90**)
- ▶ Output sent to standard output
- ▶ Compiling sources

```
user@cineca$> gfortran -c dsp.f90  
user@cineca$> gfortran -c dsp_test.f90
```

- ▶ **-c** option tells **gfortran** to only compile the source
- ▶ An object file **.o** is produced from each source file



- ▶ Linking object files together

```
user@cineca$> gfortran dsp.o dsp_test.o
```

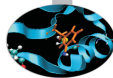
- ▶ To resolve symbols defined in external libraries, specify:
  - ▶ which libraries to use (`-l` option)
  - ▶ in which directories they are (`-L` option)
- ▶ How to link the library `libdsp.a` in `/mypath`

```
user@cineca$> gfortran file1.o file2.o -L/mypath -ldsp
```

- ▶ How to create and link the DSP library:

```
user@cineca$> gfortran -c dsp.f90  
ar curv libdsp.a dsp.o  
ranlib libdsp.a  
gfortran test_dsp.f90 -L. -ldsp
```

- ▶ `ar` create the archive `libdsp.a` containing `dsp.o`
  - ▶ `ranlib` generate index to archive
- ▶ To include file like `.mod`, specify
  - ▶ in which directories they are (`-I` option)



Introduction

Fortran Basics

More Fortran Basics

Integer Types and Iterating

More on Compiling and Linking

Homeworks





- ▶ Write a program that reads an integer value `limit` and prints the first `limit` prime numbers
  - ▶ Use the GCD function to identify those numbers
  - ▶ After testing the basic version, handle negative `limit` values: print an error message and attempt to read the value again



- ▶ Write a module containing a function that takes an integer  $n$  as input, and returns the  $n$ -th element of the Fibonacci series  $f_n$
- ▶ Hint:
  - ▶  $F_0 = 0$
  - ▶  $F_1 = 1$
  - ▶  $F_n = F_{n-1} + F_{n-2}$
- ▶ Write a main program to test your function
  - ▶ Read  $n$  from standard input
  - ▶ Try with  $n=2, 10, 40, 46, 48, \dots$
  - ▶ What's the greatest  $n := \max n$ , for which  $f_n$  is representable by a default integer? (**huge** can help to find it out)
  - ▶ Use this information to handle too large values of  $n$  in your function:
  - ▶ If  $n > \max n$  print an error message and return -1



## Part II

# A Fortran Survey 2

Passing functions as arguments of procedures. Conditional and numerical loops. Managing the precision, conversions, overflow, underflow, Inf e NaN. Expressions and subexpressions with mixed types. Types, operators and logical expressions. Type character and intrinsic functions for strings. Subroutine. Array. Default constructor for arrays and implicit loops. Assumed-shape array and automatic object. Expressions with arrays and conformity.



More Flow Control  
Numerical Integration  
Wrapping it Up 4

Fortran Intrinsic Types, Variables and Math

Arrays



The code in this section is meant for didactical purposes only.

It is deliberately naive: focus is on language aspects, not on precision or accuracy.

As a consequence, it is prone to numerical problems.



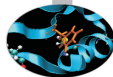
More Flow Control  
Numerical Integration  
Wrapping it Up 4

Fortran Intrinsic Types, Variables and Math

Arrays



- ▶ Let's use the trapezoidal rule to estimate  $\int_a^b f(x) dx$
- ▶ Dividing the interval  $[a, b]$  into  $n$  equal sized slices, it boils down to:  
$$\int_a^b f(x) dx \approx \frac{b-a}{n} \left( \frac{1}{2}f(a) + \frac{1}{2}f(b) + \sum_{k=1}^{n-1} f\left(a + k\frac{b-a}{n}\right) \right)$$
- ▶ And to make it more juicy, let's make a succession of estimates, doubling  $n$  each time, until the estimate seems stable



```

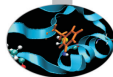
module integrals
  implicit none
  contains
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      ! recursive approximation of integral
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      real :: trap_int
      real, intent(in) :: a, b, tol
      ! integration interval and tolerance
      interface
        real function f(x)
          ! function to integrate
          real, intent(in) :: x
        end function f
      end interface
      integer, parameter :: maxsteps = 2**23
      integer :: steps, i
      real :: acc, dx, prev_estimate, estimate

      steps = 2
      prev_estimate = 0.0 ; estimate = huge(0.0)
      dx = (b - a)*0.5
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      do i=1, steps, 2
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      end do
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      trap_int = estimate
    end function trap_int
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```

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  contains
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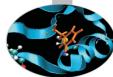
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```



- ▶ Yes, a function can be passed as an argument to another function!
- ▶ Simply pass the name on call, like this:  
`g = trap_int(-pi, pi, sinc, 0.0001)`
- ▶ And then the function can be called using the dummy argument name
- ▶ And this can be done for any procedure
- ▶ And allows for very generic code to be written
  - ▶ I.e. reuse the same routine to integrate different functions in the same program



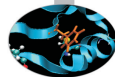
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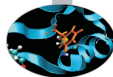
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- ▶ And then the function can be called using the dummy argument name
- ▶ And this can be done for any procedure
- ▶ And allows for very generic code to be written
  - ▶ I.e. reuse the same routine to integrate different functions in the same program
- ▶ Integer and real values can be mixed in expressions
  - ▶ As well as values of same type but different kind
  - ▶ And the right thing will be done
  - ▶ Which is: when two values of different type/kind meet each other at a binary operator, the one with smaller numeric range is converted to the other



```

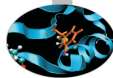
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        ! only contributions from new points
        acc = acc + f(a + i*dx)
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      steps = steps*2
      if (steps > maxsteps) exit conv
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    end do conv

      trap_int = estimate
    end function trap_int
end module
  
```

# Iterating with `do while ... end do`

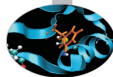


► `do while` (*logical-condition*)

*block of statements*

`end do`

1. Evaluates *logical-condition*
2. If *logical-condition* is false, goes to 5
3. Executes the *block of statements*
4. Goes back to 1
5. Execution proceeds to the statement following `end do`



```

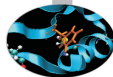
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        end function f
      end interface
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      real :: acc, dx, prev_estimate, estimate

      steps = 2
      prev_estimate = 0.0 ; estimate = huge(0.0)
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      acc = (f(a) + f(b))*0.5

      conv: do while (abs(estimate - prev_estimate) > tol)
        prev_estimate = estimate
        do i=1, steps, 2
          ! only contributions from new points
          acc = acc + f(a + i*dx)
        end do
        estimate = acc*dx
        steps = steps*2
        if (steps > maxsteps) exit conv
        dx = dx*0.5
      end do conv

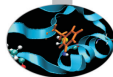
      trap_int = estimate
    end function trap_int
  end module

```

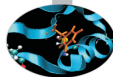


- ▶ `do while` (*logical-condition*)  
*block of statements*  
`end do`
  1. Evaluates *logical-condition*
  2. If *logical-condition* is false, goes to 5
  3. Executes the *block of statements*
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  5. Execution proceeds to the statement following `end do`
  
- ▶ `do` loops too can be given a name
  1. And it can be used on `exit` statements to make the flow more evident
  2. Particularly for nested loops





- ▶ `do while` (*logical-condition*)  
*block of statements*  
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  5. Execution proceeds to the statement following `end do`
- ▶ `do` loops too can be given a name
  1. And it can be used on `exit` statements to make the flow more evident
  2. Particularly for nested loops
- ▶ Best practices:
  1. use names to mark loops when they are long or belong to a deep nest
  2. NEVER, NEVER permit your code to loop forever for some inputs



```

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  contains
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# Iterating with Counted do



- ▶ **do** *var* = *init*, *limit* [, *step*]  
    *block of statements*  
**end do**

# Iterating with Counted do



- ▶ **do var = init, limit [, step]**  
*block of statements*  
**end do**
  1. Sets *step* to 1, if none was specified

# Iterating with Counted do



- ▶ **do** *var* = *init*, *limit* [, *step*]  
    *block of statements*  
**end do**
  1. Sets *step* to 1, if none was specified
  2. Assign the *init* value to *var*



► **do** *var* = *init*, *limit* [, *step*]  
    *block of statements*  
**end do**

1. Sets *step* to 1, if none was specified
2. Assign the *init* value to *var*
3. Evaluates  $n_{iter} = \max\{0, \lfloor (limit - init + step) / step \rfloor\}$



► **do** *var* = *init*, *limit* [, *step*]  
    *block of statements*  
**end do**

1. Sets *step* to 1, if none was specified
2. Assign the *init* value to *var*
3. Evaluates  $n_{iter} = \max\{0, \lfloor (limit - init + step) / step \rfloor\}$
4. If  $n_{iter}$  is zero goes to 6



► **do** *var* = *init*, *limit* [, *step*]  
    *block of statements*  
**end do**

1. Sets *step* to 1, if none was specified
2. Assign the *init* value to *var*
3. Evaluates  $n_{iter} = \max\{0, \lfloor (limit - init + step) / step \rfloor\}$
4. If  $n_{iter}$  is zero goes to 6
5. Executes  $n_{iter}$  times the *block of statements*, adding *step* to *var* at the end of each *block of statements*





► **do** *var* = *init*, *limit* [, *step*]  
    *block of statements*  
**end do**

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6. Execution proceeds to the statement following **end do**



- ▶ **do** *var* = *init*, *limit* [, *step*]  
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  6. Execution proceeds to the statement following **end do**
- ▶ *var*, *init*, *limit*, and *step* should be integers



- ▶ **do** *var* = *init*, *limit* [, *step*]  
    *block of statements*  
**end do**
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- ▶ **var**, *init*, *limit*, and *step* should be integers
  - ▶ Mandatory in Fortran 2003



- ▶ **do var = init, limit [, step]**  
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- ▶ **var**, *init*, *limit*, and *step* should be integers
  - ▶ Mandatory in Fortran 2003
  - ▶ Reals can be used up to Fortran 95, but a bad idea, for both performance and reliability issues

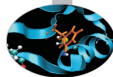


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  - ▶ Reals can be used up to Fortran 95, but a bad idea, for both performance and reliability issues
- ▶ Less flexible than a **do while** but more efficient execution (**exit** works, anyway)



- ▶ **do var = init, limit [, step]**  
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- ▶ **var, init, limit, and step** should be integers
  - ▶ Mandatory in Fortran 2003
  - ▶ Reals can be used up to Fortran 95, but a bad idea, for both performance and reliability issues
- ▶ Less flexible than a **do while** but more efficient execution (**exit** works, anyway)
- ▶ Best practice: do not give name to very tight loops

## Hands-on Session #1



```

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  contains
    function trap_int(a,b,f,tol)
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      real :: trap_int
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      end do
      estimate = acc*dx
      steps = steps*2
      if (steps > maxsteps) exit conv
      dx = dx*0.5
    end do conv

      trap_int = estimate
    end function trap_int
end module

```



- ▶ Write a program to exercise `trap_int()` on functions with known integrals





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- ▶ Then take care of what was left out



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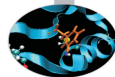
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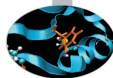
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- ▶ Counted loops looked like this:

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      do 10, i=1,10,3  
        write(*,*) i  
10    continue
```



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- ▶ Counted loops looked like this:

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      do 10, i=1,10,3  
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10    continue
```

- ▶ **do while, exit, end do** weren't there...
  - ▶ ... at least in the standard...
  - ▶ but are often found in codes, as dialect extensions.



More Flow Control

Numerical Integration

Wrapping it Up 4

Fortran Intrinsic Types, Variables and Math

Arrays



- ▶ More flow control
  - ▶ Procedure arguments
  - ▶ **do while**
  - ▶ Counted **do**
- ▶ Mixed-mode expressions
- ▶ Name your loops
  - ▶ Particularly if long or nested
  - ▶ Particularly if you **exit** them
  - ▶ But don't do it for short ones
- ▶ Prevent any loop from running forever for some program inputs



## More Flow Control

### Fortran Intrinsic Types, Variables and Math

- Integer Types

- Floating Types

- Expressions

- Arithmetic Conversions

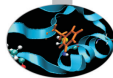
- More Intrinsic Types

## Arrays



- ▶ Computing == manipulating data and calculating results
  - ▶ Data are manipulated using internal, binary formats
  - ▶ Data are kept in memory locations and CPU registers
- ▶ Fortran doesn't make assumptions on internal data representations
  - ▶ And tries to abstract
  - ▶ Most CPU are similar but all have peculiarities
  - ▶ Some details depend on the specific executing (a.k.a. target) hardware architecture and software implementation
  - ▶ Fortran provides facilities to translate between internal formats and human readable ones
- ▶ Fortran allows programmers to:
  - ▶ think in terms of data types and named containers
  - ▶ disregard details on actual memory locations and data movements

# Fortran is a Strongly Typed Language



- ▶ Each literal constant has a type
  - ▶ Dictates internal representation of the data value
- ▶ Each variable has a type
  - ▶ Dictates content internal representation and amount of memory
  - ▶ Type must be specified in a declaration before use
  - ▶ Unless you are so naive to rely on implicit declaration
- ▶ Each expression has a type
  - ▶ And subexpressions have too
  - ▶ Depends on operators and their arguments
- ▶ Each function has a type
  - ▶ That is the type of the returned value
  - ▶ Specified in function interface
- ▶ Procedure arguments have types
  - ▶ I.e. type of arguments to be passed in calls
  - ▶ Specified in procedure interface
  - ▶ If the compiler doesn't know the interface, it will blindly pass whatever you provide



More Flow Control

Fortran Intrinsic Types, Variables and Math

Integer Types

Floating Types

Expressions

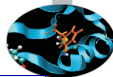
Arithmetic Conversions

More Intrinsic Types

Arrays



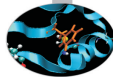
# Integer Types (as on most CPUs)



Type	Sign	Usual <b>huge ()</b>	Usual Width (bits)	Usual Size (bytes)
<code>integer(selected_int_kind(2))</code>	+/-	127	8	1
<code>integer(selected_int_kind(5))</code>	+/-	32767	16	2
<code>integer</code> <code>integer(kind(0))</code> <code>integer(selected_int_kind(9))</code>	+/-	2147483647	32	4
<code>integer(selected_int_kind(18))</code>	+/-	9223372036854775807	64	8

- ▶ `selected_int_kind(n)` returns the least type able to host  $10^n$
- ▶ `selected_int_kind(n)` returns -1 if no suitable type is available
- ▶ New platform/compiler? Always check maximum headroom with **huge ()** or **range ()**
- ▶ As we said, on most platforms **kind ()** returns the byte size, but it's not standard

# Integer Literal Constants

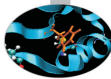


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- ▶ In a non portable way:

`-123456_8`

- ▶ Or in a portable way:

```
integer, parameter :: i8=selected_int_kind(18)  
-123456_i8
```

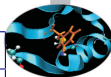


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```
- ▶ Rule of thumb:
  - ▶ write the number as is, if it is in default integer kind range
  - ▶ otherwise, specify kind



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- ▶ Rule of thumb:
  - ▶ write the number as is, if it is in default integer kind range
  - ▶ otherwise, specify kind
- ▶ Remember:
  - ▶ do not write `spokes = bicycles*2*36`
  - ▶ `integer, parameter :: SpokesPerWheel = 36`
  - ▶ code will be more readable, and you'll be ready for easy changes



Function	Returns
<b>abs (i)</b>	$ i $
<b>sign (i, j)</b>	$ i $ if $j \geq 0$ , $- i $ otherwise
<b>dim (i, j)</b>	if $i > j$ returns $i - j$ else returns 0
<b>mod (i, j)</b>	Remainder function $i - \text{int}(i/j) \times j$
<b>modulo (i, j)</b>	Modulo function $i - \text{floor}(i/j) \times j$
<b>min (i, j[, ...])</b>	$\min\{i, j[, \dots]\}$
<b>max (i, j[, ...])</b>	$\max\{i, j[, \dots]\}$

- ▶ Use like: **a = abs (b+i) + c**
- ▶ More functions are available to manipulate values
  - ▶ E.g. for bit manipulations on binary computers
  - ▶ We'll not cover them in this course, you can learn more about if you need to
- ▶ They can be found under different names (e.g. **iabs ()**): these are relics from the past



More Flow Control

Fortran Intrinsic Types, Variables and Math

Integer Types

**Floating Types**

Expressions

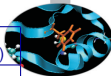
Arithmetic Conversions

More Intrinsic Types

Arrays



# Floating Types (as on most CPUs)



Type	Usual huge ()	Usual Width (bits)	Usual Size (bytes)
real real(kind(0.0)) real(selected_real_kind(6))	3.40282347e38	32	4
double precision real(kind(0.0d0)) real(selected_real_kind(15)) real(selected_real_kind(18))	1.79769313486231573e308	64	8
complex complex(kind(0.0)) complex(selected_real_kind(6))	NA	NA	8
complex(kind(0.0d0)) complex(selected_real_kind(15)) complex(selected_real_kind(18))	NA	NA	16
complex(selected_real_kind(18))	NA	NA	20 or 32

- ▶ In practice, always in IEEE Standard binary format, but not a Standard requirement
- ▶ **selected\_real\_kind()** gets number of significant decimal digits, plus a second optional argument for exponent range, returns negative result if no suitable type is available
- ▶ **tiny()** returns smallest positive value
- ▶ New platform/compiler? Always check maximum headroom with **huge()** or **range()**

## real Literal Constants

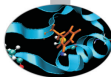


- ▶ Need something to distinguish them from integers
  - ▶ Decimal notation: **1.0**, **-17.**, **.125**, **0.22**
  - ▶ Exponential decimal notation: **2e19** ( $2 \times 10^{19}$ ), **-123.4e9** ( $-1.234 \times 10^{11}$ ), **.72e-6** ( $7.2 \times 10^{-7}$ )

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- ▶ By default, `kind(0.0)`
- ▶ Unless you specify it
  - ▶ For double precision only:  
`-1.23456d5`
  - ▶ For all kinds:  

```
integer, parameter :: r8=selected_real_kind(15)
-123456.0_r8
```



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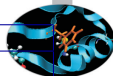
`-1.23456d5`

- ▶ For all kinds:

```
integer, parameter :: r8=selected_real_kind(15)  
-123456.0_r8
```

- ▶ Remember:

- ▶ do not write `charge = protons*1.602176487E-19`
  - ▶ `real,parameter::UnitCharge=1.602176487E-19`
  - ▶ it will come handier when more precise measurements will be available



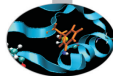
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<b>aint (x)<sup>2</sup>, int (x)<sup>1,2</sup></b>	if $x > 0$ returns $\lfloor x \rfloor$ else returns $\lceil x \rceil$
<b>anint (x)<sup>2</sup>, nint (x)<sup>1,2</sup></b>	nearest integer to $x$
<b>floor (x)<sup>1,2</sup>, ceiling (x)<sup>1,2</sup></b>	$\lfloor x \rfloor$ , $\lceil x \rceil$
<b>fraction (x)</b>	fractional part of $x$
<b>nearest (x, s)</b>	next representable value to $x$ , in direction given by the sign of $s$
<b>spacing (x)</b>	absolute spacing of numbers near $x$
<b>max (x, y[, ...])</b>	$\max\{x, y[, \dots]\}$
<b>min (x, y[, ...])</b>	$\min\{x, y[, \dots]\}$
1. Result is of integer type 2. Accept an optional argument for kind type of the result	

- ▶ They can be found under different names (e.g. **dabs ()**): these are relics from the past
- ▶ More functions are available to manipulate values
  - ▶ Mostly in the spirit of IEEE Floating Point Standard
  - ▶ We'll not cover them in this course, but encourage you to learn more about



Functions	Compute
<b>sqrt (x)</b>	$\sqrt{x}$
<b>sin (x), cos (x), tan (x), asin (x), acos (x), atan (x)</b>	Trigonometric functions
<b>atan2 (x, y)</b>	Arc tangent in $(-\pi, \pi]$
<b>exp (x), log (x), log10 (x)</b>	$e^x$ , $\log_e x, \log_{10} x$
<b>sinh (x), cosh (x), tanh (x)</b>	Hyperbolic functions

- ▶ Again, they can be found under different names (e.g. **dcos ()**): these are relics from the past



Functions	Compute
<b>abs (z)</b> ,	$ z $ ,
<b>aimag (z)</b>	imaginary part of <b>z</b> ,
<b>real (z)</b> <sup>1</sup>	real part of <b>z</b>
<b>cmplx (x, y)</b> <sup>1</sup>	converts from real to complex
<b>conj (z)</b>	Complex conjugate of <b>z</b>
<b>sqrt (z)</b>	$\sqrt{z}$
<b>sin (z)</b> , <b>cos (z)</b>	sine and cosine
<b>exp (z)</b> ,	$e^z$ ,
<b>log (z)</b>	$\log_e z$
1. Accept an optional argument for kind type of the result	

- ▶ Once again, they can be found under different names (e.g. **cabs ()**): again, these are relics from the past





- ▶ The intrinsic function **precision (x)** for real or complex **x** returns the number of significant decimal digits.
- ▶ Write a **module** which defines the **kind** constant for single, double and quadruple real precision



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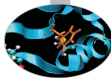
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- ▶ To gain confidence: write a small program to print out **range** and **huge** values for these kinds



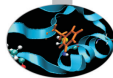
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- ▶ Something going wrong?



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- ▶ To gain confidence: write a small program to print out **range** and **huge** values for these kinds
- ▶ Something going wrong?
- ▶ GNU Fortran compiler, up to release 4.5, lacks support for the quad-precision
- ▶ If you are using, Linux load the GNU compiler version 4.6 and try again:  
**module load gcc**



- ▶ Use the `real_kinds` module to rewrite `dsp` module functions to support both single and double precision
- ▶ And make all of them generic procedures
- ▶ Modify your test program to see exercise the new `dsp` module



More Flow Control

Fortran Intrinsic Types, Variables and Math

Integer Types

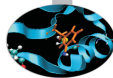
Floating Types

**Expressions**

Arithmetic Conversions

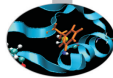
More Intrinsic Types

Arrays



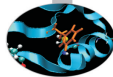
- ▶ Binary operators  $+$ ,  $-$ ,  $*$  (multiplication) and  $/$  have the usual meaning and behavior
- ▶ And so do unary operators  $-$  and  $+$
- ▶ Precedence
  - ▶  $-a*b + c/d$  same as  $((-a)*b) + (c/d)$
  - ▶  $-a + b$  same as  $(-a) + b$
- ▶ Associativity of binary ones is from left to right
  - ▶  $a + b + c$  same as  $(a + b) + c$
  - ▶  $a*b/c*d$  same as  $((a*b)/c)*d$
- ▶ Explicit  $($  and  $)$  override precedence and associativity
- ▶  $**$  is the exponentiation operator
- ▶ Assignment:  $=$ 
  - ▶ Assigns the value of expression on right hand side to a variable on the left hand side
  - ▶ Prior to first assignment, a variable content is *undefined*



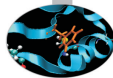


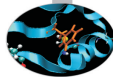
- ▶ All types are limited in range
- ▶ What about:
  - ▶ **huge** (0) + 1? (too big)
  - ▶ **-huge** (0.0) \* 3.0? (too negative)
- ▶ Technically speaking, this is an arithmetic *overflow*
- ▶ And division by zero is a problem too
- ▶ For integer types, the Standard says:
  - ▶ behavior and results are unpredictable
  - ▶ i.e. up to the implementation
- ▶ For real types, it also depends on the floating point environment
  - ▶ i.e. how behavior is configured for those cases
  - ▶ you could get **-huge** (0.0), or a **NaN**, or **-Inf**
- ▶ Best practice: NEVER rely on behaviors observed with a specific architecture and/or compiler

- ▶ Just imagine both functions `foo(x, y)` and `bar(x, y)` modify their actual arguments, or do I/O

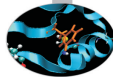


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  - ▶ As you'll remember, these are known as *side effects*



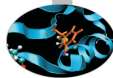


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- ▶ Now imagine you meet code like this:  
`t = foo(a, b) - bar(b, a)`  
`q = mod(foo(a, b), bar(a, b))`



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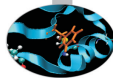
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- ▶ Code like this is evil!



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- ▶ Order of subexpressions evaluation is implementation dependent!

# Order of Subexpressions Evaluation



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q = mod(foo(a, b), bar(a, b))
```
- ▶ Code like this is evil!
- ▶ Order of subexpressions evaluation is implementation dependent!
  - ▶ You don't know in advance the order in which `foo()` and `bar()` are called



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  - ▶ As you'll remember, these are known as *side effects*
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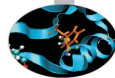
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- ▶ Just imagine both functions `foo(x, y)` and `bar(x, y)` modify their actual arguments, or do I/O
  - ▶ As you'll remember, these are known as *side effects*
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- ▶ Ditto for order of evaluation of function arguments!
- ▶ NEVER! NEVER write code that relies on order of evaluation of subexpressions, or actual arguments!



More Flow Control

Fortran Intrinsic Types, Variables and Math

Integer Types

Floating Types

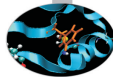
Expressions

**Arithmetic Conversions**

More Intrinsic Types

Arrays

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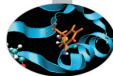


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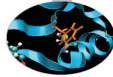
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  - ▶ while:  
**a = b\*2\_i8 + 1**  
is OK



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# Subtle Type Conversion Traps



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# Subtle Type Conversion Traps

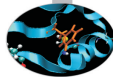


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- ▶ A 32 bits `real` has fewer digits of precision than a 32 bits `integer`
- ▶ And a 64 bits `real` has fewer digits of precision than a 64 bits `integer`
- ▶ The result of a conversion could actually be smaller than expected!



- ▶ Do not blindly rely on implementation dependent chance!



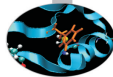
- ▶ Do not blindly rely on implementation dependent chance!
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- ▶ Type conversion functions are not magic
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- ▶ Do not abuse them
  - ▶ Make codes unreadable
  - ▶ Could be evidence of design mistakes
  - ▶ Or that your Fortran needs a refresh





More Flow Control

Fortran Intrinsic Types, Variables and Math

Integer Types

Floating Types

Expressions

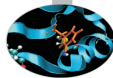
Arithmetic Conversions

**More Intrinsic Types**

Arrays



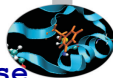
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  - ▶ May have `.false.` or `.true.` value
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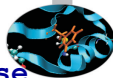
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  - ▶ `==` (equal), `/=` (not equal), `>`, `<`, `>=`, `<=`
  - ▶ or, in ancient Fortran, `.eq.`, `.ne.`, `.gt.`, `.lt.`, `.ge.`, `.le.`



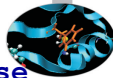
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- ▶ Logical expressions
  - ▶ `.not.` is unary NOT, `.and.` and `.or.` are binary AND and OR respectively, `.eqv.` is logical equivalence (`.true.` if operands both `.true.` or both `.false.`)
  - ▶ `.not. a .and. b .or. a .and. .not. b`  
means  
`((.not.a) .and.b) .or. (a .and. (.not.b))`
  - ▶ In doubt, add parentheses, but be sober



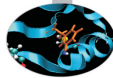
- ▶ Logical friends from `ieee_arithmetic` module (simply use it)
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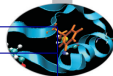


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- ▶ As usual, order of subexpressions evaluation is implementation dependent
- ▶ But it's worse:
  - ▶ if `test()` is a function returning a logical type value
  - ▶ and `a` is `.true.`
  - ▶ and `b` is `.false.`
  - ▶ implementation is free (but not forced!) to not call `test()` at all in `a.or.test(x)` and `b.and.test(x)`
  - ▶ Again, do not rely on expressions side effects



- ▶ Fortran is not that good at manipulating text
- ▶ But it has some **character**:
  - ▶ **character** :: **c** defines a variable holding a single character, like ' **f**'
  - ▶ **character(len=80)** :: **s1, s2, s3** defines three variables holding strings of up to 80 characters, like ' **Fortran 2003**'
- ▶ There are character expressions, like:
  - ▶ **s3(1:40) = s1(1:20)//s2(21:40)**  
which assigns to first half of **s3** the first quarter of **s1** and second quarter of **s2**
- ▶ On assignment of a character expression to a longer variable, blank filling will take place
- ▶ On assignment of a character expression to a shorter variable, truncation will happen





Function	Returns
<code>len(s)</code>	string length
<code>len_trim(s)</code>	string length with trailing blanks ignored
<code>trim(s)</code>	string with trailing blanks removed
<code>repeat(s, n)</code>	string made of <code>n</code> copies of <code>s</code>
<code>adjustl(s)</code>	move leading blanks to trailing position
<code>adjustr(s)</code>	move trailing blanks to leading position
<code>lge(s1, s2),</code> <code>lgt(s1, s2),</code> <code>lle(s1, s2),</code> <code>llt(s1, s2)</code>	string comparisons
<code>index(s, subs)</code>	starting position of <code>subs</code> in <code>s</code> , 0 if not found
<code>scan(s, set)</code>	first position in <code>s</code> of a character matching <code>set</code> , 0 if none found
<code>verify(s, set)</code>	first position in <code>s</code> of a character not matching <code>set</code> , 0 if all match
<code>achar(i)</code>	character with ASCII code <code>i</code>
<code>iachar(c)</code>	ASCII code of character <code>c</code>

- ▶ Our advice:
  - ▶ For most practical purposes, use I/O statements to manipulate strings as internal files (more on this later)
  - ▶ If you are really serious about textual data, learn more
  - ▶ Or switch to a different language



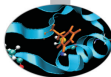
More Flow Control

Fortran Intrinsic Types, Variables and Math

Arrays

Smoothing Signals

A More Compact Notation



More Flow Control

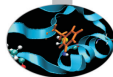
Fortran Intrinsic Types, Variables and Math

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Smoothing Signals

A More Compact Notation

# In Place Smoothing of a Periodic Signal



```

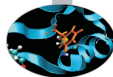
module smoothing
  implicit none
  contains
    subroutine smooth(v, k)
      real,intent(inout) :: v(:)
      integer,intent(in)  :: k
      integer             :: n, l, i, j
      real               :: work(size(v))

      n=size(v)
      l = 2*k +1
      work = 0.0
      do i=1,n
        do j=i-k,i+k
          work(i) = work(i) + v(1+mod(n-1+j, n))
        enddo
      enddo
      v = work/l
    end subroutine smooth
end module smoothing

program test_smooth
  use smoothing
  implicit none
  integer, parameter :: n=10
  integer :: i, k
  real    :: x(n)

  k = 2
  x = (/ (real(mod(i,n/2)), i=1,n) /)
  if ( k > n) stop 'More smoothing points than array elements'
  call smooth(x,k)
  write(*,*) x
end program test_smooth
  
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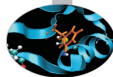
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end program test_smooth
  
```



- ▶ Subroutines are procedures, like functions, except they do not return any value
- ▶ They are invoked by:  
`call subroutine-name (argument-list)`
- ▶ Like functions, they have *dummy* arguments that will be associated to *actual* arguments at call time
- ▶ Unlike functions, they can not be used inside expressions
- ▶ Their use is to be preferred to functions when:
  - ▶ actual arguments must be modified
  - ▶ more than one result needs to be returned

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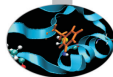
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end program test_smooth
  
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- ▶ `real :: x(n)`
  - ▶ Declares an array named `x`
  - ▶ A collection of variables of the same type (elements), laid out contiguously in memory
  - ▶ `i`-th element can be accessed with `x(i)`
  - ▶ `n` must be an integer expression whose value must be known at declaration time



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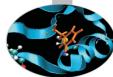
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- ▶ **real :: x(n)**
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  - ▶ **i**-th element can be accessed with **x(i)**
  - ▶ **n** must be an integer expression whose value must be known at declaration time
  
- ▶ What's that **x = (/.../)** ?
  - ▶ **(/.../)** is an array constructor
  - ▶ I.e. a sequence of values forming an array
  - ▶ Assigned to array in a single statement
  - ▶ **(expression, index=initial, final)** evaluates **expression** for each value of **index** as in a do-loop (hence is termed *implied do-loop*)

# In Place Smoothing of a Periodic Signal



```

module smoothing
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      do i=1,n
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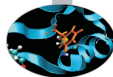
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  write(*,*) x
end program test_smooth
    
```



- ▶ Arrays can be passed as arguments to procedures
- ▶ How can subroutine *smooth* know the size of the actual argument passed as  $\mathbf{v}$ ?
  - ▶ `real :: v(:)` states that size of  $\mathbf{v}$  will be that of the actual argument
  - ▶  $\mathbf{v}$  is termed an *assumed-shape* array
  - ▶ This only works if the subroutine has explicit interface
- ▶ Otherwise, you can still use the good ol' way:

```
subroutine smooth(v,k,n)
  integer n
  real v(n)
  ...
```

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      integer :: n, l, i, j
      real :: work(size(v))

      n=size(v)
      l = 2*k +1
      work = 0.0
      do i=1,n
        do j=i-k,i+k
          work(i) = work(i) + v(1+mod(n-1+j, n))
        enddo
      enddo
      v = work/l
    end subroutine smooth
  end module smoothing

program test_smooth
  use smoothing
  implicit none
  integer, parameter :: n=10
  integer :: i, k
  real :: x(n)

  k = 2
  x = (/ (real(mod(i,n/2)), i=1,n) /)
  if ( k > n) stop 'More smoothing points than array elements'
  call smooth(x,k)
  write(*,*) x
end program test_smooth
  
```



- ▶ Arrays can be passed as arguments to procedures
- ▶ How can subroutine *smooth* know the size of the actual argument passed as  $\mathbf{v}$ ?
  - ▶ **real :: v(:)** states that size of  $\mathbf{v}$  will be that of the actual argument
  - ▶  $\mathbf{v}$  is termed an *assumed-shape* array
  - ▶ This only works if the subroutine has explicit interface

- ▶ Otherwise, you can still use the good ol' way:

```
subroutine smooth(v,k,n)
  integer n
  real v(n)
  ...
```

- ▶ How can subroutine *smooth* declare a local array matching in size the actual argument?
  - ▶ **size(v)** returns the number of elements (size) of  $\mathbf{v}$
  - ▶ **real :: work(size(v))** gives **work** same size as  $\mathbf{v}$
  - ▶ **work** is termed an *automatic object*

## WARNING: NO BOUNDS CHECKING!



- ▶ In Fortran, there is no bounds checking on array access
- ▶ And it is possible for something like this to happen

```
real :: a(10)
...
do i=-100,100
  a(i) = i
end do
```

- ▶ If you are lucky, you'll get a runtime error, otherwise you'll corrupt surrounding memory areas, with really puzzling behavior
- ▶ Once upon a long ago, it used to be a 'feature':

```
subroutine smooth(v,k,n)
  integer n
  real v(1)
  ...
```

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- ▶ Once upon a long ago, it used to be a 'feature':

```
subroutine smooth(v,k,n)
  integer n
  real v(1)
  ...
```
- ▶ Use compiler options to enable runtime detection of out of bounds accesses
  - ▶ But execution is incredibly slowed down
  - ▶ Just a debugging tool, do not use it in production





- ▶ The intrinsic subroutine `cpu_time()` is used to time code regions

```
real :: t1, t2
...
call cpu_time(t1)
...           ! code to be timed
call cpu_time(t2)
write(*,*) 'Execution time for section 1: ', t2-t1, 'seconds'
```



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- ▶ And returns in it processor time consumed by the program in seconds



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- ▶ Takes a default real argument
  - ▶ And returns in it processor time consumed by the program in seconds
- ▶ Use it to measure execution time of `test_smooth` program
- ▶ Can we use less operations to get the same results (within round-off errors)?



More Flow Control

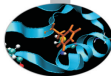
Fortran Intrinsic Types, Variables and Math

**Arrays**

Smoothing Signals

**A More Compact Notation**

# Same Smoothing in a Different Idiom



```
module smoothing
  implicit none
  contains

  subroutine smoothinplace(v, k)
    implicit none
    real,intent(inout) :: v(:)
    integer,intent(in)  :: k
    real                :: work(-k+1:size(v)+k)
    integer :: i, j, l, n

    n=size(v)
    l = 2*k + 1
    work(1:n) = v
    work(-k+1:0) = v(n-k+1:n)
    work(n+1:n+k) = v(1:k)

    do j=1, k
      v = v + work(1-j:n-j) + work(1+j:n+j)
    end do
    v = v/l

  end subroutine smoothinplace

end module smoothing
```

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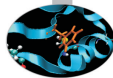
    n=size(v)
    l = 2*k + 1
    work(1:n) = v
    work(-k+1:0) = v(n-k+1:n)
    work(n+1:n+k) = v(1:k)

    do j=1, k
      v = v + work(1-j:n-j) + work(1+j:n+j)
    end do
    v = v/l

  end subroutine smoothinplace

end module smoothing
```

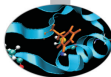




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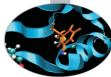
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  - ▶ **work(n+1:n+k)** selects...
- ▶ Arrays and array sections are assigned to by = in a natural manner (more on this later)

# Same Smoothing in a Different Idiom



```
module smoothing
  implicit none
  contains

  subroutine smoothinplace(v, k)
    implicit none
    real,intent(inout) :: v(:)
    integer,intent(in)  :: k
    real                :: work(-k+1:size(v)+k)
    integer :: i, j, l, n

    n=size(v)
    l = 2*k + 1
    work(1:n) = v
    work(-k+1:0) = v(n-k+1:n)
    work(n+1:n+k) = v(1:k)

    do j=1, k
      v = v + work(1-j:n-j) + work(1+j:n+j)
    end do
    v = v/l

  end subroutine smoothinplace

end module smoothing
```



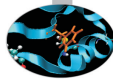
- ▶ Arrays and array sections may be
  - ▶ referenced and used in expressions
  - ▶ passed as arguments to procedures

```
do j=1, k
  v = v + work(1-j:n-j) + work(1+j:n+j)
end do
```

- ▶ Without array expressions, this code would look like:

```
do j=1, k
  do i=1, n
    v(i) = v(i) + work(i-j) + work(i+j)
  end do
end do
```

- ▶ In an array expression, result must not depend in any way on the order of evaluation of elements
- ▶ You should think of array expressions as if all elements were computed at the same time

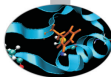


- ▶ The size of a one-dimensional array is its *shape*
- ▶ Arithmetic operators act on arrays element by element
- ▶ Binary operators combine pairs of corresponding elements from the operands
- ▶ With binary operators and assignments, you must use *conformable*, i.e. identically shaped, arrays
- ▶ Except for scalar values (not variables!), that match any shape, as if they were replicated

```

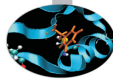
real, dimension(4) :: u, v, w
real :: t(1), s
t = s ! it's right
s = t ! it's wrong
w = (u-v)**2 ! it's right
w = s*u+v+2.3 ! it's OK
w = u+v(1:2) ! it's wrong
  
```

- ▶ By the way, **dimension** attribute lets you specify bounds and dimensions for a list of identical arrays

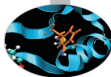


- ▶ Intrinsic subroutine **random\_number (x)** returns pseudo-random numbers uniformly distributed in  $[0, 1)$  interval
  - ▶ Takes an argument of type **real**, that can be either a scalar or an array
  - ▶ Returns one random number if **x** is a scalar
  - ▶ Returns an array of random numbers if **x** is an array
  
- ▶ Is **random\_number ()** as uniform as advertised? Let's check...

# Let's Build An Histogram

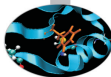


- ▶ Write a program that:
  1. reads an integer **niter** from standard input
  2. generates **niter** random numbers in interval  $[0, 10)$
  3. builds an histogram and computes their average
  4. Prints out results
  
- ▶ To build the histogram:
  1. Initialize to 0s an array **hist** of 20 **integers** to hold the bin count, then, at each iteration:
  2. generate a random number
  3. find out the bin it belongs to (i.e. its index in the array **hist**)
  4. intrinsic **ceiling(x)** function helps: it returns  $\lceil x \rceil$
  5. increment the corresponding array element and compute the percentages
  6. accumulate the sum of the random numbers to compute the average value



- ▶ A prime number is a natural number which has only two distinct natural divisors: 1 and itself
- ▶ Find all primes less than or equal to a given  $n$  by Eratosthenes' algorithm:
  1. create a list of consecutive integers from 2 to  $n$
  2. let be  $p \leftarrow 2$  the first prime
  3. strike from the list all multiples of  $p$  up to  $n$
  4. let  $p \leftarrow$  next number still in the list after  $p$
  5. if  $2p < n$ , get back to step 3
  6. all remaining numbers in the list are primes

Try it now!

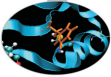


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Try it now!

- ▶ How could you spare iterations?
- ▶ How could you spare memory?





## Part III

# Array Syntax and I/O

Multidimensional arrays, array-syntax, array-value function, temporary array, shape, data, reshape, constant array and elemental procedure. Coconstructs where, forall, array reduction. Advanced I/O: formats and descriptors, I/O to/from file, namelist, internal file, unformatted I/O, positioning instructions, stream access. Managing errors.



## Array Syntax

More dimensions

Not a Panacea

Arrays of Constants

Elemental Procedures

More Array Syntax

Input/Output



## Array Syntax

**More dimensions**

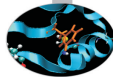
Not a Panacea

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More Array Syntax

Input/Output



```
function avgk(v, k)

    implicit none

    real,intent(in) :: v(:, :)
    integer,intent(in) :: k
    real :: avgk(size(v,1)/k,size(v,2)/k)

    integer :: i, j, n, m

    n = (size(v,1)/k)*k
    m = (size(v,2)/k)*k

    avgk = 0.0

    do j=1, k
        do i=1, k
            avgk = avgk + v(i:n:k,j:m:k)
        end do
    end do

    avgk = avgk/k**2

end function avgk
```



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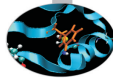
    n = (size(v,1)/k)*k
    m = (size(v,2)/k)*k

    avgk = 0.0

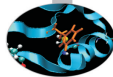
    do j=1, k
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            avgk = avgk + v(i:n:k,j:m:k)
        end do
    end do

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end function avgk
```



- ▶ Arrays may have up to 7 dimensions
- ▶ Lower bounds default to 1, but you can specify them as for one-dimensional arrays, like in `q(-k:k, 11:20)`
- ▶ Elements are referenced by a list of indices: `v(1, 1)`
- ▶ The sequence of extents of an array is termed its *shape*, e.g. if `a` is `real :: a(3, 2:5)` then:
  - ▶ `shape(a)` returns the array of extents `(/3, 4/)`
  - ▶ whereas `size(a)` returns `12`
- ▶ Multidimensional (i.e. `rank>1`) arrays and array sections may be involved in array expressions
- ▶ As in the case of rank 1 arrays, they must be conformable when needed:  
`avgk(1:3, :) = avgk(5:9, :)` is wrong



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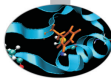
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- ▶ Yes, a function may return an array
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- ▶ **`size(array, dim)`** returns the integer extent of **`array`** along dimension **`dim`**
- ▶ Number of dimensions (a.k.a. rank) is mandatory in assumed shape arrays



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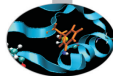
end function avgk
```

# Pay Attention to Conformability



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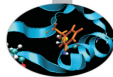
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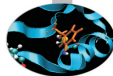


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- ▶ what if  $\mathbf{v}$  extents aren't multiple of  $\mathbf{k}$ ?
  - ▶  $\mathbf{v}(\mathbf{i}:\mathbf{n}:\mathbf{k}, \mathbf{j}:\mathbf{m}:\mathbf{k})$  and  $\mathbf{avgk}$  would not be conformable

# Pay Attention to Conformability



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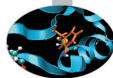
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► Good ol' style:

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do i=1,n
  x(i) = b(i) / a(i,i)
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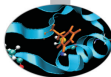
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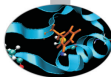
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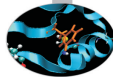
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▶ What happens for  $i==n$ ?

- ▶ the array section  $b(n+1:n)$  has zero size:  
lower bound > upper bound
- ▶ No operation is performed

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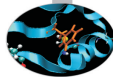
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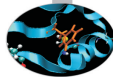
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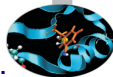
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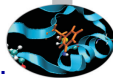
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# A Closer Look To Array Expressions

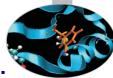


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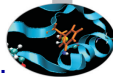
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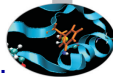


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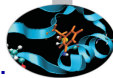


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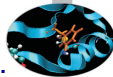


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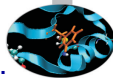
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- ▶ Array syntax can be very compact and elegant
- ▶ But temporary copies may impact performance, use your compiler options to spot them



## Array Syntax

More dimensions

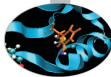
**Not a Panacea**

Arrays of Constants

Elemental Procedures

More Array Syntax

Input/Output



```
function trace(matrix)
implicit none
real, intent(in) :: matrix(:, :)
real :: trace
integer :: i
integer :: dim(2)

dim = shape(matrix)
trace = 0.0
if (dim(1) /= dim(2)) return

do i=1,dim(1)
    trace = trace + matrix(i,i)
enddo
end function trace
```

- ▶ Not all operations on arrays can easily be expressed in array syntax
- ▶ Do you remember **shape ()** ? It returns an array whose elements are the extents of its argument



```
subroutine smooth(v, k)
  implicit none
  real,intent(inout) :: v(:)
  integer,intent(in) :: k
  integer :: n, l, i, j
  real :: work(size(v))

  n=size(v)
  l = 2*k +1
  work(1) = 0.0
  do j=1-k,1+k
    work(1) = work(1) + v(1+mod(n-1+j, n))
  enddo
  do i=2,n
    work(i)=work(i-1)+v(1+mod(n-1+i+k, n))-v(1+mod(n-2+i-k, n))
  enddo
  v = work/l
end subroutine smooth
```

- ▶ The above code does the smoothing with minimal operations count
- ▶ And cannot be expressed at all in array syntax
- ▶ This is a quite common situation: optimal algorithms operating on arrays often sports dependencies in elements evaluations and updates





## Array Syntax

More dimensions

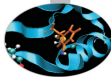
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! Polynomial approximation of  $J_0(x)$  for  $-3 \leq x \leq 3$   
 ! See Abramowitz&Stegun for details

```

function j0(x)
  implicit none
  real :: j0
  real, intent(in) :: x

  integer, parameter :: order = 6
  real, parameter, dimension(0:order) :: coeff = &
    (/ 1.0000000, &
      -2.2499997, &
       1.2656208, &
      -0.3163866, &
       0.0444479, &
      -0.0039444, &
       0.0002100 /)
  real :: xo3sq
  integer :: i

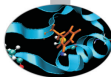
  xo3sq = (x/3.0)**2
  j0 = coeff(order)

  ! horner method
  do i=order, 1, -1
    j0 = j0*xo3sq + coeff(i-1)
  end do
end function j0
  
```



- ▶ **parameter** arrays are very good at storing tables of:
  - ▶ polynomial coefficients
  - ▶ physical measurements
  - ▶ function values at discrete points
  
- ▶ In the past, **data** statements were used:  

```
data coeff /1.0,-2.2499997,1.2656208,-0.3163866, &  
          0.0444479,-0.0039444,0.0002100/
```
  
- ▶ **data** statements:
  - ▶ are very versatile
  - ▶ very difficult to decipher
  - ▶ and tend to float away from variable declaration
  
- ▶ Use initialization instead



## Array Syntax

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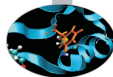
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Arrays of Constants

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```
program array_swap

  implicit none
  integer :: i, j
  real    :: a(0:10,10), b(11,10)

  a=reshape( (/ (i*0.1, i=1,110) /), (/11,10/) )
  b=reshape( (/ ((i*j+i, i=1,11), j=1,10) /), (/11,10/) )
  call swap(a,b)

end program array_swap

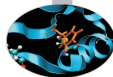
subroutine swap(a,b)

  implicit none

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  real, dimension(size(a,1),size(a,2)) :: tmp

  tmp = a
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end subroutine swap
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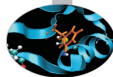
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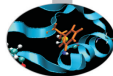
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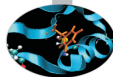
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  - ▶ Other variables with same names are unaffected
- ▶ **reshape** (*source*, *new\_shape*) returns an array with shape given by the rank one integer array *new\_shape*, and elements taken from *source* in array element order



```

program array_swap

  implicit none
  integer :: i, j
  real    :: a(0:10,10), b(11,10)

  interface
    subroutine swap(a,b)
      real, intent(inout)      :: a(:, :), b(:, :)
      real, dimension(size(a,1),size(a,2)) :: tmp
    end subroutine swap
  end interface

  a=reshape( (/ (i*0.1, i=1,110) /), (/11,10/) )
  b=reshape( (/ ((i*j+i, i=1,11), j=1,10) /), (/11,10/) )
  call swap(a,b)

end program array_swap

subroutine swap(a,b)

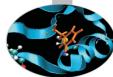
  implicit none

  real, intent(inout) :: a(:, :),b(:, :)
  real, dimension(size(a,1),size(a,2)) :: tmp

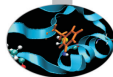
  tmp = a
  a = b
  b = tmp

end subroutine swap

```



- ▶ The scope of the implied do loop indices  $i$  and  $j$  is the loop itself
  - ▶ Other variables with same names are unaffected
- ▶ **reshape** (*source*, *new\_shape*) returns an array with shape given by the rank one integer array *new\_shape*, and elements taken from *source* in array element order
- ▶ Interface is as always mandatory for assumed shape arguments, so the compiler knows that additional information must be passed in to the function



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- ▶ Interface is as always mandatory for assumed shape arguments, so the compiler knows that additional information must be passed in to the function
- ▶ But life can be simpler...



```
program array_swap
  implicit none
  integer :: i, j
  real    :: a(0:10,10), b(11,10)

  interface
    elemental subroutine swap(a,b)
      real, intent(inout) :: a, b
      real                :: tmp
    end subroutine swap
  end interface

  a = reshape( (/ (i*0.1, i=1,110) /), (/11,10/) )
  b = reshape( (/ ((i*j+i, i=1,11), j=1,10) /), (/11,10/) )

  call swap(a,b)

end program array_swap

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  implicit none

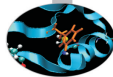
  real, intent(inout) :: a, b
  real                :: tmp

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- ▶ To define one, it has to be pure
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  - ▶ If a subroutine, side effects shall be restricted to **intent (out)** and **intent (inout)** arguments
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  - ▶ And some more constraints ensure the different procedure calls can be safely executed in any order
- ▶ An explicit interface is mandatory
  - ▶ It must specify the procedure as **elemental**
  - ▶ It must specify **intent ()** attribute for all arguments





## Array Syntax

More dimensions

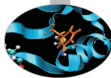
Not a Panacea

Arrays of Constants

Elemental Procedures

**More Array Syntax**

Input/Output



- ▶ Logical array expressions like `a (: ) > 0 . 0` are often termed *masks*

## Masks and **where**



- ▶ Logical array expressions like  $\mathbf{a}(:,) > 0.0$  are often termed *masks*
- ▶ They come useful to restrict computations to specific array elements, as in the **where** statement:  
**where (abs(a) > abs(b)) a = b**  
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- ▶ The general form is the **where** construct

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elsewhere  
    c=a  
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- ▶ **where** constructs can be nested and given a name

# Say it With `forall`s



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- ▶ In its construct form, it looks like:

```
forall(i = 2:n-1, j = 2:n-1)  
  a(i,j) = a(i,j-1) + a(i,j+1) + a(i-1,j) + a(i+1,j)  
  b(i,j) = a(i,j)  
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- ▶ `forall` constructs can be nested and given a name



- Using `do` loops (dependencies! loop order is crucial)

```
do j=2,n-1
  do i=2,n-1
    T(i,j) = ( T(i-1,j) + T(i+1,j) + &
              T(i,j-1) + T(i,j+1) )/4.0
  enddo
enddo
```



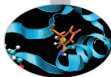
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  enddo
enddo
```

- ▶ Using array syntax (compiler enforces correct semantics)

```
T(2:n-1,2:n-1) = ( T(1:n-2,2:n-1) + T(3:n,2:n-1) &
                  + T(2:n-1,1:n-2) + T(2:n-1,3:n) )/4.0
```





- ▶ Using **do** loops (dependencies! loop order is crucial)

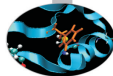
```
do j=2,n-1
  do i=2,n-1
    T(i,j) = ( T(i-1,j) + T(i+1,j) + &
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  enddo
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```

- ▶ Using **forall** (ditto, but more readable)

```
forall (i=2:n-1, j=2:n-1)
  T(i,j) = ( T(i-1,j) + T(i+1,j) + &
            T(i,j-1) + T(i,j+1) )/4.0
end forall
```



```

integer, parameter :: maxn=768, maxm=939, R=3
real, parameter :: sd=10.0, sr=10.0
real, parameter :: sd22=2.0*sd**2, sr22=2.0*sr**2
integer :: i, j, m, n
real :: B(maxn,maxm), A(maxn,maxm)
real :: zsum, z
real, dimension(-R:R,-R:R), parameter :: z0=&
  reshape((/ ((exp(-(m**2 + n**2)/sr22), m=-R, R), n=-R,R) /), (/ 2*R+1, 2*R+1 /))
...
B = 0.0
do i=1,maxn
  do j=1,maxm
    zsum = 0.0
    do m= -R,R
      if (i+m >= 1 .and. i+m <= maxn) then
        do n = -R,R
          if(j+n >= 1 .and. j+n <= maxm) then
            z = exp(-(A(i+m, j+n)-A(i, j))**2/sd22)*z0(m,n)
            zsum = zsum + z
            B(i, j) = B(i, j) + z*A(i+m, j+n)
          end if
        end do
      end if
    end do
    B(i, j) = B(i, j) / zsum
  end do
end do

```

# Bilateral Filter Using forall



```

integer, parameter :: maxn=768, maxm=939, R=3
real, parameter :: sd=10.0, sr=10.0
real, parameter :: sd22=2.0*sd**2, sr22=2.0*sr**2
integer :: i, j, m, n
real :: B(maxn,maxm), A(maxn,maxm)
real :: z(-R:R,-R:R), aw(-R:R,-R:R)
real, dimension(-R:R,-R:R), parameter :: z0=&
  reshape(/ ((exp(-(m**2 + n**2)/sr22), m=-R, R), n=-R,R) /), (/ 2*R+1, 2*R+1 /))
...
do i=1,maxn      ! These two cannot be changed into forall
  do j=1,maxm    ! Why?
    z = 0.0
    forall (m=max(1,i-R):min(maxn,i+R))
      forall (n=max(1,j-R):min(maxm,j+R))
        aw(m-i,n-j) = A(m,n)
        z(m-i,n-j) = exp(-(aw(m-i,n-j)-A(i,j))**2/sd22)*z0(m-i,n-j)
      end forall
    end forall
    B(i,j) = sum(z*aw)/sum(z)
  end do
end do
  
```

# Bilateral Filter Using forall

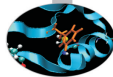


```

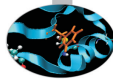
integer, parameter :: maxn=768, maxm=939, R=3
real, parameter :: sd=10.0, sr=10.0
real, parameter :: sd22=2.0*sd**2, sr22=2.0*sr**2
integer :: i, j, m, n
real :: B(maxn,maxm), A(maxn,maxm)
real :: z(-R:R,-R:R), aw(-R:R,-R:R)
real, dimension(-R:R,-R:R), parameter :: z0=&
  reshape(/ ((exp(-(m**2 + n**2)/sr22), m=-R, R), n=-R, R) /), (/ 2*R+1, 2*R+1 /))
...
do i=1,maxn      ! These two cannot be changed into forall
  do j=1,maxm    ! Why?
    z = 0.0      ! Because this happens at every iteration, it's a dependency!
    forall (m=max(1,i-R):min(maxn,i+R))
      forall (n=max(1,j-R):min(maxm,j+R))
        aw(m-i,n-j) = A(m,n)
        z(m-i,n-j) = exp(-(aw(m-i,n-j)-A(i,j))**2/sd22)*z0(m-i,n-j)
      end forall
    end forall
    B(i,j) = sum(z*aw)/sum(z)
  end do
end do
  
```



- ▶ Reductions squeeze an array to a scalar
  - ▶ **all** (*mask*) returns true if all the elements of mask are true
  - ▶ **any** (*mask*) returns true if any of the elements of mask are true
  - ▶ **count** (*mask*) returns the number of **.true.** elements in mask
  - ▶ **maxval** (*array*) returns the maximum value of array
  - ▶ **minval** (*array*) returns the minimum value of array
  - ▶ **sum** (*array*) returns the sum of the elements of array
  - ▶ **product** (*array*) returns the product of the elements of array
- ▶ Or to an array of rank reduced by one, if you specify an optional dimension to perform reduction along, like in **sum(a(:, :, :), dim=2)**



- ▶ More functions, good to know:
  - ▶ **maxloc()** and **minloc()** return locations of maximum and minimum value respectively
  - ▶ **cshift()** performs a *circular* shift along an array dimension
  - ▶ **eoshift()** perform a *end-off* shift along an array dimension
  - ▶ **spread()** increases by one the rank of an array expression
  - ▶ **pack()** selects elements from an array according to a mask and packs them in a rank-1 array
  - ▶ And **unpack()** does the reverse
- ▶ But too much detail to cover in this introduction, look for them on your compiler documentation, and experiment



- ▶ Vector and matrix multiplication functions
  - ▶ `dot_product(vector_a, vector_b)`
  - ▶ `matmul(matrix_a, matrix_b)`
- ▶ But the BLAS libraries are around
  - ▶ Widely used
  - ▶ Highly optimized implementations available
- ▶ Outstanding compilers include special purpose, optimized BLAS version for those calls
- ▶ Good compilers do not include BLAS, but give option to link them for those calls
- ▶ Average compilers do not shine for those calls
- ▶ Our advice: install a reputedly good BLAS version and use it
- ▶ There is more to matrix algebra than matrix multiplies and vector products



- ▶ Re-write the Sieve of Eratosthenes algorithm using array syntax





## Array Syntax

## Input/Output

- Formatted I/O

- File I/O

- Namelist

- Internal Files

- Unformatted I/O

- Robust I/O



Array Syntax

Input/Output

**Formatted I/O**

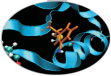
File I/O

Namelist

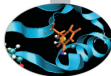
Internal Files

Unformatted I/O

Robust I/O



- ▶ Data are manipulated in internal (usually binary) format
- ▶ Fortran Standard leaves internal format details up to the implementation
- ▶ Formatted I/O translates internal representation of variables into human readable format
- ▶ Best practices:
  - ▶ Use formatted I/O just for small amount of data meant to be read by humans
  - ▶ Beware: human readable representation may cause problems because of rounding or not enough digits
  - ▶ Do not use I/O inside heavy computations: inhibits some code optimizations, and significantly affects performance



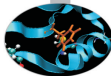
```
program golden_ratio
! experiments with the golden ratio iterative relation
implicit none
integer, parameter :: rk = kind(1.0d0)
real(rk) :: phi, phi_old
real(rk) :: phi_start, tol
integer :: i, max_iter

write(*,*) 'Enter start value, tol, max iterations'
read(*,*) phi_start, tol, max_iter

phi_old = phi_start
do i=1,max_iter
  phi = 1.0d0/phi_old + 1.0d0
  if (abs(phi - phi_old) < tol) exit
  phi_old = phi
end do

write(*,100) 'Start value:',phi_start
write(*,100) 'Tolerance:',tol
write(*,'(2(A," ",I11," "))') 'Ended at iteration:', i, 'of', max_iter
write(*,100) 'Final value:',phi

100 format(A," ",F13.10)
end program golden_ratio
```



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! experiments with the golden ratio iterative relation
implicit none
integer, parameter :: rk = kind(1.0d0)
real(rk) :: phi, phi_old
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write(*,*) 'Enter start value, tol, max iterations'
read(*,*) phi_start, tol, max_iter

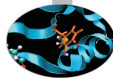
phi_old = phi_start
do i=1,max_iter
  phi = 1.0d0/phi_old + 1.0d0
  if (abs(phi - phi_old) < tol) exit
  phi_old = phi
end do

write(*,100) 'Start value:',phi_start
write(*,100) 'Tolerance:',tol
write(*,'(2(A," ",I11," "))') 'Ended at iteration:', i, 'of', max_iter
write(*,100) 'Final value:',phi

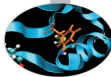
100 format(A," ",F13.10)
end program golden_ratio
```



- ▶ The easiest way to do formatted I/O
- ▶ Specified using `*`
- ▶ Values are translated according to their types
- ▶ In the order they are listed on I/O statements
- ▶ No-nonsense, implementation dependent format
- ▶ Often outputs more digits than you actually care of
- ▶ Best practices:
  - ▶ Use it for terminal input
  - ▶ Use it for input of white-space separated values
  - ▶ Use it for quick output
  - ▶ Not suitable for rigid tabular formats



- ▶ Put you in total control of what is read/written



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! experiments with the golden ratio iterative relation
  implicit none
  integer, parameter :: rk = kind(1.0d0)
  real(rk) :: phi, phi_old
  real(rk) :: phi_start, tol
  integer :: i, max_iter

  write(*,*) 'Enter start value, tol, max iterations'
  read(*,*) phi_start, tol, max_iter

  phi_old = phi_start
  do i=1,max_iter
    phi = 1.0d0/phi_old + 1.0d0
    if (abs(phi - phi_old) < tol) exit
    phi_old = phi
  end do

  write(*,100) 'Start value:',phi_start
  write(*,100) 'Tolerance:',tol
  write(*,'(2(A," ",I11," "))') 'Ended at iteration:', i, 'of', max_iter
  write(*,100) 'Final value:',phi

100 format(A," ",F13.10)
end program golden_ratio
  
```



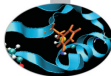
## Explicit formats



- ▶ Put you in total control of what is read/written
- ▶ Specified by (*format-list*)



- ▶ Put you in total control of what is read/written
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- ▶ Where *format-list* is a comma separated list of items, which can be:
  - ▶ string literals, usually in double quotes, emitted *as-is*
  - ▶ or proper *edit descriptors*, which dictate how a corresponding element on the I/O list should be converted



```

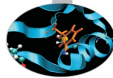
program golden_ratio
! experiments with the golden ratio iterative relation
  implicit none
  integer, parameter :: rk = kind(1.0d0)
  real(rk) :: phi, phi_old
  real(rk) :: phi_start, tol
  integer :: i, max_iter

  write(*,*) 'Enter start value, tol, max iterations'
  read(*,*) phi_start, tol, max_iter

  phi_old = phi_start
  do i=1,max_iter
    phi = 1.0d0/phi_old + 1.0d0
    if (abs(phi - phi_old) < tol) exit
    phi_old = phi
  end do

  write(*,100) 'Start value:',phi_start
  write(*,100) 'Tolerance:',tol
  write(*,'(2(A," ",I11," "))') 'Ended at iteration:', i, 'of', max_iter
  write(*,100) 'Final value:',phi

100 format(A," ",F13.10)
end program golden_ratio
  
```



- ▶ Put you in total control of what is read/written
- ▶ Specified by (*format-list*)
- ▶ Where *format-list* is a comma separated list of items, which can be:
  - ▶ string literals, usually in double quotes, emitted *as-is*
  - ▶ or proper *edit descriptors*, which dictate how a corresponding element on the I/O list should be converted
- ▶ *Repeat counts* can be used
  - ▶ Like in `5I3`, which will convert 5 **integer** values
  - ▶ Like in `2(I3,F7.4)`, which will convert 2 pairs, each made of an **integer** and a **real** value



```

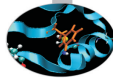
program golden_ratio
! experiments with the golden ratio iterative relation
  implicit none
  integer, parameter :: rk = kind(1.0d0)
  real(rk) :: phi, phi_old
  real(rk) :: phi_start, tol
  integer :: i, max_iter

  write(*,*) 'Enter start value, tol, max iterations'
  read(*,*) phi_start, tol, max_iter

  phi_old = phi_start
  do i=1,max_iter
    phi = 1.0d0/phi_old + 1.0d0
    if (abs(phi - phi_old) < tol) exit
    phi_old = phi
  end do

  write(*,100) 'Start value:',phi_start
  write(*,100) 'Tolerance:',tol
  write(*,'(2(A," ",I11," "))') 'Ended at iteration:', i, 'of', max_iter
  write(*,100) 'Final value:',phi

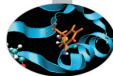
100 format(A," ",F13.10)
end program golden_ratio
  
```



- ▶ Put you in total control of what is read/written
- ▶ Specified by (*format-list*)
- ▶ Where *format-list* is a comma separated list of items, which can be:
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- ▶ *Repeat counts* can be used
  - ▶ Like in `5I3`, which will convert 5 **integer** values
  - ▶ Like in `2(I3,F7.4)`, which will convert 2 pairs, each made of an **integer** and a **real** value
- ▶ Formats must be specified on I/O statements
  - ▶ As a literal string, usually in single quotes
  - ▶ As a character expression
  - ▶ As a numeric label of a **format** statement in the same program unit (traditionally, before its end), reusable in many statements



- ▶ **A** is used to translate **character** values
  - ▶ **A** will emit the value as is
  - ▶ **A10** will emit 10 characters, truncating the value if longer, right justifying it if shorter
  - ▶ Beware: leading white-space skipped on input
  - ▶ Beware: **A10** and **10A** mean very different things!



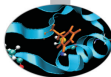
- ▶ **A** is used to translate **character** values
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  - ▶ **A10** will emit 10 characters, truncating the value if longer, right justifying it if shorter
  - ▶ Beware: leading white-space skipped on input
  - ▶ Beware: **A10** and **10A** mean very different things!
  
- ▶ **I** is used to translate **integer** values
  - ▶ **I6** will emit up to 6 characters (sign included!), right justified with blanks
  - ▶ **I6.3** will emit 6 characters (sign included!), containing at least 3 (possibly zero) digits, right justified with blanks
  - ▶ Beware: again, **I10** and **10I** mean very different things!





- ▶ `F` can be used to translate `real` values
  - ▶ `F8.3` will emit up to 8 characters (sign and decimal point included!) in total, with 3 decimal digits (possibly zero), right justified with blanks
  - ▶ Beware: if `F6.2` is specified in input, and `-12345` is met, the value `-123.45` will be read in!
  - ▶ Beware: if `F6.2` is specified in input, and `-1.234` is met, the value `-1.234` will be read in anyhow!
- ▶ Beware of rounding: internal representation could have more precision than specified in edit descriptors

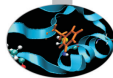
# More Edit Descriptors for `real`s



- ▶ **E** (or **D**) can also be used to translate `real` values
  - ▶ Exponential form is used (mantissa in the  $[0,1)$  range)
  - ▶ Values  $|x| < 10^{99}$ , as  $-1.5372 \times 10^{98}$ , will be converted like:  
**-.15372E+99**
  - ▶ Values  $|x| \geq 10^{99}$ , as  $-1.5372 \times 10^{99}$ , will be converted like:  
**-.15372+100**
  - ▶ **E15.7** will emit up to 15 characters (sign, decimal point, and exponent field included!), with 7 decimal mantissa digits (possibly zero), right justified with blanks
  - ▶ Ditto for **E15.7E4**, except that 4 digits will be used for exponent
  - ▶ Again, input is more liberal



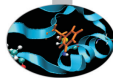
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**-.15372+100**
  - ▶ **E15.7** will emit up to 15 characters (sign, decimal point, and exponent field included!), with 7 decimal mantissa digits (possibly zero), right justified with blanks
  - ▶ Ditto for **E15.7E4**, except that 4 digits will be used for exponent
  - ▶ Again, input is more liberal
- ▶ And more can be used to the same purpose
  - ▶ Like **EN** (engineering notation), same as **E**, with exponent always multiple of 3
  - ▶ Like **G**, which uses the most suitable between **F** and **E**, depending on the value magnitude



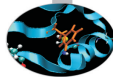
- ▶ /
  - ▶ Forces a new line on output
  - ▶ Skips to next line on input



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  - ▶ Skips to next line on input
  
- ▶ Leading sign of numeric values
  - ▶ **SP** forces following numeric conversions to emit a leading + character for positive values
  - ▶ **SS** restores the default (sign is suppressed for positive values)



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  - ▶ **SS** restores the default (sign is suppressed for positive values)
  
- ▶ Embedded blanks in numeric input fields
  - ▶ **BZ** forces embedded blanks to be treated as 0 digits
  - ▶ **BN** restores the default (blanks are skipped)

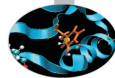


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- ▶ Embedded blanks in numeric input fields
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  - ▶ **BN** restores the default (blanks are skipped)
- ▶ And more... browse your compiler manuals



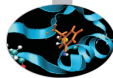
- ▶ **complex** values are made of two reals
  - ▶ Thus two edit descriptors must be provided
  - ▶ First one for real part, second one for imaginary part





- ▶ **complex** values are made of two reals
  - ▶ Thus two edit descriptors must be provided
  - ▶ First one for real part, second one for imaginary part
- ▶ Arrays are indexed collections of elements
  - ▶ Thus a proper edit descriptor must be provided for each element
  - ▶ And if elements are of **complex**, or derived types, see above

- ▶ What if more characters than needed are present on an input line?

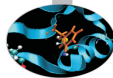


# Fortran I/O is Robustly Designed



- ▶ What if more characters than needed are present on an input line?
  - ▶ After `read`, remaining ones are ignored up to end of line

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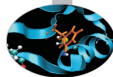
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  - ▶ Following edit descriptors are ignored
- ▶ What if the list of edit descriptors in a format is exhausted before end of items to read/write?

# Iterative Matrix Inversion



```

program iterative_inversion
! experiments with matrix iterative inversion
implicit none
real, dimension(4,4) :: a, x, x_old, x_start
real :: tol, err
integer :: i, max_iter

write(*,*) 'Enter 4x4 matrix to invert'
read(*,*) a
write(*,*) 'Enter 4x4 start matrix'
read(*,*) x_start
write(*,*) 'Enter tol, max iterations'
read(*,*) tol, max_iter

x_old = x_start
do i=1,max_iter
  x = 2.0*x_old - matmul(x_old,matmul(a,x_old))
  err = maxval(abs(x - x_old))
  if (err < tol) exit
  x_old = x
end do

write(*,('("Matrix to invert:")'))
write(*,100) a
write(*,('/',"Start matrix:")')
write(*,100) x_start
write(*,('/',A," ",E15.7')) 'Tolerance:',tol
write(*,('/',2(A," ",I11," "))') 'Ended at iteration:', i, 'of', max_iter
write(*,('("Final matrix:")'))
write(*,100) x

100 format(4(E15.7," "))
end program iterative_inversion
  
```



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  - ▶ Easy answer: I/O continues on a new line, reapplying the format list from its beginning, quite handy for arrays





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  - ▶ Could be more complex, look for *reversion* to know more

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  - ▶ The field is filled with asterisks (i.e. `*`)
- ▶ What if a type mismatch happens between an item to read/write and its corresponding edit descriptor?
  - ▶ Your fault, you are in for a runtime, implementation defined surprise!



- ▶ Play with `golden.f90` and `itinv.f90`:
  - ▶ trying good and bad inputs
  - ▶ giving less or more inputs than needed
  - ▶ changing format descriptors



Array Syntax

Input/Output

Formatted I/O

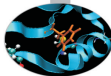
**File I/O**

Namelist

Internal Files

Unformatted I/O

Robust I/O



```
program golden_ratio
! experiments with the golden ratio iterative relation
implicit none
integer, parameter :: rk = kind(1.0d0)
real(rk) :: phi, phi_old
real(rk) :: phi_start, tol
integer :: i, max_iter

open(11,FILE='golden.in',STATUS='old')
read(11,*) phi_start, tol, max_iter
close(11)

phi_old = phi_start
do i=1,max_iter
  phi = 1.0d0/phi_old + 1.0d0
  if (abs(phi - phi_old) < tol) exit
  phi_old = phi
end do

open(12,FILE='golden.out')
write(12,100) 'Start value:',phi_start
write(12,100) 'Tolerance:',tol
write(12,'(2(A," ",I11," "))') 'Ended at iteration:', i, 'of', max_iter
write(12,100) 'Final value:',phi
close(12)

100 format(A," ",F13.10)
end program golden_ratio
```





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! experiments with the golden ratio iterative relation
implicit none
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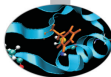
open(11,FILE='golden.in',STATUS='old')
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close(11)

phi_old = phi_start
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write(12,100) 'Final value:',phi
close(12)

100 format(A," ",F13.10)
end program golden_ratio
```

## opening a File for I/O



```
open (u, FILE=file_name[, option][, option][...])
```

- ▶ *u* is an integer, positive expression specifying a *file handle*

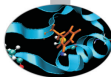
## opening a File for I/O



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- ▶ ***file\_name*** is a string specifying file name (and possibly path) in your file system

## opening a File for I/O



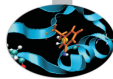
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  - ▶ When you pass a **\*** instead, you are using pre-opened units mapping to user terminal



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- ▶ *file handle* is then used as first argument to **read** and **write**
  - ▶ When you pass a **\*** instead, you are using pre-opened units mapping to user terminal
  - ▶ Which usually means **5** for **read** and **6** for **write**, but **\***, or **input\_unit** and **output\_unit** from **iso\_fortran\_env** Fortran 2003 module are more portable



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```

- ▶ ***u*** is an integer, positive expression specifying a *file handle*
- ▶ ***file\_name*** is a string specifying file name (and possibly path) in your file system
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  - ▶ When you pass a **\*** instead, you are using pre-opened units mapping to user terminal
  - ▶ Which usually means **5** for **read** and **6** for **write**, but **\***, or **input\_unit** and **output\_unit** from **iso\_fortran\_env** Fortran 2003 module are more portable
  - ▶ For error messages, **0** is commonly used, but **error\_unit** from **iso\_fortran\_env** module is portable

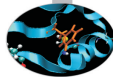


- ▶ **ACTION=act** specifies allowed actions
  - ▶ use '**read**' to only read
  - ▶ use '**write**' to only write
  - ▶ use '**readwrite**' (the default) to allow both





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- ▶ **STATUS=st** tells how to behave wrt. file existence:
  - ▶ use '**old**' to open a file that must already exist
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  - ▶ use '**replace**' to open a new file, even if one already exists
  - ▶ use '**unknown**' (the default) to leave it up to the implementation (in all cases we know of, this means '**replace**')



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  - ▶ use '**replace**' to open a new file, even if one already exists
  - ▶ use '**unknown**' (the default) to leave it up to the implementation (in all cases we know of, this means '**replace**')
- ▶ **POSITION=pos** tells where to start I/O on an existing file
  - ▶ use '**rewind**' (the default) to start at beginning of file
  - ▶ use '**append**' to start at end of file



```
program golden_ratio
! experiments with the golden ratio iterative relation
implicit none
integer, parameter :: rk = kind(1.0d0)
real(rk) :: phi, phi_old
real(rk) :: phi_start, tol
integer :: i, max_iter

open(11,FILE='golden.in',STATUS='old')
read(11,*) phi_start, tol, max_iter
close(11)

phi_old = phi_start
do i=1,max_iter
  phi = 1.0d0/phi_old + 1.0d0
  if (abs(phi - phi_old) < tol) exit
  phi_old = phi
end do

open(12,FILE='golden.out')
write(12,100) 'Start value:',phi_start
write(12,100) 'Tolerance:',tol
write(12,'(2(A," ",I11," "))') 'Ended at iteration:', i, 'of', max_iter
write(12,100) 'Final value:',phi
close(12)

100 format(A," ",F13.10)
end program golden_ratio
```



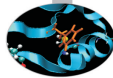
`close (u[, STATUS=st])`

- ▶ `close` completes all pending I/O operations and disassociates the file from the unit



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- ▶ `close` completes all pending I/O operations and disassociates the file from the unit
- ▶ `close` is automatically executed on all open files at program end, but closing a file explicitly when you are done with it is a good practice
- ▶ `st` tells what to do with the file after closing it
  - ▶ use '`keep`' to preserve the file (it's the default)
  - ▶ use '`delete`' to remove it (good for files used for temporary storage)



Array Syntax

Input/Output

Formatted I/O

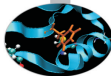
File I/O

**Namelist**

Internal Files

Unformatted I/O

Robust I/O



```
program golden_ratio
! experiments with the golden ratio iterative relation
implicit none
integer, parameter :: rk = kind(1.0d0)
real(rk) :: phi, phi_old
real(rk) :: phi_start, tol
integer :: i, max_iter

namelist /golden_inputs/ phi_start, tol, max_iter

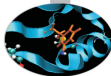
open(11,FILE='golden.in',STATUS='old')
read(11,golden_inputs)
close(11)

phi_old = phi_start
do i=1,max_iter
  phi = 1.0d0/phi_old + 1.0d0
  if (abs(phi - phi_old) < tol) exit
  phi_old = phi
end do

open(12,FILE='golden.out')
write(12,100) 'Start value:',phi_start
write(12,100) 'Tolerance:',tol
write(12,'(2(A," ",I11," "))') 'Ended at iteration:', i, 'of', max_iter
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close(12)

100 format(A," ",F13.10)
end program golden_ratio
```





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program golden_ratio
! experiments with the golden ratio iterative relation
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namelist /golden_inputs/ phi_start, tol, max_iter

open(11,FILE='golden.in',STATUS='old')
read(11,golden_inputs)
close(11)

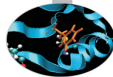
phi_old = phi_start
do i=1,max_iter
  phi = 1.0d0/phi_old + 1.0d0
  if (abs(phi - phi_old) < tol) exit
  phi_old = phi
end do

open(12,FILE='golden.out')
write(12,100) 'Start value:',phi_start
write(12,100) 'Tolerance:',tol
write(12,'(2(A," ",I11," "))') 'Ended at iteration:', i, 'of', max_iter
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end program golden_ratio
```



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- ▶ File content is structured, self-describing, order independent, comments are allowed:

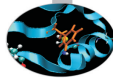
```
&golden_inputs  
tol=1.e-4      ! tolerance  
phi_start=5.0 ! 0th iteration  
max_iter=1000000 /
```



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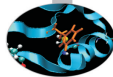
- ▶ Items missing in the input will retain previous value



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- ▶ Items missing in the input will retain previous value
- ▶ Items can be added to a namelist in different statements, but code like this easily misleads readers (and you read your own codes, don't you?)
- ▶ Use them to make input robust, in output mostly good for debugging



Array Syntax

Input/Output

Formatted I/O

File I/O

Namelist

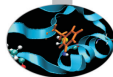
**Internal Files**

Unformatted I/O

Robust I/O



## Iterative search for the Golden Ratio



```
program golden_ratio
! experiments with the golden ratio iterative relation
implicit none
integer, parameter :: rk = kind(1.0d0)
real(rk) :: phi, phi_old
real(rk) :: phi_start, tol
integer :: i, max_iter, test_no
character(15) :: outfilename

namelist /golden_inputs/ phi_start, tol, max_iter, test_no

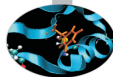
test_no = 1
open(11,FILE='golden.in',STATUS='old')
read(11,golden_inputs)
close(11)

phi_old = phi_start
do i=1,max_iter
  phi = 1.0d0/phi_old + 1.0d0
  if (abs(phi - phi_old) < tol) exit
  phi_old = phi
end do

write(outfilename,'("golden",I5.5,".out")') test_no
open(12,FILE=outfilename)
write(12,100) 'Start value:',phi_start
write(12,100) 'Tolerance:',tol
write(12,'(2(A," ",I11," "))') 'Ended at iteration:', i, 'of', max_iter
write(12,100) 'Final value:',phi
close(12)

100 format(A," ",F13.10)
end program golden_ratio
```

## Iterative search for the Golden Ratio



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program golden_ratio
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integer, parameter :: rk = kind(1.0d0)
real(rk) :: phi, phi_old
real(rk) :: phi_start, tol
integer :: i, max_iter, test_no
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namelist /golden_inputs/ phi_start, tol, max_iter, test_no

test_no = 1
open(11,FILE='golden.in',STATUS='old')
read(11,golden_inputs)
close(11)

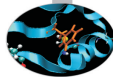
phi_old = phi_start
do i=1,max_iter
  phi = 1.0d0/phi_old + 1.0d0
  if (abs(phi - phi_old) < tol) exit
  phi_old = phi
end do

write(outfilename,'("golden",I5.5,".out")') test_no
open(12,FILE=outfilename)
write(12,100) 'Start value:',phi_start
write(12,100) 'Tolerance:',tol
write(12,'(2(A," ",I11," "))') 'Ended at iteration:', i, 'of', max_iter
write(12,100) 'Final value:',phi
close(12)

100 format(A," ",F13.10)
end program golden_ratio
```



- ▶ **character** variables of default kind can be specified in place of units in **read** and **write** statements



- ▶ **character** variables of default kind can be specified in place of units in **read** and **write** statements
- ▶ Writing to internal files is good to:
  - ▶ dynamically build file names according to a pattern (like number of iterations)
  - ▶ dynamically assemble complex I/O formats, depending on actual data
  - ▶ prepare complex labels for plot data formats
  - ▶ build commands to be sent to hardware devices
  - ▶ ...



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- ▶ Writing to internal files is good to:
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  - ▶ dynamically assemble complex I/O formats, depending on actual data
  - ▶ prepare complex labels for plot data formats
  - ▶ build commands to be sent to hardware devices
  - ▶ ...
- ▶ Reading from internal files can be useful to read complex inputs
  - ▶ You have a textual input file sporting different formats
  - ▶ And the right format depends on actual data in the file
  - ▶ Just read each line in a **character** variable, suitably sized
  - ▶ Pick the suitable format
  - ▶ And use it to read from the variable itself



- ▶ Play with `goldenfile.f90`, `goldenfn1.f90`, and `goldenio.f90`:
  - ▶ writing input files
  - ▶ writing good and bad data in input files
  - ▶ giving input files wrong file names



Array Syntax

Input/Output

Formatted I/O

File I/O

Namelist

Internal Files

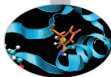
**Unformatted I/O**

Robust I/O



- ▶ Formatted I/O is good, but:
  - ▶ internal data format is much more compact
  - ▶ and roundoff may happen, making recovery of original values impossible
  - ▶ and conversion takes time





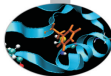
```
program iterative_inversion
! experiments with matrix iterative inversion
implicit none
real, dimension(4,4) :: a, x, x_old, x_start
real :: tol, err
integer :: i, max_iter

open(21,FILE='input.dat',FORM='unformatted',STATUS='old')
read(21) a
read(21) x_start
read(21) tol,max_iter
close(21)

x_old = x_start
do i=1,max_iter
  x = 2.0*x_old - matmul(x_old,matmul(a,x_old))
  err = maxval(abs(x - x_old))
  if (err < tol) exit
  x_old = x
end do

open(22,FILE='itinv.dat',FORM='unformatted')
write(22) a
write(22) x_start
write(22) tol,max_iter
write(22) i
write(22) x
close(22)

end program iterative_inversion
```



```
program iterative_inversion
! experiments with matrix iterative inversion
implicit none
real, dimension(4,4) :: a, x, x_old, x_start
real :: tol, err
integer :: i, max_iter

open(21,FILE='input.dat',FORM='unformatted',STATUS='old')
read(21) a
read(21) x_start
read(21) tol,max_iter
close(21)

x_old = x_start
do i=1,max_iter
  x = 2.0*x_old - matmul(x_old,matmul(a,x_old))
  err = maxval(abs(x - x_old))
  if (err < tol) exit
  x_old = x
end do

open(22,FILE='itinv.dat',FORM='unformatted')
write(22) a
write(22) x_start
write(22) tol,max_iter
write(22) i
write(22) x
close(22)

end program iterative_inversion
```



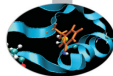
- ▶ Formatted I/O is good, but:
  - ▶ internal data format is much more compact
  - ▶ and roundoff may happen, making recovery of original values impossible
  - ▶ and conversion takes time
- ▶ Unformatted I/O is used to store and recover data in internal representation
  - ▶ Just give **FORM='unformatted'** option when opening the file
  - ▶ And omit format in **read** and **write** statements



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  - ▶ internal data format is much more compact
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  - ▶ and conversion takes time
- ▶ Unformatted I/O is used to store and recover data in internal representation
  - ▶ Just give **FORM='unformatted'** option when opening the file
  - ▶ And omit format in **read** and **write** statements
- ▶ Unformatted I/O is performed on a *record* basis
  - ▶ In unformatted mode, each **write** writes a *record*
  - ▶ As we'll see, this allows walking your files backward and forward
  - ▶ But has interesting consequences, as more than your data is written to your file...



- ▶ Modify `itinv.f90` to perform unformatted I/O
- ▶ To test it, you'll need an additional program:
  - ▶ taking text input from keyboard or initializing all needed data
  - ▶ to write a good unformatted input file for the new version of `itinv.f90`



▶ Try different ways to output the results:

▶ element-wise

```
do j=1,n
  do i=1,n
    write(79) a(i,j)
  end do
end do
```

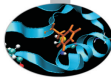
▶ column-wise, using an implied do-loop:

```
do j=1,n
  write(79) (a(i,j), i=1,n)    ! a(:,j) will also do
end do
```

▶ with two implied do-loops:

```
write(79) ((a(i,j), i=1,n), j=1,n)
```

▶ Can you spot the difference?



▶ Try different ways to output the results:

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```
do j=1,n
  do i=1,n
    write(79) a(i,j)
  end do
end do
```

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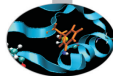
▶ Can you spot the difference?

▶ Not a big issue for  $4 \times 4$  matrices, but think of a  $256 \times 256 \times 1024$  grid!



- ▶ **read** always advance to next record, even if you read only part of the record (or possibly nothing)
- ▶ **backspace (u)** moves position for subsequent I/Os to the record preceding the current one
- ▶ **rewind (u)** moves position for subsequent I/Os to file beginning
- ▶ To allow positioning back and forth, a four bytes record marker is added in 32 bit mode (eight bytes in 64 bit mode) before and after each record
- ▶ Best practice: write data in whole blocks





- ▶ Record markers added in unformatted I/O make exchanging data with other programs (notably C ones) troublesome
- ▶ `open (unit, ..., ACCESS='stream', ...)` is a new method to access external files
- ▶ No record markers are written before or after a `write`
  - ▶ Thus, advancing or backspacing over records is not possible
  - ▶ But required position may be specified by:  

```
write(unit,POS=position) x  
read(unit,POS=position) y
```
- ▶ Best practice: if you are really serious about data exchanges, across different programs and systems, use libraries like HDF5, VTK, CGNS



Array Syntax

Input/Output

Formatted I/O

File I/O

Namelist

Internal Files

Unformatted I/O

**Robust I/O**



- ▶ You may happen to:
  - ▶ Try to open a new file, when one with same name already exists
  - ▶ Look for an existing file, which is missing
  - ▶ Encounter an unexpected end of record in a **read**
  - ▶ Encounter an unexpected end of file while reading
  - ▶ Run out of disk space while writing
  - ▶ Try writing to a read-only file
  - ▶ ...



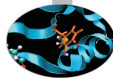
- ▶ You may happen to:
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- ▶ And get an unfriendly runtime error



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  - ▶ Encounter an unexpected end of file while reading
  - ▶ Run out of disk space while writing
  - ▶ Try writing to a read-only file
  - ▶ ...
- ▶ And get an unfriendly runtime error
- ▶ Or you may need to open a file in a library you are writing
  - ▶ And use a unit already opened in a calling program
  - ▶ The previously opened unit is automatically closed
  - ▶ With surprising consequences on program behavior

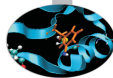


- ▶ All I/O statements accept an **IOSTAT=*ios*** option
  - ▶ ***ios*** must be an integer variable of default kind
  - ▶ Set to zero on success
  - ▶ Set to negative values on end of file or record (in Fortran 2003, **`iostat_end`** and **`iostat_eor`** respectively, from **`iso_fortran_env`** module)
  - ▶ Set to positive values on error
  - ▶ Execution will not stop
- ▶ Use it to identify the issue, and recover or fail gracefully



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  - ▶ **ios** must be an integer variable of default kind
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  - ▶ Set to positive values on error
  - ▶ Execution will not stop
- ▶ Use it to identify the issue, and recover or fail gracefully
- ▶ All I/O statements accept an **ERR=err-label** option
  - ▶ **err-label** is a statement label in the same program unit
  - ▶ Flow control jumps to **err-label** in case of error
- ▶ Use it to centralize error management and recovery
- ▶ Together with **iostat**, of course

In Doubt? *inquire!*



- ▶ Let's assume **ans** is a logical variable, **k** is an integer variable, and **s** is a character variable of suitable length



## In Doubt? *inquire!*



- ▶ Let's assume **ans** is a logical variable, **k** is an integer variable, and **s** is a character variable of suitable length
- ▶ **inquire (FILE='input.dat', EXIST=ans)** will set **ans** to **.true.** if file **input.dat** exists

## In Doubt? *inquire!*



- ▶ Let's assume **ans** is a logical variable, **k** is an integer variable, and **s** is a character variable of suitable length
- ▶ `inquire(FILE='input.dat',EXIST=ans)` will set **ans** to `.true.` if file `input.dat` exists
- ▶ `inquire(FILE='input.dat',OPENED=ans)` will set **ans** to `.true.` if file `input.dat` is already opened

## In Doubt? *inquire!*

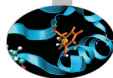


- ▶ Let's assume **ans** is a logical variable, **k** is an integer variable, and **s** is a character variable of suitable length
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- ▶ **inquire(FILE='input.dat',OPENED=ans)** will set **ans** to **.true.** if file **input.dat** is already opened
- ▶ **inquire(15,OPENED=ans)** will set **ans** to **.true.** if a file is already opened on unit **15**



- ▶ Let's assume **ans** is a logical variable, **k** is an integer variable, and **s** is a character variable of suitable length
- ▶ **inquire(FILE='input.dat',EXIST=ans)** will set **ans** to **.true.** if file **input.dat** exists
- ▶ **inquire(FILE='input.dat',OPENED=ans)** will set **ans** to **.true.** if file **input.dat** is already opened
- ▶ **inquire(15,OPENED=ans)** will set **ans** to **.true.** if a file is already opened on unit 15
- ▶ **inquire(FILE='input.dat',NUMBER=k)** will set **k** to **-1** if file **input.dat** is not opened, to connected unit otherwise

# More Doubts? *inquire* More!



- ▶ `inquire(15, FORM=s)` will set `s` to 'FORMATTED' or 'UNFORMATTED' if unit 15 is connected for formatted or unformatted I/O respectively, to 'UNDEFINED' otherwise

## More Doubts? *inquire* More!



- ▶ **`inquire(15, FORM=s)`** will set **`s`** to '**`FORMATTED`**' or '**`UNFORMATTED`**' if unit 15 is connected for formatted or unformatted I/O respectively, to '**`UNDEFINED`**' otherwise
- ▶ **`inquire(15, ACTION=s)`** will set **`s`** to '**`READ`**' or '**`WRITE`**' or '**`READWRITE`**', depending on what actions are allowed on unit 15, to '**`UNDEFINED`**' if unconnected

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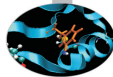
## More Doubts? *inquire* More!



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- ▶ And many more variations, look to manuals
- ▶ Of course, **IOSTAT** and **ERR** can be useful on ***inquire*** too



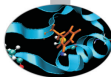
- ▶ Write a program that:
  - ▶ reads an 'arbitrarily' long column of real numbers from an ASCII file
  - ▶ prints maximum, minimum, average of the numbers
  - ▶ and prints the  $\lfloor n/2 \rfloor$ -th row where  $n$  is the length of the column



## Part IV

# Derived Types and Memory Management

Derived types, operators overloading, parametric types and inheritance. Memory management, dynamic allocation and memory heap. Pointers. C and Fortran interoperability.



Extending the Language  
Derived Types  
Operators Overloading  
Parameterized Types  
Extending Types, and Objects

Managing Memory

Conclusions



## Extending the Language

### Derived Types

Operators Overloading

Parameterized Types

Extending Types, and Objects

## Managing Memory

## Conclusions

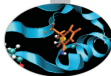


- ▶ Fortran allows programmers to add new types, built as assemblies of existing ones

```
type position
  real :: x, y, z
end type position
```

```
type velocity
  real :: x, y, z
end type velocity
```

- ▶ Components in different derived types may have the same name (not a surprise!)
- ▶ **type(position) :: r** declares a variable of type **position**
- ▶ Components of a derived type can be accessed like this:  
**r%y = 0.0**



- ▶ Derived types are not second class citizens
- ▶ Thus derived types (also termed *structures*) can be assembled from other derived types too

```

type particle
  type(position) :: r
  type(velocity) :: v
  real :: mass
end type particle
  
```

```

type atom
  type(position) :: r
  type(velocity) :: v
  real :: mass ! In atomic units
  integer :: an ! Atomic number
end type atom
  
```

- ▶ `type(particle) :: p` declares a variable of type `particle`
- ▶ Components of a component of a variable can be accessed like this: `p%v%z = 0.0`



```
type(atom) :: h1, h2, he

h1%r = position(0.0, 0.0, 0.0)
h1%v = velocity(1.0, -1.0, 0.0)
h1%mass = 1.00794
h1%an = 1 ! Assigns atomic number

h2 = h1 ! Intrinsic assignment

he = atom(position(1.0, 0.0, -1.0), h2%v, 4.002602, 2)
```

- ▶ Derived type name can be used to construct values of the type
- ▶ Unsurprisingly, `velocity()` is termed a *constructor*
- ▶ Values passed as argument to constructors must be ordered as in type definition
- ▶ Assignment is intrinsically available





- ▶ Derived types boil down (possibly recursively) to collections of intrinsic types
- ▶ And behavior is coherent with I/O of complex values and arrays
- ▶ All single intrinsic type (sub)components will be processed in sequence
- ▶ If you want control of the conversion:
  - ▶ a proper edit descriptor must be provided for each component
  - ▶ in same order as components are declared in type declaration
- ▶ Fortran 2003 introduces the **D<sub>T</sub>** edit descriptor to give users total control



## Extending the Language

Derived Types

**Operators Overloading**

Parameterized Types

Extending Types, and Objects

## Managing Memory

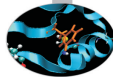
## Conclusions



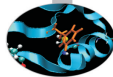
- ▶ Binary operator + can be used to add:



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  - ▶ a pair of integer values



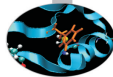
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- ▶ Binary operator  $+$  can be used to add:
  - ▶ a pair of integer values
  - ▶ a pair of real values
  - ▶ a pair of complex values



- ▶ Binary operator  $+$  can be used to add:
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  - ▶ two integer values of different kinds



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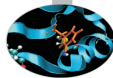




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  - ▶ two real values of different kinds
  - ▶ two complex values of different kinds
  - ▶ an integer and a real value



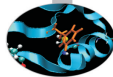
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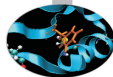


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- ▶ It's like the meaning of  $+$  is 'overloaded'
  - ▶ Different machine code is generated depending on operand types
- ▶ And ditto for  $-$ ,  $*$ ,  $/$ ,  $>$ ,  $>=$ , ...

# Bringing Abstractions Further



Wouldn't it be nice to have arithmetic operators work on structures?

```
interface operator(-)
  function subvel(p1, p2)
    type(velocity), intent(in) :: p1, p2
    type(velocity) :: subvel
  end function
end interface operator(-)
```

```
interface operator(-)
  function chsvel(p)
    type(velocity), intent(in) :: p
    type(velocity) :: chsvel
  end function
end interface operator(-)
```

```
function subvel(p1, p2)
  implicit none
  type(velocity), intent(in) :: p1, p2
  type(velocity) :: subvel
  subvel%x = p1%x-p2%x; subvel%y = p1%y-p2%y; subvel%z = p1%z-p2%z
end function subvel
```

```
function chsvel(p)
  implicit none
  type(velocity), intent(in) :: p
  type(velocity) :: chsvel
  chsvel%x = -p%x; chsvel%y = -p%y; chsvel%z = -p%z
end function chsvel
```





- ▶ We are fitting an infinite space into a finite box with periodic boundary conditions
- ▶ Wouldn't it be nice to define our operators with custom functionality?

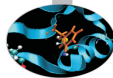
```
interface operator(+)  
  function addpos(p1, p2)  
    type(position), intent(in) :: p1, p2  
    type(position) :: addpos  
  end function  
end interface operator(+)  
  
function addpos(p1, p2) ! Adds positions with periodic boundary conditions  
  implicit none  
  type(position), intent(in) :: p1, p2  
  type(position) :: addpos  
  real,parameter :: boxwidth = 128.0  
  
  addpos%x = modulo(p1%x+p2%x, boxwidth)  
  addpos%y = modulo(p1%y+p2%y, boxwidth)  
  addpos%z = modulo(p1%z+p2%z, boxwidth)  
end function addpos
```



- ▶ **interface operator(*op-name*)** lets you overload *op-name* with a generic procedure
  - ▶ Arguments must be **intent (in)** and can be either one or two
  - ▶ *op-name* may be an intrinsic operator, or a **.new\_name.**



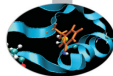
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  - ▶ highest for new unary operators
  - ▶ lowest for new binary operators



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- ▶ Now velocities may be added as intrinsic arithmetic types
- ▶ And defining subtraction is an easy job
- ▶ Positions may be added as usual intrinsic variables and boundary conditions are automatically imposed



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- ▶ Now velocities may be added as intrinsic arithmetic types
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- ▶ Positions may be added as usual intrinsic variables and boundary conditions are automatically imposed
- ▶ Time for a module



```

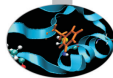
module periodic_box
  implicit none
  real, private, parameter :: boxwidth = 128.0
  private addpos, addvel, chsvel, subvel, subpos

  type position
    real :: x, y, z
  end type position

  type velocity
    real :: x, y, z
  end type velocity

  interface operator(+)
    module procedure addpos
    module procedure addvel
  end interface operator(+)
! ...
contains
  function addpos(p1, p2) ! Adds positions with periodic boundary conditions on x
    type(position), intent(in) :: p1, p2
    type(position) :: addpos
    addpos%x = modulo(p1%x+p2%x, boxwidth)
    addpos%y = modulo(p1%y+p2%y, boxwidth)
    addpos%z = modulo(p1%z+p2%z, boxwidth)
  end function addpos

  function addvel
    !...
  end function addvel
! ...
end module periodic_box
  
```



- ▶ Again, modules are the best way of grouping related stuff
- ▶ Again, with modules and module procedures we don't need to write interface blocks
- ▶ Modules let us hide implementation details
- ▶ Best practice: put structure definitions and related functions and operators in modules
  - ▶ Anyway, they will be used together
  - ▶ When dealing with nested types with many related functions, a hierarchy of modules would probably help
  - ▶ Because, of course, you can **use** modules in a module



- ▶ Write a module that defines:
  - ▶ A new type *vector* made up of three real components
  - ▶ Operator `.cross.` for cross product
  - ▶ Operator `+` to sum two *vectors*
- ▶ Write a program to test your module

```

program test_class_vector
  use class_vector

  implicit none

  type(vector) :: v, w, z

  v=vector(1.d0,0.d0,0.d0)
  w=vector(0.d0,1.d0,0.d0)
  z=vector(0.d0,0.d0,1.d0)

  write(*,*) v+w.cross.z

end program test_class_vector
  
```

- ▶ Definition of cross product:

$$a \times b = (a_2 b_3 - a_3 b_2) \hat{i} + (a_3 b_1 - a_1 b_3) \hat{j} + (a_1 b_2 - a_2 b_1) \hat{k}$$

- ▶ Then extend operators to have them work with array of vectors: it's elementary!





## Extending the Language

Derived Types

Operators Overloading

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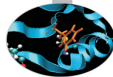
- ▶ What if we wanted different **kinds of points**?

- ▶ This is a possibility:

```
type point
  real( selected_real_kind(5) ) :: x, y, z
end type point
```

```
type widepoint
  real( selected_real_kind(12) ) :: x, y, z
end type widepoint
```

- ▶ But not very elegant, nor easy to manage



- ▶ In Fortran 2003, types may have kind type parameters:

```
type point(point_kind)
  integer, kind :: point_kind = kind(0.0)
  real(point_kind) :: x, y, z
end type point
```

```
type(point(point_kind=kind(0.0))) :: apoint
type(point) :: anotherpoint
type(point(selected_real_kind(12))) :: awiderpoint
```

- ▶ **kind** states that this type parameter behaves as a kind
- ▶ And it works as **kind** does for intrinsic types



- ▶ Structures may have array components

```

type segments(point_kind)
  integer, kind :: point_kind = kind(0.0)
  type(point(point_kind)), dimension(100) :: start_point
  type(point(point_kind)), dimension(100) :: end_point
end type segments
  
```

- ▶ Our `segments` type looks a bit rigid, doesn't it?

- ▶ Derived type parameters come to rescue:

```

type segments(point_kind, n)
  integer, kind :: point_kind = kind(0.0)
  integer, len :: n
  type(point(point_kind)), dimension(n) :: start_point
  type(point(point_kind)), dimension(n) :: end_point
end type segments
  
```

```

type(segments(n=100)) :: ahundredsegments
type(segments(n=1000)) :: athousandsegments
  
```



## Extending the Language

Derived Types

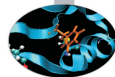
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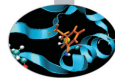
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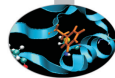
Conclusions



- ▶ So, we are able to define new types, and specialized procedures and operators to use them

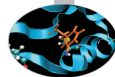


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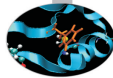


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- ▶ But **point**, **position**, and **velocity** have the same components
  - ▶ And that's always true, whatever the space dimensions
  - ▶ But they are conceptually (and dimensionally!) different things

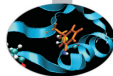




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  - ▶ And a **ion** would simply add a **charge** component



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- ▶ And **particle**, and **atom** share identical components
  - ▶ And a **ion** would simply add a **charge** component
- ▶ Wouldn't it be nice to 'inherit' from one type to another?
  - ▶ Yeah, and easier to manage, too!
  - ▶ And this is what CS priests call *Object-Oriented* programming, and is so trendy!



```
type point
  real :: x, y, z
end type point

type, extends(point) :: position
end type position

type, extends(point) :: velocity
end type velocity

type particle
  type(position) :: r
  type(velocity) :: v
  real :: mass
end type particle

type, extends(particle) :: atom
  integer :: an ! atomic number
end type atom

type, extends(atom) :: ion
  integer :: charge ! in units of elementary charge
end type ion
```

- ▶ **extends** means that the new type has the same components, and possibly more
- ▶ Now we still have to write procedures and operators, don't we?

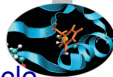


- ▶ A type extension includes an implicit component with the same name and type as its parent type
  - ▶ this can come in handy when the programmer wants to operate on components specific to a parent type

```
type(ion) :: p                ! declare p as a ion object

p%mass                       ! access mass component for p
p%atom%mass                  ! another way
p%atom%particle%mass        ! ...
```

- ▶ We often say the child and parent types have a “is a” relationship
  - ▶ an atom “is” a particle
  - ▶ but a particle is not an atom because the atomic component may be found in atom but not in particle



- ▶ Consider the case you have to evolve the position of a particle according to a given velocity field
  - ▶ atoms or ions may behave in the (nearly) same way wrt this evolution
  - ▶ and you do not want to write two (nearly) identical procedures for the two types
  
- ▶ Polymorphic procedures are the right way
  - ▶ i.e., procedures which can take one or more polymorphic variables as arguments
  - ▶ “polymorphic variable” = variable whose data type is dynamic at runtime
  - ▶ the **class** keyword allows F2003 programmers to create polymorphic variables
  - ▶ use it for dummy arguments (the simplest usage, not the only one)



```

subroutine setMass(p, m)
class(particle) :: p
real :: m
p%mass = m
end subroutine setMass
  
```

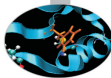
- ▶ The **p** dummy argument is polymorphic, based on the usage of **class(particle)**
- ▶ The subroutine can operate on objects that satisfy the "is a" particle relationship
  - ▶ **setMass** can be called passing a particle, atom, ion, or any future type extension of particle

```

type(particle) :: pa      ! declare an instance of particle
type(atom)      :: at     ! declare an instance of atom
type(ion)       :: io     ! declare an instance of ion
  
```

```

call setMass(pa, mm)      ! set the mass for a particle
call setMass(at, mm)     ! set the mass for an atom
call setMass(io, mm)     ! set the mass for a ion
  
```



- ▶ By default, only those components found in the declared type of an object are accessible



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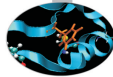




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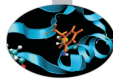


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  - ▶ and optional arguments come in handy



```

subroutine initialize(p, mm, rr, vv, aan, ccharge)
class(particle) :: p
real :: mm
type(position) :: rr
type(velocity) :: vv
integer, optional :: aan, ccharge
p%mass = mm
p%r = rr
p%v = vv
select type (p)
type is (particle)
  ! no further initialization required
class is (atom)
  ! atom or ion specific initializations
  if (present(aan)) then
    p%an = aan
  else
    p%an = 1
  endif
class is (ion)
  if (present(ccharge)) then
    p%charge = ccharge
  else
    p%charge = 0
  endif
class default
  ! give error for unexpected/unsupported type
  stop 'initialize: unexpected type for p object!'
end select
end subroutine initialize
  
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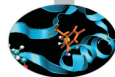


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- ▶ There are two styles of type checks that we can perform
  - ▶ **type is**: satisfied if the dynamic type of the object is the same as the type specified in parentheses
  - ▶ **class is**: satisfied if the dynamic type of the object is the same or an extension of the specified type in parentheses

# Selecting type



```
subroutine initialize(p, mm, rr, vv, aan, ccharge)
class(particle) :: p
real :: mm
type(position) :: rr
type(velocity) :: vv
integer, optional :: aan, ccharge
p%mass = mm
p%r = rr
p%v = vv
select type (p)
type is (particle)
    ! no further initialization required
class is (atom)
    ! atom or ion specific initializations
    if (present(aan)) then
        p%an = aan
    else
        p%an = 1
    endif
class is (ion)
    if (present(ccharge)) then
        p%charge = ccharge
    else
        p%charge = 0
    endif
class default
    ! give error for unexpected/unsupported type
    stop 'initialize: unexpected type for p object!'
end select
end subroutine initialize
```



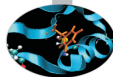
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  - ▶ **class is**: satisfied if the dynamic type of the object is the same or an extension of the specified type in parentheses
- ▶ Best practice: add a **class default** branch and print error when p is not an extension of **particle** type
  - ▶ an empty **type is (particle)** branch may be required to avoid getting error when p is only a particle

# Selecting type



```

subroutine initialize(p, mm, rr, vv, aan, ccharge)
class(particle) :: p
real :: mm
type(position) :: rr
type(velocity) :: vv
integer, optional :: aan, ccharge
p%mass = mm
p%r = rr
p%v = vv
select type (p)
type is (particle)
    ! no further initialization required
class is (atom)
    ! atom or ion specific initializations
    if (present(aan)) then
        p%an = aan
    else
        p%an = 1
    endif
class is (ion)
    if (present(ccharge)) then
        p%charge = ccharge
    else
        p%charge = 0
    endif
class default
    ! give error for unexpected/unsupported type
    stop 'initialize: unexpected type for p object!'
end select
end subroutine initialize

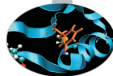
```



- ▶ Objects in Fortran 2003
  - ▶ A Fortran 90/95 module can be viewed as an object because it can encapsulate both data and procedures
  - ▶ But, derived types in F2003 are considered objects because they now can encapsulate data as well as procedures
  - ▶ Modules and types work together...
- ▶ Procedures encapsulated in a derived type are called type-bound procedures (“methods” in OO jargon)

```
type particle
  type(position) :: r
  type(velocity) :: v
  real :: mass
  contains
    procedure :: initialize => initialize_particle
end type particle
```

- ▶ **initialize\_particle** is the name of the underlying procedure to be implemented
- ▶ Explicit interface is required: wrap in a module!



► Employing modules and types to design objects

```

module particle_mod
  type particle
    type(position) :: r
    type(velocity) :: v
    real :: mass
  contains
    procedure :: initialize => initialize_particle
  end type particle
  type, extends(particle) :: atom
  ...
end type atom
  type, extends(atom) :: ion
  ...
end type ion
  contains
  ! insert the implementation or at least the interface of initialize_particle
  subroutine initialize_particle(p, mm, rr, vv, aan, ccharge)
  class(particle) :: p
  ...
  end subroutine initialize_particle
end module particle_mod
  
```



- ▶ `initialize` is the name to be used to invoke the type bound procedure

```

use particle_mod
type(particle) :: p           ! declare an instance of particle
call p%initialize(mas, pos, vel) ! initialize particle
  
```

- ▶ What about the first dummy argument of `initialize`?
  - ▶ it is known as the *passed-object* dummy argument
  - ▶ must be declared `class` and of the same type as the derived type that defined the type-bound procedure
  - ▶ by default, it is the first dummy argument in the type-bound procedure: it receives the object that invoked the type-bound procedure

- ▶ It is possible to pass another argument in place of the first one

```
procedure, pass(p) :: initialize
```

- ▶ ...or to avoid passing it at all

```
procedure, nopass :: initialize
```

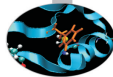


- ▶ A child type inherits or reuses components from their parent or ancestor types: this applies to both data and procedures

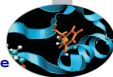
```

type(particle) :: pa           ! declare an instance of particle
type(atom)     :: at           ! declare an instance of atom
type(ion)      :: io           ! declare an instance of ion
call pa%initialize(mas, pos, vel) ! initialize a particle
call at%initialize(mas, pos, vel, anu) ! initialize an atom
call io%initialize(mas, pos, vel, anu, cha) ! initialize a ion
  
```

- ▶ **initialize** behaves accordingly to the passed arguments, i.e. using **optional** and **select type** features
- ▶ Sometimes, another approach may be more appropriate: overriding!



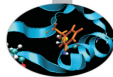
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module particle_mod
type particle
  type(position) :: r
  type(velocity) :: v
  real :: mass
  contains
    procedure :: initialize => initialize_particle
end type particle
type, extends(particle) :: atom
  ...
  contains
    procedure :: initialize => initialize_atom
end type atom
type, extends(atom) :: ion
  ...
end type ion
contains
! insert the implementation or at least the interface of initialize
subroutine initialize_particle(p, mm, rr, vv, aan, cch)
class(particle) :: p
  ...
end subroutine initialize_particle
subroutine initialize_atom(p, mm, rr, vv, aan, cch)
class(atom) :: p
  ...
end subroutine initialize_atom
end module particle_mod
```



```
type(particle) :: pa           ! declare an instance of particle
type(atom)     :: at           ! declare an instance of atom
type(ion)      :: io           ! declare an instance of ion
call pa%initialize(mas, pos, vel) ! calls initialize_particle
call at%initialize(mas, pos, vel, anu) ! calls initialize_atom
call io%initialize(mas, pos, vel, anu, cha) ! calls initialize_atom
```

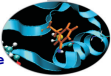
- ▶ Beware: an overriding type-bound procedure must have exactly the same interface as the overridden procedure except for the passed-object dummy argument which must be **class(new-type)**
  - ▶ optional arguments may hide useless arguments





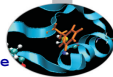
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type particle
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  ...
end type ion
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! insert the implementation or at least the interface of initialize
subroutine initialize_particle(p, mm, rr, vv, aan, cch)
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end subroutine initialize_particle
subroutine initialize_atom(p, mm, rr, vv, aan, cch)
class(atom) :: p
  ...
end subroutine initialize_atom
end module particle_mod
```

# Override with care



```
type(particle) :: pa           ! declare an instance of particle
type(atom)      :: at         ! declare an instance of atom
type(ion)       :: io         ! declare an instance of ion
call pa%initialize(mas, pos, vel) ! calls initialize_particle
call at%initialize(mas, pos, vel, anu) ! calls initialize_atom
call io%initialize(mas, pos, vel, anu, cha) ! calls initialize_atom
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- ▶ Beware: an overriding type-bound procedure must have exactly the same interface as the overridden procedure except for the passed-object dummy argument which must be **class (new-type)**
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- ▶ Of course, it is still possible to explicitly invoke the version defined by a parent type instead of the overridden one



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type(particle) :: pa           ! declare an instance of particle
type(atom)      :: at         ! declare an instance of atom
type(ion)       :: io         ! declare an instance of ion
call pa%initialize(mas, pos, vel) ! calls initialize_particle
call at%initialize(mas, pos, vel, anu) ! calls initialize_atom
call io%initialize(mas, pos, vel, anu, cha) ! calls initialize_atom
```

- ▶ Beware: an overriding type-bound procedure must have exactly the same interface as the overridden procedure except for the passed-object dummy argument which must be **class (new-type)**
  - ▶ optional arguments may hide useless arguments
- ▶ Of course, it is still possible to explicitly invoke the version defined by a parent type instead of the overridden one
- ▶ And it is possible to prevent any type extensions from overriding a particular type-bound procedure

```
procedure, non_overridable :: initialize
```



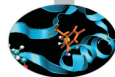
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- ▶ Hiding data:
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  - ▶ changes to the data structure will not affect codes using our class provided that we don't change interfaces
- ▶ Hiding procedures: e.g., prevent users from calling low-level procedures



- ▶ Fortran 2003 adds “private” and “public” keywords for derived types
  - ▶ beware of the placement of the keywords, in modules and/or in types: confused?

```

module particle_mod
private ! hide the implementation of type-bound procedures
public :: average_position_particle ! allow access to particle averaging position
type, public :: particle
  private ! hide the data underlying details
    type(position) :: r
    type(velocity) :: v
    real :: mass
  contains
    private ! hide the type bound procedures by default
      procedure :: check_init => check_init_particle ! private type-bound procedure
      procedure, public :: initialize => initialize_particle ! allow access to TBP
end type particle
contains
! implementation of type-bound procedures
subroutine initialize_particle(p, mm, rr, vv, aan, cch)
...

subroutine check_init_particle(p)
...

subroutine average_position_particle(p1,p2)
class(particle) :: p1, p2
...
end subroutine average_position_particle
end module particle_mod
    
```



- ▶ Data Polymorphism:
  - ▶ as how polymorphic dummy arguments form the basis to procedure polymorphism...
  - ▶ ...polymorphic non-dummy variables form the basis to data polymorphism





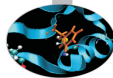
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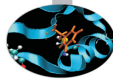
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- ▶ Abstract types and deferred bindings
- ▶ Finalization



## Extending the Language

### Managing Memory

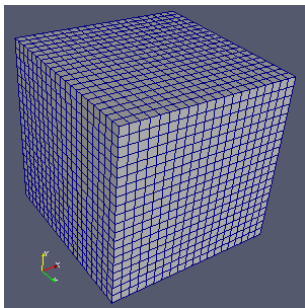
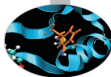
- Dynamic Memory Allocation

- Fortran Pointers

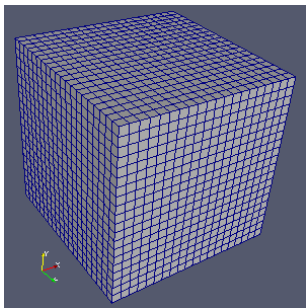
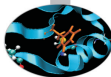
- Bridging the Gap with C

- Sketchy Ideas on Data Structures

## Conclusions



- ▶ Let's imagine we have to solve a PDE



- ▶ Let's imagine we have to solve a PDE
- ▶ On a dense, Cartesian, uniform grid
  - ▶ Mesh axes are parallel to coordinate ones
  - ▶ Steps along each direction have the same size
  - ▶ And we have some discretization schemes in time and space to solve for variables at each point





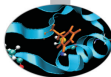
```
integer, parameter :: NX = 200
integer, parameter :: NY = 450
integer, parameter :: NZ = 320

integer, parameter :: rk = selected_real_kind(12)

real(rk) :: deltax ! Grid steps
real(rk) :: deltay
real(rk) :: deltaz

real(rk) :: u(NX,NY,NZ)
real(rk) :: v(NX,NY,NZ)
real(rk) :: w(NX,NY,NZ)
real(rk) :: p(NX,NY,NZ)
```

- We could write something like that in a module, and use it everywhere



```
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- ▶ We could write something like that in a module, and use it everywhere
- ▶ But it has annoying consequences
  - ▶ Recompile each time grid resolution changes
  - ▶ A slow process, for big programs
  - ▶ And error prone, as we may forget about



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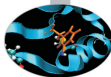
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- ▶ We could write something like that in a module, and use it everywhere
- ▶ But it has annoying consequences
  - ▶ Recompile each time grid resolution changes
  - ▶ A slow process, for big programs
  - ▶ And error prone, as we may forget about
- ▶ Couldn't we size data structures according to user input?

# A Recurrent Issue: SoA or AoS



```
► type flow
   real(rk) :: u(NX,NY,NZ)
   real(rk) :: v(NX,NY,NZ)
   real(rk) :: w(NX,NY,NZ)
   real(rk) :: p(NX,NY,NZ)
end type
```

```
type(flow) :: f
```

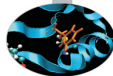
Or

```
type flow
   real(rk) :: u,v,w,p
end type
```

```
type(flow) :: f(NX,NY,NZ)
```

Which one is best?

# A Recurrent Issue: SoA or AoS



```
► type flow
  real(rk) :: u(NX,NY,NZ)
  real(rk) :: v(NX,NY,NZ)
  real(rk) :: w(NX,NY,NZ)
  real(rk) :: p(NX,NY,NZ)
end type
```

```
type(flow) :: f
```

Or

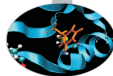
```
type flow
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Which one is best?

► Both have merits

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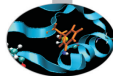
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type flow
  real(rk) :: u,v,w,p
end type
```

```
type(flow) :: f(NX,NY,NZ)
```

Which one is best?

- Both have merits
- The choice strongly depends on the computer architecture

# A Recurrent Issue: SoA or AoS



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► type flow
   real(rk) :: u(NX,NY,NZ)
   real(rk) :: v(NX,NY,NZ)
   real(rk) :: w(NX,NY,NZ)
   real(rk) :: p(NX,NY,NZ)
end type
```

```
type(flow) :: f
```

Or

```
type flow
   real(rk) :: u,v,w,p
end type
```

```
type(flow) :: f(NX,NY,NZ)
```

Which one is best?

- Both have merits
- The choice strongly depends on the computer architecture
  - for cache-based CPUs the choice is difficult (it depends on the order of the accesses of your numerical scheme)

# A Recurrent Issue: SoA or AoS



```
▶ type flow
  real(rk) :: u(NX,NY,NZ)
  real(rk) :: v(NX,NY,NZ)
  real(rk) :: w(NX,NY,NZ)
  real(rk) :: p(NX,NY,NZ)
end type
```

```
type(flow) :: f
```

Or

```
type flow
  real(rk) :: u,v,w,p
end type
```

```
type(flow) :: f(NX,NY,NZ)
```

Which one is best?

- ▶ Both have merits
- ▶ The choice strongly depends on the computer architecture
  - ▶ for cache-based CPUs the choice is difficult (it depends on the order of the accesses of your numerical scheme)
  - ▶ but using GPUs or MICs the first one is usually better!





```
subroutine my_pde_solver(nx, ny, nz)
  integer, intent(in) :: nx, ny, nz

  integer, parameter :: rk = selected_real_kind(12)
  real(rk) :: deltax, deltay, deltaz ! Grid steps

  real(rk) :: u(nx,ny,nz)
  real(rk) :: v(nx,ny,nz)
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  real(rk) :: p(nx,ny,nz)
```

- ▶ We could think of declaring automatic arrays inside a subroutine



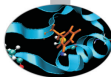
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```

- ▶ We could think of declaring automatic arrays inside a subroutine
- ▶ This is unwise
  - ▶ Automatic arrays are usually allocated on the process stack
  - ▶ Which is a precious resource
  - ▶ And limited in most system configurations

# A Bad, Old, Common approach



```
program pde_solve
  parameter (MAXNX=400, MAXNY=400, MAXNZ=400)
  parameter (MAXSIZE=MAXNX*MAXNX*MAXNZ)

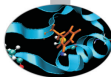
  real*8 u(MAXSIZE), v(MAXSIZE), w(MAXSIZE), p(MAXSIZE)

  common u, v, w, p
! ...
  call my_pde_solver(nx, ny, nz, u, v, w, p)
! ...
end

subroutine my_pde_solver(nx, ny, nz, u, v, w, p)
  real*8 u(nx, ny, nz), v(nx, ny, nz), w(nx, ny, nz), p(nx, ny, nz)
!...
```

- ▶ We could give a different shape to dummy arguments

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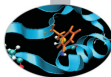
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- ▶ But this only works if interface is implicit
  - ▶ Which is dangerous

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```

- ▶ We could give a different shape to dummy arguments
- ▶ But this only works if interface is implicit
  - ▶ Which is dangerous
- ▶ Maximum problem size still program limited:  **$nx*ny*nz$**  must be less than **MAXSIZE**



- ▶ Being program limited is annoying
- ▶ It's much better to accommodate to any user specified problem size
  - ▶ Right, as long as there is enough memory
  - ▶ But if memory is not enough, not our fault
  - ▶ It's computer or user's fault
- ▶ And there are many complex kinds of computations
  - ▶ Those in which memory need cannot be foreseen in advance
  - ▶ Those in which arrays do not fit
  - ▶ Those in which very complex data structures are needed



Extending the Language

Managing Memory

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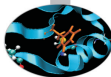
Fortran Pointers

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# Enter Allocatable Arrays



```
integer, parameter :: rk = selected_real_kind(12)

real(rk), dimension(:, :, :), allocatable :: u, v, w, p

allocate (u (nx, ny, nz), v (nx, ny, nz), w (nx, ny, nz), p (nx, ny, nz))
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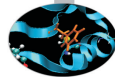


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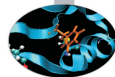
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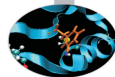
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- ▶ Best practice: use `STAT=` and, on failure, provide information to users before terminating execution



- ▶ Where all these 'dynamic allocated memory' comes from?
  - ▶ From an internal area, often termed "*memory heap*"
  - ▶ When that is exhausted, OS is asked to give the process more memory
  - ▶ And if OS is short of memory, or some configuration limit is exhausted...

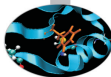


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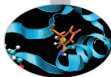
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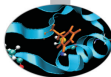


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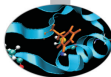
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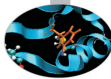
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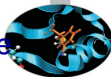


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- ▶ Mistaking allocatables for a substitute to procedure automatic arrays
  - ▶ Dynamic allocation incurs costs
  - ▶ Only worth for big arrays that would not fit program stack

## Automatic assignment (F2003)



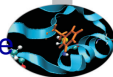
- ▶ When assigning an array value to a not allocated allocatable array, the allocatable array gets automatically allocated
- ▶ This simplifies the use of array functions which return a variable-sized result

```
real, dimension(100) :: x  
real, allocatable, dimension(:) :: nonzero_values  
nonzero_values = pack(x, x/=0)
```

- ▶ Also useful when dealing with allocatable components in a derived type



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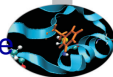


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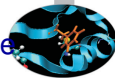
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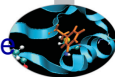
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  - ▶ if you don't need it, disable it using compiler options



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Conclusions



- ▶ Fortran pointers are aliases to other objects
- ▶ Declared like regular variables, with attribute **pointer**
- ▶ Associated to actual objects with pointer assignment =>
- ▶ To be associated with a pointer, variables must have the **target** attribute
  - ▶ But compilers are often liberal (sloppy?) on this
- ▶ Disassociated by actual objects with **nullify** statement or by pointer assignment of **null ()**

```
real, dimension(:, :, :), pointer :: r
```

```
real, target :: a(5,15,6), b(3,22,7)
```

```
r => a
```

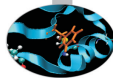
```
! now r is an alias of a  
nullify(r)
```

```
r => b
```

```
! now r is an alias of b  
r => null()
```

# More Fortran Pointers

- Pointers can point to whatever type, scalar or array



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type atom_list
  type(atom) :: a
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end type
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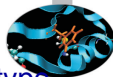
which comes in handy to define complex data structures, like lists

- ▶ Pointers may also alias subobjects

```

real, dimension(:, :, :), pointer :: r
type(velocity), pointer :: v
real, target :: a(5,15,6)
type(atom), target :: oneatom

r => a(2:4, 1:10, 3:6) ! r(1,1,1) aliases a(2,1,3)
                       ! r(3,10,4) aliases a(4,10,6)
v => oneatom%velocity
  
```



- ▶ If you allocate a pointer, an unnamed object of the pointee type is created, and associated with the pointer itself

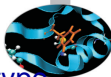
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real, dimension(:, :, :), pointer :: r  
type(atom_list), pointer :: first
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```
allocate(r(5, 15, 6))
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! now r refers an unnamed array allocated on the heap
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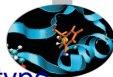
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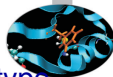
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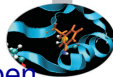
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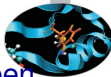
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- ▶ You can deallocate the pointee by specifying the pointer in a **`deallocate`** statement

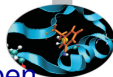


- ▶ Referencing an undefined pointer (strange things may happen, it may also seem to work)
  - ▶ Good practice: initialize pointers to `nullptr` ( )

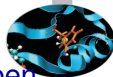


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  - ▶ This is a memory leak, and programmers causing memory leaks have really bad reputation
- ▶
 

```

real, dimension(:, :), pointer :: r, p
! ...
allocate(r(n, m))
p => r
! ...
deallocate(r)
p(k, 1) = p(k, 1) + 1
      
```

Now you'll be in troubles with `p`, with really strange behavior



► Let us clarify

```
real, dimension(:), pointer :: p
```

does not declare an array of pointers, but a pointer capable of aliasing an array

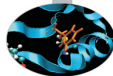


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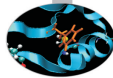
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- ▶ as such are not allowed in Fortran, but the equivalent effect can be achieved by creating a type containing a pointer component and building an array of this type



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▶ What about array of pointers?

- ▶ as such are not allowed in Fortran, but the equivalent effect can be achieved by creating a type containing a pointer component and building an array of this type

▶ For example, a lower-triangular matrix may be held using a pointer for each row

```

type row
  real, dimension(:), pointer :: r
end type row
type(row), dimension(n) :: t
do i=1,n
  allocate(t(i)%r(1:i)) ! Allocate row i of length i
enddo
  
```

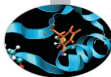


- ▶ Discretization on Cartesian 2D grid with Dirichlet Boundary Conditions

$$\begin{cases}
 f(x_{i+1,j}) + f(x_{i-1,j}) - 2f(x_{i,j}) + \\
 f(x_{i,j+1}) + f(x_{i,j-1}) - 2f(x_{i,j}) = 0 & \forall x_{i,j} \in (a,b)^2 \\
 f(x_{i,j}) = \alpha(x_{i,j}) & \forall x_{i,j} \in \partial[a,b]^2
 \end{cases}$$

- ▶ Iterative advancement using Jacobi method

$$\begin{cases}
 f_{n+1}(x_{i,j}) = \frac{1}{4} [ & f_n(x_{i+1,j}) + f_n(x_{i-1,j}) + \\
 & f_n(x_{i,j+1}) + f_n(x_{i,j-1}) ] & \forall n > 0 \\
 f_0(x_{i,j}) = 0 & \forall x_{i,j} \in (a,b)^2 \\
 f_n(x_{i,j}) = \alpha(x_{i,j}) & \forall x_{i,j} \in \partial[a,b]^2, & \forall n > 0
 \end{cases}$$



```

program laplace
  implicit none
  integer, parameter :: dp=kind(1.d0), n = 100
  integer             :: maxIter = 100000, i, j, iter = 0
  real(dp), dimension(0:n+1,0:n+1) :: T, Tnew
  real(dp)           :: tol = 1.d-4, var = 1.d0, top = 100.d0
  T(0:n,0:n) = 0.d0
  T(n+1,1:n) = (/ (i, i=1,n) /) * (top / (n+1))
  T(1:n,n+1) = (/ (i, i=1,n) /) * (top / (n+1))
  do while (var > tol .and. iter <= maxIter)
    iter = iter + 1;    var = 0.d0
    do j = 1, n
      do i = 1, n
        Tnew(i,j) = 0.25d0*( T(i-1,j) + T(i+1,j) + &
                             T(i,j-1) + T(i,j+1) )
        var = max(var, abs( Tnew(i,j) - T(i,j) ))
      end do
    end do
    if (mod(iter,100)==0) &
      write(*,"(a,i8,e12.4)") ' iter, variation:', iter, var
    T(1:n,1:n) = Tnew(1:n,1:n)
  end do
end program laplace
  
```



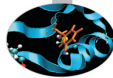


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        var = max(var, abs( Tnew(i,j) - T(i,j) ))
      end do
    end do
    if (mod(iter,100)==0) &
      write(*,"(a,i8,e12.4)") ' iter, variation:', iter, var
    T(1:n,1:n) = Tnew(1:n,1:n)
  end do
end program laplace
  
```



- ▶ Modify the code using advanced Fortran features:
  - ▶ array syntax
  - ▶ allocatable arrays
  - ▶ pointer arrays
  
- ▶ Try to list pros and cons of each approach



- ▶ Re-write a program that:
  - ▶ reads an 'arbitrarily' long column of real numbers from an ASCII file
  - ▶ copy the values to a suitably allocated array
  - ▶ prints maximum, minimum, average of the numbers
  - ▶ and prints the  $\lfloor n/2 \rfloor$ -th row where  $n$  is the length of the columnavoiding the constraint of the declaration of a static maximal array, e.g. `a(1000)`
- ▶ Use allocatables: you need to read the file twice
- ▶ Use pointers, first store the values in a linked list and then copy them to an allocatable array
  - ▶ How to spare memory?



Extending the Language

**Managing Memory**

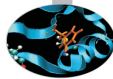
Dynamic Memory Allocation

Fortran Pointers

**Bridging the Gap with C**

Sketchy Ideas on Data Structures

Conclusions



- ▶ You may want to call a C function from a Fortran program
- ▶ Or call a Fortran procedure from a C program
- ▶ And you don't want to translate and re-debug
- ▶ Or you can't, as you don't have sources
  
- ▶ You may also want to share global data among C and Fortran program units
  
- ▶ This has been done in the past with non-standard tricks
- ▶ Fortran 2003 offers a better, standard way
- ▶ Let's look at it in steps

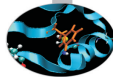


- Imagine you have this C function:

```
double avg_var(int n, const double a[], double *var) {  
    double avg = 0.0;  
    double avg2 = 0.0;  
    for(int i=0;i<n;i++) {  
        avg += a[i];  
        avg2 += a[i]*a[i];  
    }  
    avg = avg/n;  
    *var = avg2/n - avg*avg;  
    return avg;  
}
```

and you want to call it from your Fortran code like:

```
avg = avg_var(m,b,var)
```



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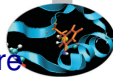
- Ore you have your favorite, thoroughly tested Poisson solver:

```

interface
    subroutine myPoissonSolver(l, m, n, f)
        integer, intent(in) :: l, m, n
        real(kind(1.0D0)), intent(inout) :: f(l,m,n)
    end subroutine myPoissonSolver
end interface
  
```

and you want to call it from your C code like:

```
myPoissonSolver(nx, ny, nz, field);
```



- ▶ We could think that Fortran interfaces and C declarations are enough

- ▶ And write, to call C from Fortran:

```

interface
  function avg_var(n, a, var)
    integer, intent(in) :: n
    real(kind(1.0D0)), intent(in) :: a(*)
    real(kind(1.0D0)), intent(out) :: var
    real(kind(1.0D0)) :: avg_var
  end function avg_var
end interface
  
```

- ▶ And to call Fortran from C, add on Fortran side:

```

interface
  subroutine myPoissonSolver(l, m, n, f)
    integer, intent(in) :: l, m, n
    real(kind(1.0D0)), intent(inout) :: f(l,m,n)
  end subroutine myPoissonSolver
end interface
  
```

and on the C side, the declaration:

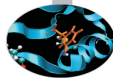
```
void myPoissonSolver(int nx, int ny, int nz, field[nz][ny][nx]);
```

- ▶ This is the right track, but still half way from our destination



## Thou Shalt Not Mangle Names

- ▶ Fortran compilers mangle procedure names
  - ▶ All uppercase or all lowercase
  - ▶ Compilers may append/prepend one or two `_` characters
  - ▶ And for module procedures is even worse
  - ▶ Used to be sorted out on the C side, in non-portable ways



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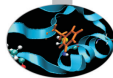


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- ▶ Enter Fortran 2003 `bind` attribute



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- ▶ Enter Fortran 2003 `bind` attribute
- ▶ For C to Fortran:

```
interface
  function avg_var(n, a, var) bind(c)
    integer, intent(in) :: n
    real(kind(1.0D0)), intent(in) :: a(*)
    real(kind(1.0D0)), intent(out) :: var
    real(kind(1.0D0)) :: avg_var
  end function avg_var
end interface
```



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  end function avg_var
end interface
  
```

▶ For Fortran to C, Fortran side:

```

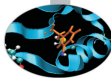
interface
  subroutine myPoissonSolver(l, m, n, f) bind(c)
    integer, intent(in) :: l, m, n
    real(kind(1.0D0)), intent(inout) :: f(l,m,n)
  end subroutine myPoissonSolver
end interface
  
```

and on the C side, the declaration:

```
void myPoissonSolver(int nx, int ny, int nz, field[nz][ny][nx]);
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## Thou Shalt Care for Argument Passing

- ▶ Fortran passes arguments by reference
  - ▶ Under the hood, it's like a C pointer
  - ▶ Works for C arrays and pointers to scalar variables
  - ▶ But usually scalars are passed by value in C



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```
interface
  function avg_var(n, a, var) bind(c)
    integer, value :: n
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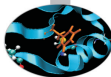
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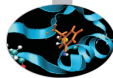
```
void myPoissonSolver(int nx, int ny, int nz, field[nz][ny][nx]);
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  - ▶ Implementations have a lot of freedom
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- ▶ For C to Fortran:

```
interface
  function avg_var(n, a, var) bind(c)
    use iso_c_binding
    integer(c_int), value :: n
    real(c_double), intent(in) :: a(*)
    real(c_double), intent(out) :: var
    real(c_double) :: avg_var
  end function avg_var
end interface
```



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- ▶ For Fortran to C, Fortran side:

```

interface
  subroutine myPoissonSolver(l, m, n, f) bind(c)
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end interface
  
```

and on the C side, still the declaration:

```
void myPoissonSolver(int nx, int ny, int nz, field[nz][ny][nx]);
```

## More from `iso_c_binding`



- ▶ `iso_c_binding` defines named constants holding kind type parameter values for intrinsic types for the platform
- ▶ `integer(c_int)` is the kind value corresponding to a C `int`
- ▶ Negative values are used for unsupported C types, so the compiler will flag the problem

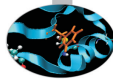
- ▶ A few of them:

Type	Kind	C type
<code>integer</code>	<code>c_int</code>	<code>int</code>
	<code>c_short</code>	<code>short int</code>
<code>real</code>	<code>c_float</code>	<code>float</code>
	<code>c_double</code>	<code>double</code>
<code>complex</code>	<code>c_float_complex</code>	<code>float _Complex</code>
	<code>c_double_complex</code>	<code>double _Complex</code>
<code>logical</code>	<code>c_bool</code>	<code>_Bool</code>
<code>character</code>	<code>c_char</code>	<code>char</code>

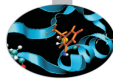
- ▶ Fortran 2008 adds `c_sizeof()`, check with your compiler!



- ▶ Fortran has multidimensional arrays
- ▶ C has arrays of arrays (of arrays...)
- ▶ Thus the mapping of array indexes to actual data layout in memory is inverted
  - ▶ Fortran array  $\mathbf{a}(\mathbf{L}, \mathbf{M}, \mathbf{N})$
  - ▶ maps to C array  $\mathbf{a}[\mathbf{N}][\mathbf{M}][\mathbf{L}]$



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- ▶ Before C99, the leading dimension of an array function parameter could not be specified in C
  - ▶ C array parameter  $\mathbf{a}[\ ]$
  - ▶ maps to Fortran assumed size array parameter  $\mathbf{a}(\star)$



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  - ▶ Fortran array **a (L, M, N)**
  - ▶ maps to C array **a [N] [M] [L]**
- ▶ Before C99, the leading dimension of an array function parameter could not be specified in C
  - ▶ C array parameter **a []**
  - ▶ maps to Fortran assumed size array parameter **a (\*)**
- ▶ In C99, Variable Length Arrays were introduced
  - ▶ C99 array parameter **a [nz] [ny] [nx]**
  - ▶ maps to Fortran array parameter **a (nx, ny, nz)**





- ▶ **bind** also helps for derived types and global data
- ▶ For derived types, each component must be interoperable

Fortran

```

type, bind(c) :: particle
  integer(c_int) :: n
  real(c_float) :: x,y,z
  real(c_float) :: vx,vy,vz
end type particle
  
```

C

```

typedef struct particle {
  int n;
  float x,y,z;
  float vx,vy,vz;
} particle;
  
```

- ▶ For module variables or common blocks, use

Fortran

```

integer(c_long), bind(c) :: n

real(c_double) :: m,k
common /com_mk/ m,k
bind(c) :: /com_mk/
  
```

C

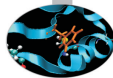
```

extern long n;

extern struct mk {
  double m, k;
} com_mk;
  
```

- ▶ Note: common blocks become C **structs**

# Fortran Pointers vs. C Pointers



- ▶ As of argument passing, not a problem
- ▶ But Fortran pointers are not interoperable with C
- ▶ Fortran pointers sport richer semantics, notably:
  - ▶ multidimensional arrays
  - ▶ non-contiguous memory areas
- ▶ C functions returning a pointer must have `type(c_ptr)` type (from `iso_c_binding`)
- ▶ Ditto for C pointer variables and pointer members of C `structs`:

## Fortran

```

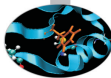
type, bind(c) :: block
  integer(c_int) :: n_neighbors
  type(c_ptr) :: neighbors
  type(c_ptr) :: grid
end type block
  
```

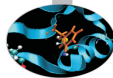
## C

```

typedef struct {
  int n_neighbors;
  int *neighbors;
  mesh *grid;
} block;
  
```

- ▶ **iso\_c\_binding** module provides much needed help

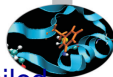




- ▶ `iso_c_binding` module provides much needed help
- ▶ `c_loc(x)` returns a valid C pointer to the content of variable `x`
- ▶ `c_f_pointer(cptr, fptr[, shape])` performs the opposite translation, writing the result in the Fortran pointer `fptr`
  - ▶ An optional `shape` argument like `(/n/)` or `(/1,m,n/)` gives it a shape for array pointers



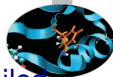
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  - ▶ An optional `shape` argument like `(/n/)` or `(/1,m,n/)` gives it a shape for array pointers
- ▶ If `f_proc` is an interoperable Fortran procedure, `c_funloc(f_proc)` returns a valid C pointer (`type(c_funptr)`) to it
- ▶ `c_f_procpointer(cfptr, fpptr)` performs the opposite translation, writing the result in the Fortran procedure pointer `fpptr`



- ▶ Obviously, C and Fortran sources must be separately compiled and then linked

```
user@cineca$> gcc -c fun_cmd.c
user@cineca$> gfortran -c main_cmd.f90
user@cineca$> gfortran fun_cmd.o main_cmd.o -o main_cmd
```

- ▶ Easy, if calling C functions from a Fortran program
  - ▶ Fortran Runtime Library is usually built on top of C one



- ▶ Obviously, C and Fortran sources must be separately compiled and then linked

```
user@cineca$> gcc -c fun_cmd.c
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```

- ▶ Easy, if calling C functions from a Fortran program
  - ▶ Fortran Runtime Library is usually built on top of C one
- ▶ Less so if calling Fortran procedures from a C program
  - ▶ Fortran compiler might insert calls to its Runtime Library
- ▶ Best practice:

```
user@cineca$> gcc -lgfortran procedures.o main.c
```

- ▶ Your mileage may vary, browse your compiler manuals



Extending the Language

**Managing Memory**

Dynamic Memory Allocation

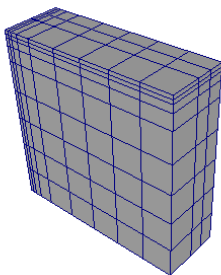
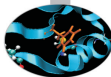
Fortran Pointers

Bridging the Gap with C

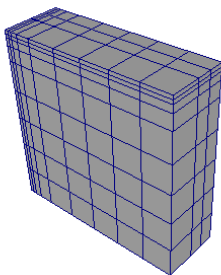
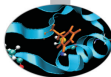
**Sketchy Ideas on Data Structures**

Conclusions





- ▶ Let's imagine we have to solve a PDE



- ▶ Let's imagine we have to solve a PDE
- ▶ On a dense, Cartesian, non uniform grid
  - ▶ Mesh axes are parallel to coordinate ones
  - ▶ Steps along each direction differ in size from point to point



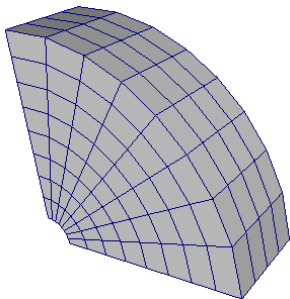
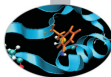
```
type nonuniform_grid
  integer :: nx, ny, nz
  ! Grid steps
  real(rk), dimension(:), allocatable :: deltax
  real(rk), dimension(:), allocatable :: deltay
  real(rk), dimension(:), allocatable :: deltaz
end type
!...
type(nonuniform_grid) :: my_grid
integer :: alloc_stat
!...
allocate(my_grid%deltax(nx),my_grid%deltay(ny), &
         my_grid%deltaz(nz), STAT=alloc_stat)
if (alloc_stat > 0) then
  ! graceful failure
end if
```

- ▶ Related information is best kept together

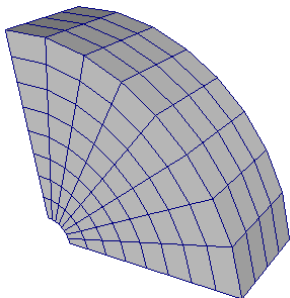
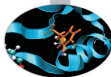


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```

- ▶ Related information is best kept together
- ▶ Grid size and grid steps are related information



- ▶ Let's imagine we have to solve a PDE



- ▶ Let's imagine we have to solve a PDE
- ▶ On a dense structured mesh
  - ▶ Could be continuously morphed to a Cartesian grid
  - ▶ Need to know coordinates of each mesh point

# Sketching a Mesh Description



```
type meshpoint
  real(rk) :: x, y, z
end type

type, extends(meshpoint) :: normal
end type

type mesh
  integer :: nx, ny, nz

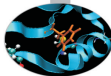
  type(meshpoint), dimension(:,:,:), allocatable :: coords

  type(normal), dimension(:,:,:), allocatable :: xnormals
  type(normal), dimension(:,:,:), allocatable :: ynormals
  type(normal), dimension(:,:,:), allocatable :: znormals

  real(rk), dimension(:,:,:), allocatable :: volumes
end type
!...
type(mesh) :: my_mesh

! allocate my_mesh components with extents nx, ny, nz
! immediately checking for failures!
```

# A Recurrent Issue, Again



```
► real(rk) :: x(NX,NY,NZ)
   real(rk) :: y(NX,NY,NZ)
   real(rk) :: z(NX,NY,NZ)
```

Or

```
type meshpoint
  real(rk) :: x, y, z
end type
```

```
type(meshpoint), dimension(NX,NY,NZ) :: coords
```

Which one is best?



# A Recurrent Issue, Again



- ▶ `real(rk) :: x(NX,NY,NZ)`  
`real(rk) :: y(NX,NY,NZ)`  
`real(rk) :: z(NX,NY,NZ)`

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```

```
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Which one is best?

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  - ▶ The former (if done properly) allows hardware to play efficient tricks in memory accesses
  - ▶ The latter brings in cache all values related to a grid point as soon as one component is accessed

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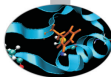
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- ▶ Again, both have merits
  - ▶ The former (if done properly) allows hardware to play efficient tricks in memory accesses
  - ▶ The latter brings in cache all values related to a grid point as soon as one component is accessed
- ▶ Here, we lean to the latter
  - ▶ As in most numerical schemes,  $x$ ,  $y$ , and  $z$  components of the same mesh point are accessed together

# Multiblock Meshes and More



- ▶ A multiblock mesh is an assembly of connected structured meshes



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  - ▶ You could dynamically allocate a **mesh** array
  - ▶ Or build a **block** type including a **mesh** and connectivity information



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- ▶ Adaptive Mesh Refinement





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  - ▶ You want your blocks resolution to adapt to dynamical behavior of PDE solution
  - ▶ Which means splitting blocks to substitute part of them with more resolved meshes



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  - ▶ You could dynamically allocate a **mesh** array
  - ▶ Or build a **block** type including a **mesh** and connectivity information
- ▶ Adaptive Mesh Refinement
  - ▶ You want your blocks resolution to adapt to dynamical behavior of PDE solution
  - ▶ Which means splitting blocks to substitute part of them with more resolved meshes
- ▶ Eventually, you'll need more advanced data structures
  - ▶ Like lists
  - ▶ Like binary trees, oct-trees, n-ary trees

# If You Read Code Like This...



```
type block_item
  type(block), pointer :: this_block

  type(block_item), pointer :: next
end type

!...
do while (associated(p))
  call advance_block_in_time(p%this_block)
  p => p%next
end do
```

# If You Read Code Like This...



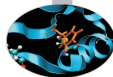
```
type block_item
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  type(block_item), pointer :: next
end type

!...
do while (associated(p))
  call advance_block_in_time(p%this_block)
  p => p%next
end do
```

- ▶ It is processing a singly-linked list of mesh blocks
- ▶ You need to learn more on abstract data structures
- ▶ Don't be afraid, it's not that difficult

# And If You Read Code Like This...



```

type block_tree_node
  type(block), pointer :: this_block

  integer :: children_no
  type(block_tree_node), pointer :: childrens

  type(block_tree_node), pointer :: next_sibling
end type

!...
recursive subroutine tree_advance_in_time(n)
  type(block_tree_node) :: n
  type(block_tree_node), pointer :: p
  integer :: i

  p => n%childrens
  do i=0,n%children_no
    call tree_advance_in_time(p)
    p => p%next_sibling
  end do

  call advance_block_in_time(n%this_block)
end subroutine tree_advance_in_time

```

# And If You Read Code Like This...



```
type block_tree_node
  type(block), pointer :: this_block

  integer :: children_no
  type(block_tree_node), pointer :: childrens

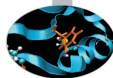
  type(block_tree_node), pointer :: next_sibling
end type

!...
recursive subroutine tree_advance_in_time(n)
  type(block_tree_node) :: n
  type(block_tree_node), pointer :: p
  integer :: i

  p => n%childrens
  do i=0,n%children_no
    call tree_advance_in_time(p)
    p => p%next_sibling
  end do

  call advance_block_in_time(n%this_block)
end subroutine tree_advance_in_time
```

- ▶ It is processing a tree of mesh blocks (AMR, probably)
- ▶ You need to learn more on abstract data structures
- ▶ Don't be afraid, it's not that difficult



Extending the Language

Managing Memory

Conclusions





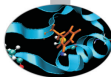
- ▶ More Fortran practice
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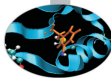
- ▶ More Fortran practice
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- ▶ More Fortran practice
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- ▶ More Fortran
  - ▶ Full object oriented programming
  - ▶ Floating point environment
  - ▶ Direct I/O
  - ▶ Asynchronous I/O
  - ▶ Submodules
  - ▶ Even more format edit descriptors
  - ▶ A few more statements and quite a few intrinsics



-  J3 US Fortran Standards Committee  
<http://www.j3-fortran.org/>
-  ISO WG5 Committee  
<http://www.nag.co.uk/sc22wg5/>
-  Fortran 2003 Standard Final Draft  
Search Internet for `n3661.pdf`
-  Fortran Wiki  
<http://fortranwiki.org/>
-  M. Metcalf, J. Reid, M. Cohen  
*Fortran 95/2003 Explained*  
Oxford University Press, corrected ed., 2008
-  M. Metcalf, J. Reid, M. Cohen  
*Modern Fortran Explained*  
Oxford University Press, 2011
-  S. Chapman  
*Fortran 95/2003 for Scientists & Engineers*  
McGraw-Hill, 3d ed., 2007
-  Adams, J.C., Brainerd, W.S., Hendrickson, R.A., Maine, R.E., Martin, J.T., Smith, B.T.  
*The Fortran 2003 Handbook*  
Springer, 2009



Salvatore Filippone's Home Page

[www.ce.uniroma2.it/people/filippone.html](http://www.ce.uniroma2.it/people/filippone.html)



Parallel Sparse Basic Linear Algebra Subroutines

[www.ce.uniroma2.it/psblas/index.html](http://www.ce.uniroma2.it/psblas/index.html)



Numerical Engine (for) Multiphysics Operators

[www.ce.uniroma2.it/nemo/index.html](http://www.ce.uniroma2.it/nemo/index.html)



Portable Fortran Interfaces to the Trilinos C++ Package

[trilinos.sandia.gov/packages/fortrilinos/](http://trilinos.sandia.gov/packages/fortrilinos/)



Stefano Toninel

*Development of a New Parallel Code for Computational Continuum Mechanics  
Using Object-Oriented Techniques*

PhD Thesis, University of Bologna, 2006.



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